**IEEE P802.15**

**Wireless Personal Area Networks**

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| Re: |  |
| Abstract | Informative text on BCH Code Construction for Information Reconciliation  |
| Purpose | Approval |
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# Annex E

(informative)

**E.1 Design procedure of a primitive *(n,k,t)* BCH code**

* Determine field size *m* where *n = 2m -* 1
* Determine error correction capability *t*
* Find a primitive *n*-th root α of unity in a field GF(2*m*)
* Select 2*t* consecutive powers of α, starting with for some non-negative integer

For uniqueness, , (i.e. narrow sense)

* Let be the least common multiple of the minimal polynomials for the selected powers of with respect to . (Each of the minimal polynomials should appear only once in the product)

where is the minimal polynomial for .

**E.2 Example of generation of parity bits of BCH Code**

Assume *n=*15 and *t=*2. Then, to generate a narrow sense primitive BCH code with *n=*15 and *t=*2,

1. Find , the root of , i.e. .
2. Select consecutive powers of starting with
3. Calculate . The minimal polynomials for are shown in Table 1. (See E.3 for how to construct minimal polynomials.)

Table —Minimal polynomials for

|  |  |
| --- | --- |
| Power of root,  | Minimal polynomial,  |
|  |  |
|  |  |
|  |  |
|  |  |
|  |  |

1. From Table 1, , since . By substituting and , we have

To obtain the parity bits of BCH coding

1. Multiply the message polynomial by .
2. Divide the result by the generator polynomial .
3. Let be the remainder, then is the parity for .

In case of (15,7,2) BCH code, assuming the message bits are [1 0 1 1 0 1 1],

and since

we have

From , the parity bits are [0 1 1 0 1 1 0 1].

**E.3 Generation of minimal polynomials**

Assume GF(23) = GF(8).

**E.3.1 Cyclotomic coset**

Cyclotomic cosets on are sets for all . Assuming , .

* If , then .
* If , then since in .
* If , then since in .

**E.3.2 Minimal polynomial**

Minimal polynomial for each cyclotomic coset on is defined as

where is the cyclotomic coset on which includes . For example, , the minimal polynomial for each cyclotomic coset is as follows.

* Minimal polynomial for :
* Minimal polynomial for :
* Minimal polynomial for :

Table —Elements of GF(23) with primitive polynomial .

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| Since , (negative sign is ignored on binary calculation).From , ,,since 2 equals to 0 on binary calculation |