

Project: IEEE P802.15 Working Group for Wireless Personal Area Networks (WPANs)

Submission Title: [Samsung Pre-proposal to TG8 CFC: Overall PAC Procedures]

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Source: [Seung-Hoon Park, Kyungkyu Kim, Anil Agiwal, Youngbin Chang, Hyunseok Ryu, Daegyun Kim and Won-il Roh]

Company [Samsung Electronics]

Address [416, Maetan-3Dong, Yeongtong-Gu, Suwon-Si, Gyeonggi-Do, 443-742, Korea]

Voice:[+82-10-9349-9845]¹, **FAX:** [+82-31-279-0813]¹, **E-Mail:**[shannon.park@samsung.com]

Re: [.]

Abstract: [Presentation of PAC procedures to meet functional requirements including identified features from PFD]

Purpose: [Corresponding to Call for Contribution]

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Samsung Pre-proposal to CFC: Overall PAC Procedures

March, 2014

Samsung

1. Overall

- High Competition for Next Big Trend
 - Proximity-based Service

| Title | Company/ Organization | Note |
|------------------|-----------------------|--------------------------------|
| iBeacon | Apple | Launched in December 2013 |
| LTE Direct | 3GPP | Specified until September 2014 |
| NAN [†] | WiFi Alliance | Spec. 1.0 |

[†]NAN(Neighbor Awareness Networking)

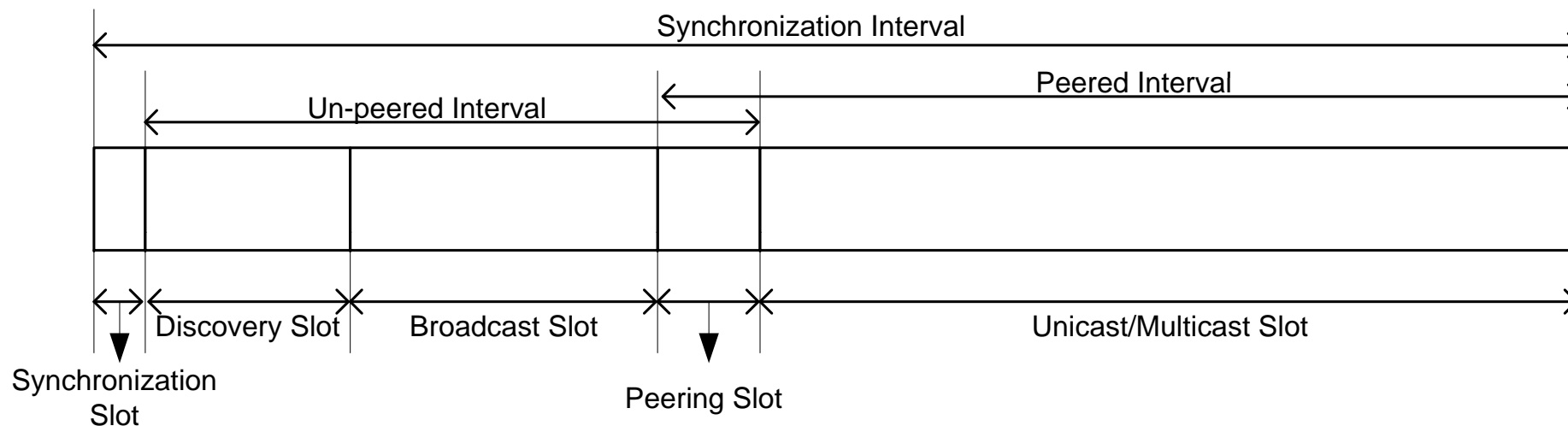
- PAC Competitive Features
 - Low power consumption for peer discovery
 - Large number of detected discovery information
 - Higher throughput for communication

1.1. PAC Procedures

- **Synchronization**
 - Based on Pulse-Coupled Oscillator (PCO) algorithm
- **Peer Discovery**
 - Broadcast Discovery Information via selected resource
- **Broadcast**
 - Broadcast data traffic without peering
- **Peering**
 - Link establishment for unicast and multicast links
- **Unicast/Multicast**
 - Request/Response-based resource assignment

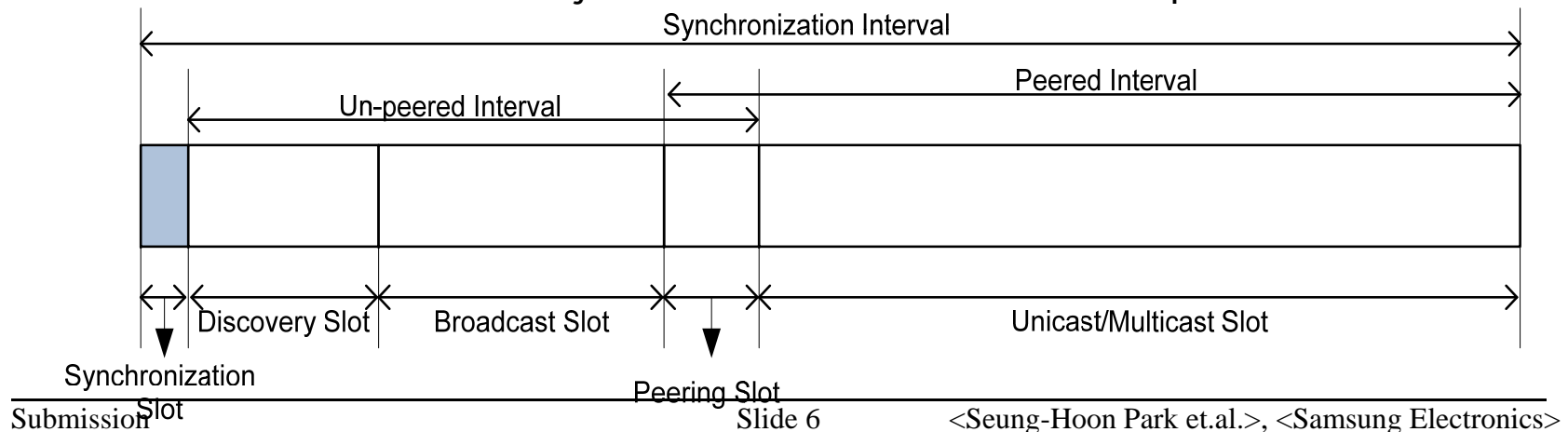
1.2 Frame Structure

- Un-peered Interval
 - Synchronization / Peer discovery / Broadcast / Peering
- Peered Interval
 - Peering / Unicast / Multicast



2. Synchronization

- Synchronization for Scalable Network
 - Distributed synchronization
 - Master-slave synchronization should be avoided
 - PDs between two different synchronized group happen
 - It is matched well to flat architecture (no hierarchy)
 - Synchronization should be done before peer discovery
 - Peer discovery prior to link connection (peering)
 - Broadcast-based synchronization mechanism is required



2. Synchronization

- Proposal
 - Physical layer signaling based synchronization
 - Details referred from 15-13-0376-01
- Synchronization Signal (SS)
 - ZC-sequence is suggested
 - Low PAPR with high detection probability
- Two Type of SS
 - Type-1 SS: for initial synchronization
 - PD performs just synchronization without frame structure
 - Short SS interval: e.g. 10 ms
 - Type-2 SS: for maintaining synchronization
 - PD follows operations defined in frame structure
 - Long SS interval: e.g. 1000ms SS

2. Synchronization

- Technical Issues
 - Low detection probability in a high dense environment

- Proposal 1
 - A PD has triggering condition to take a role of SS transmitter
 - e.g.
 - 1) When the PD is a discovery information transmitter
 - 2) Based on the detected number of SS transmitter

- Proposal 2
 - SS transmitter may transmit SS with different offset in the synchronization interval

3. Peer Discovery

■ Design Considerations

- Discovery Information (DI)

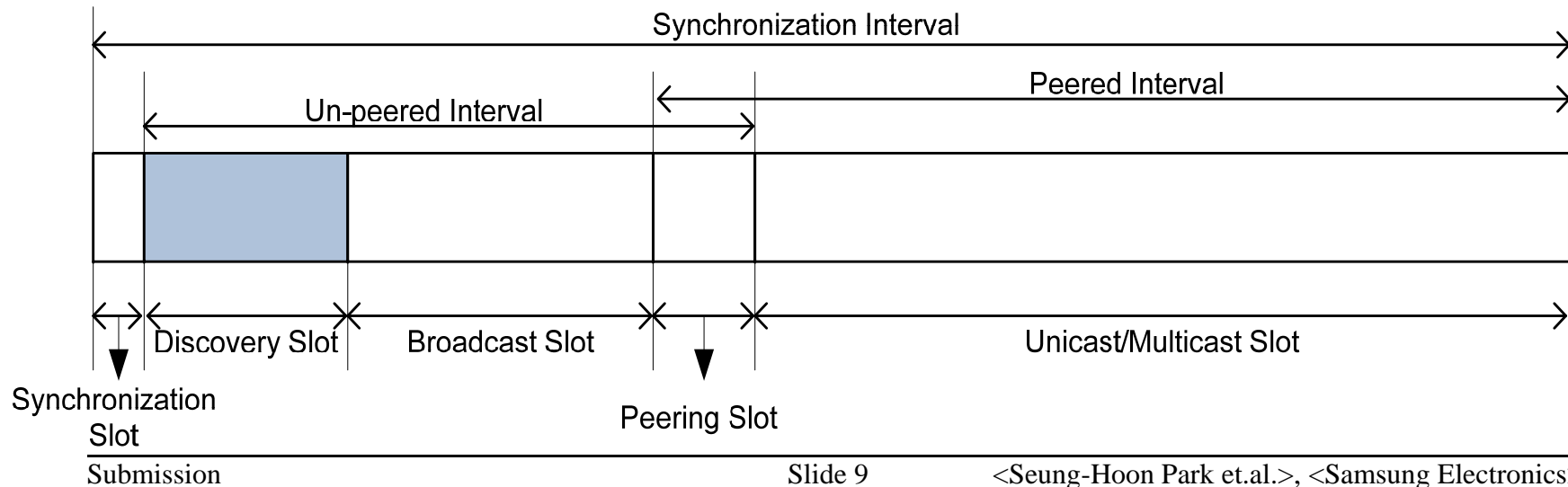
- Came from application or middleware

- Plain DI: Application ID (PACbook), or User ID (Bob@PACbook), or etc

- Coded DI: generated by middleware or retrieved from server

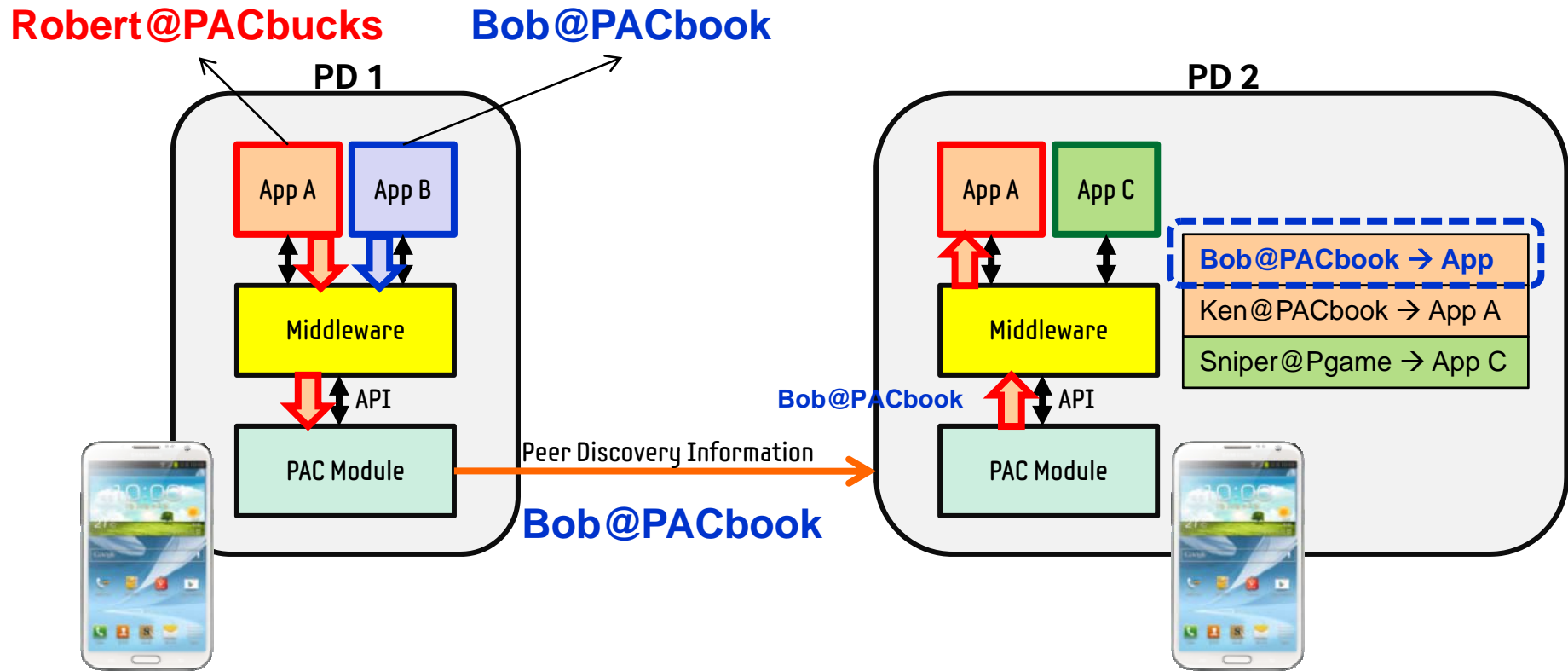
- Discovery matching

- PD A is matched by other PDs storing DI representing PD A

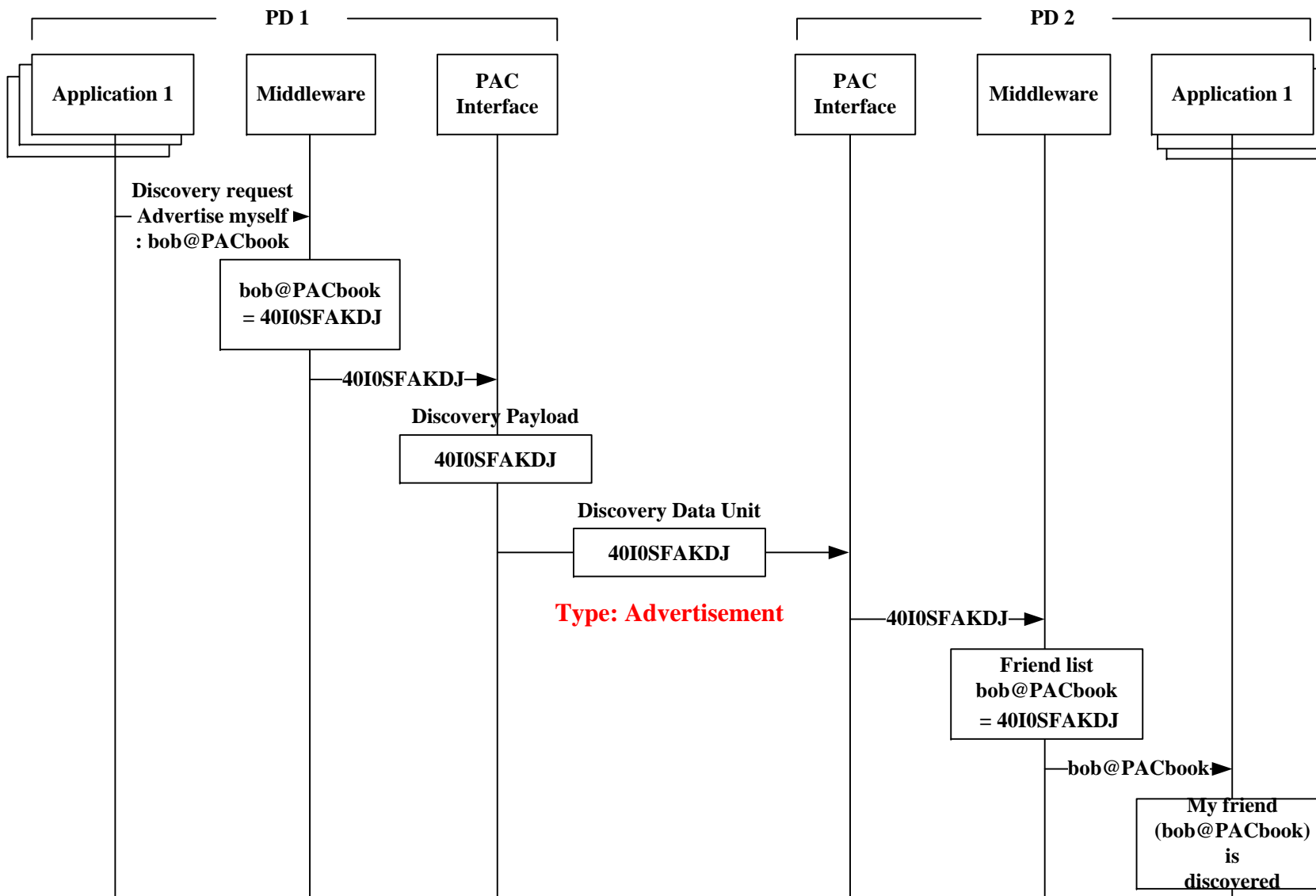


3. Peer Discovery

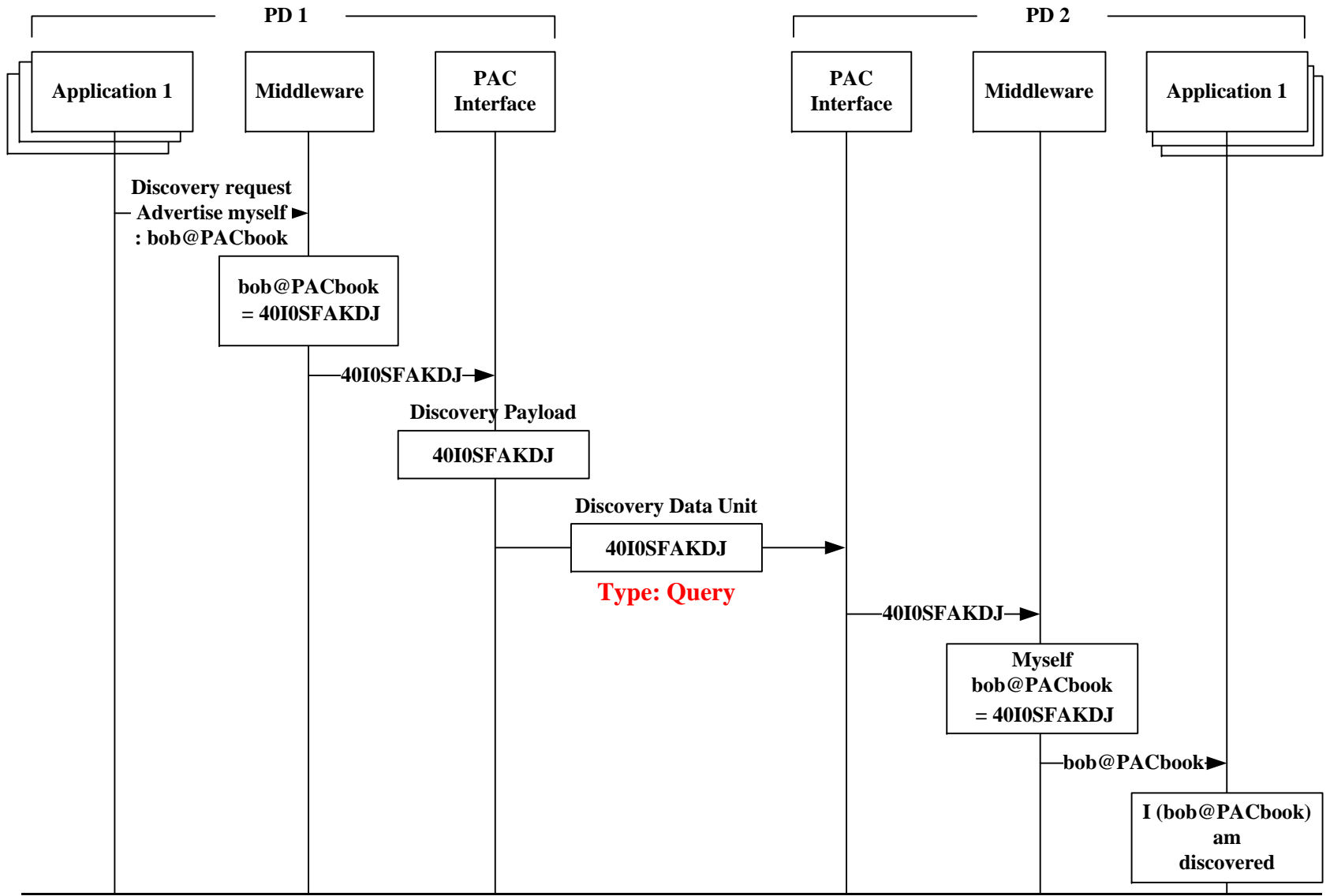
- What is Peer Discovery?
 - A peer represents an application-specific entity, not a device



Advertisement/Publish Scenario



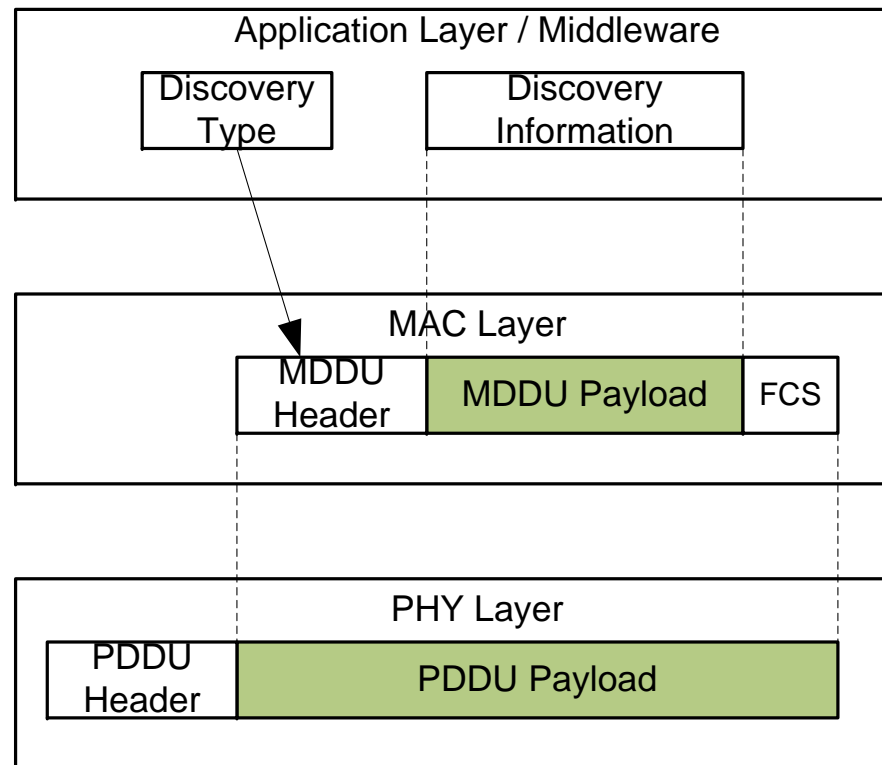
Query Scenario



3. Peer Discovery

- Protocol Stack for Peer Discovery
 - MAC/PHY Discovery Data Unit (MDDU/PPDU)

- Discovery Type
- Advertisement
 - Publish/Subscribe
 - Query/Reply



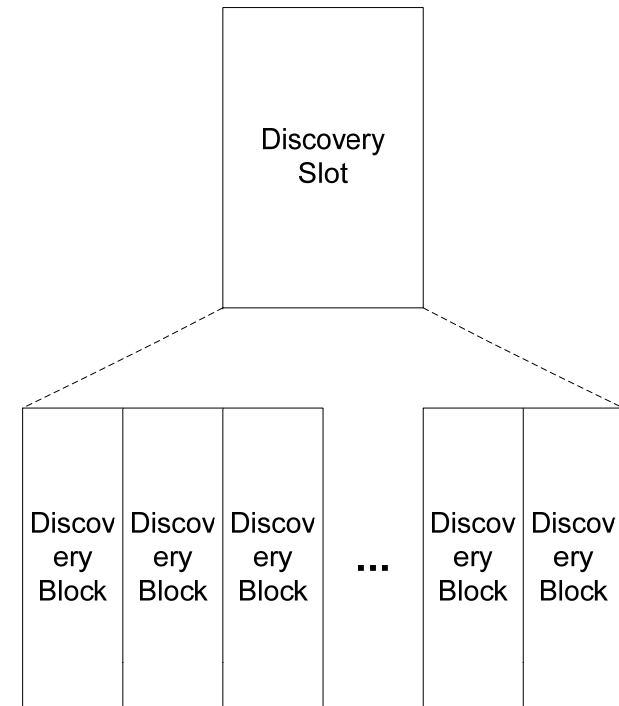
3. Peer Discovery

■ Discovery Slot

- It is comprised of multiple Discovery Blocks (DBs)
- Basic procedure
 - PD selects one DB in a Discovery Slot
 - PD broadcasts DDU (Discovery Data Unit)
 - At selected DB

■ Proposal

- Congestion-aware DB selection
- Hashing-based DB selection



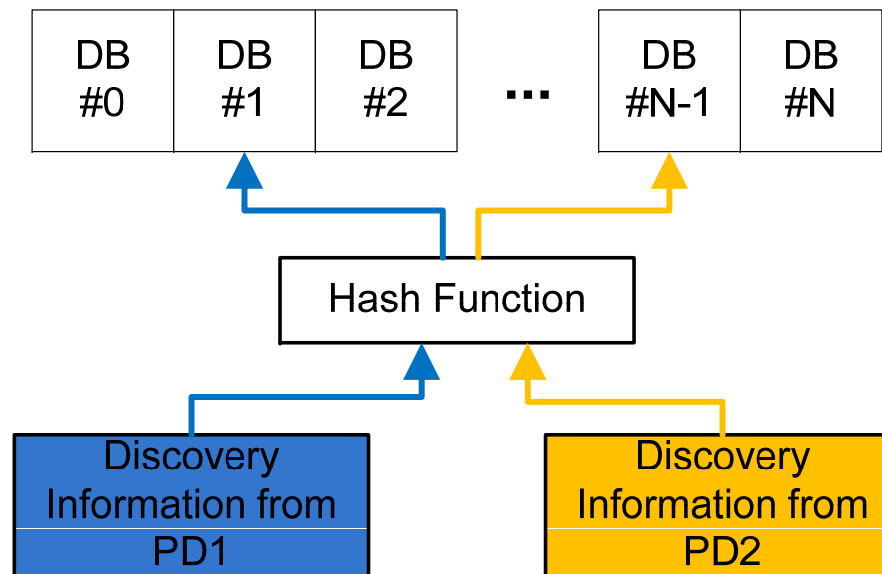
3. Peer Discovery

- Congestion-aware DB selection
 - Based on the received power
 - Compare received power between current DB and candidate DB
 - Details referred from 15-13-0376-01
 - Based on congestion condition
 - Discovery Transmission Interval (DTI) control
 - Depending on the number of detected DDU
 - e.g.) DTI is increased when the number of detected DDU is high

3. Peer Discovery

■ Hashing-based DB Selection

- DI index is determined based on the hashed DI
 - A receiving PD monitors only DB with hashed index based on monitoring DI
 - Benefit for receiver PD in power consumption perspective



3. Peer Discovery

- Technical Issues

- Support of various length of discovery information
- Subject to provide low power discovery

- Background

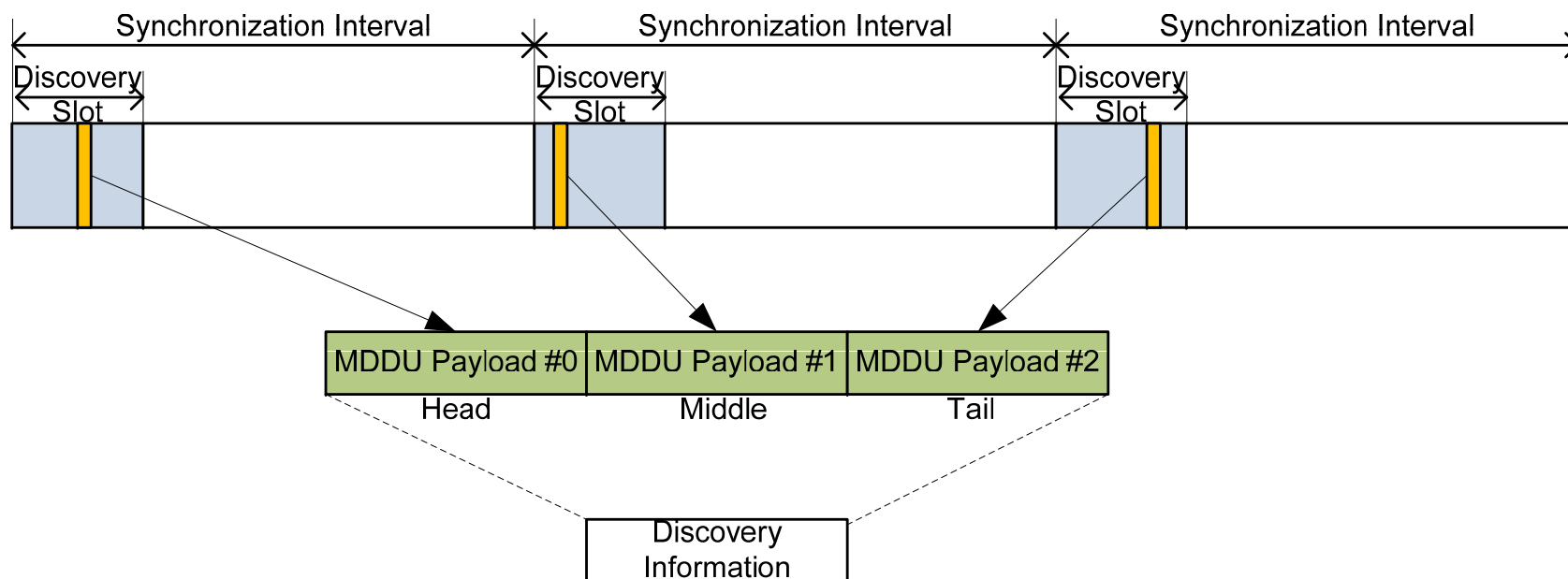
- Current distributed middleware platform in the industry

| Platform | Feature | Required size of discovery information |
|----------|---|--|
| XMPP[1] | Decentralized protocol for instant messaging and presence | Up to 3071 bytes |
| UUID[2] | Distributed systems to uniquely identify information without significant central coordination | 16 bytes |

3. Peer Discovery

■ Proposal

- Discovery information fragmentation
 - to support long discovery information
 - Discovery Information is divided into multiple MDDU payloads



3. Peer Discovery

- Design Consideration to link MDDU Payloads
 - MDDU payload size is defined to fit minimum length of DI
 - 16 bytes (e.g. UUID)
 - TX ID (MAC Address) is large compared to the length of MDDU
 - MAC address: 6 bytes
 - The size of DI: from 16 to 3071 bytes
 - Overhead analysis when using TX ID
 - To support 3071 bytes fragmentation, DI is fragmented to 192 MDDU payloads
 - TX ID overhead: $6 \times 192 = 1152$ bytes → **37 % overhead!**

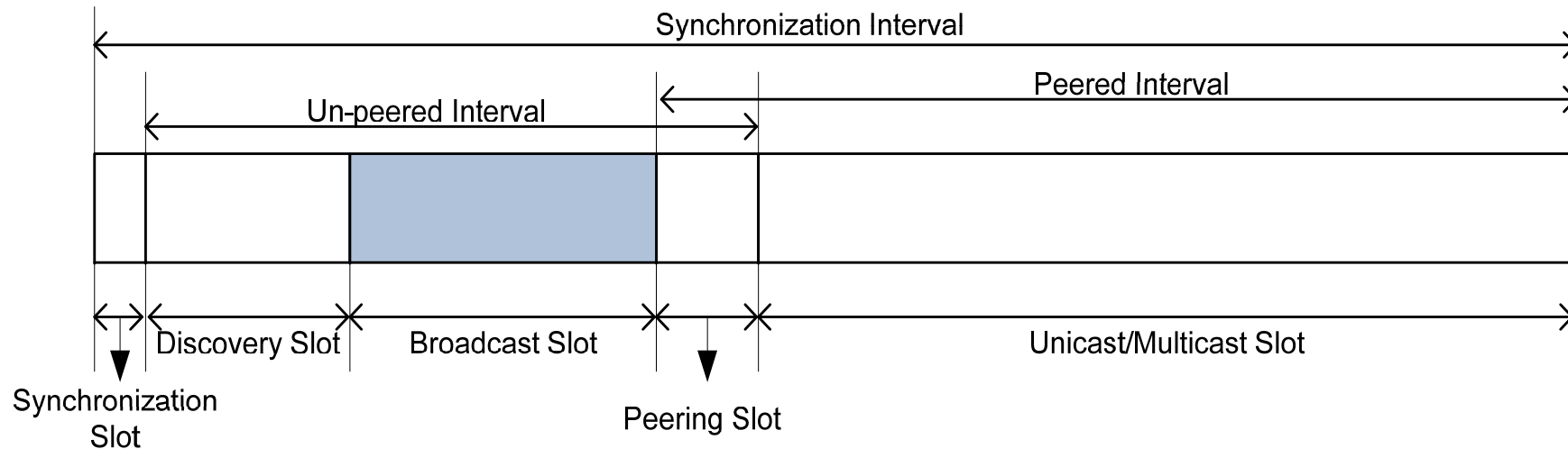
3. Peer Discovery

- Possible Solution to Link MDDU Payloads
 - Low overheads required
 - Option 1
 - Discovery Session ID (DSID)
 - to distinguish different DIs from different PDs as well as same PD
 - Potential problem
 - Who coordinates and assign DSID?
 - How to guarantee collision avoidance when using small length of DSID?
 - Option 2
 - Location Indicator
 - Each MDDU indicates the location of next MDDU
 - Receiving PD can aggregate based on the known location

4. Broadcast

■ Features of Broadcast

- Broadcast data transmission within un-peered PDs
- Contention-based access scheme
- No multi-hop allowed
 - (because it happens before peering for authentication)

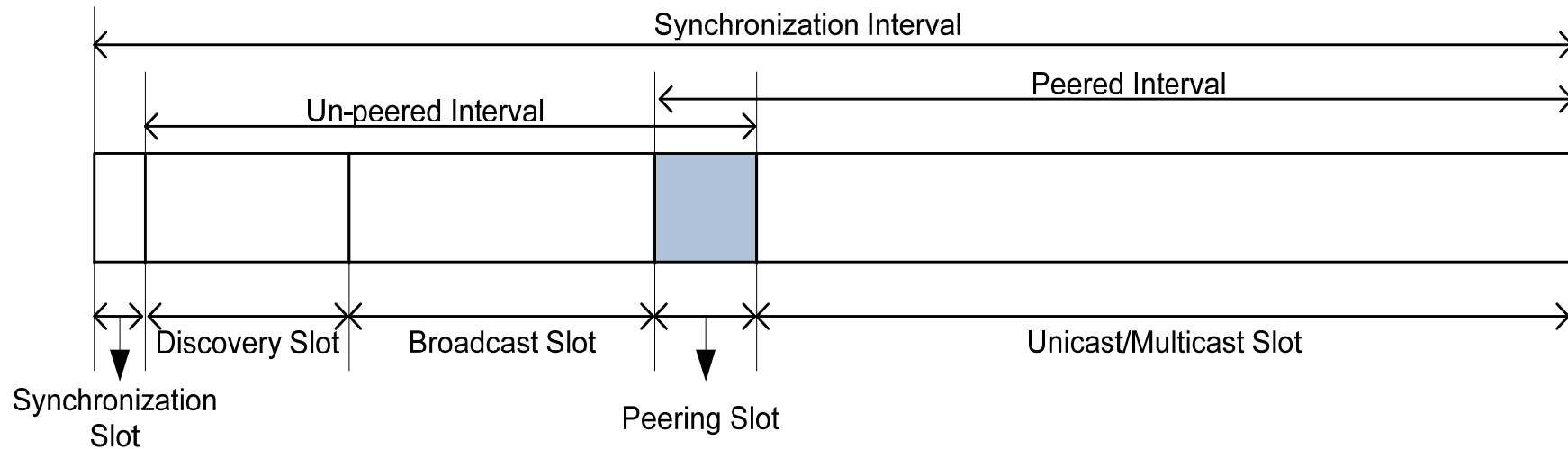


5. Peering

- The role of Peering
 - A procedure to connect to discovered peer
 - Triggered by application automatically or by user manually
 - **No MAC-level triggering**
 - Link establishment for unicast/multicast link
 - Between a TX PD and RX PD(s)
 - Exchange of information for setup
 - TX/RX ID (MAC address), capability, or etc
 - Determine link related parameters
 - Link ID, QoS class, link range, or etc
 - Messages
 - Peering Request
 - Peering Response

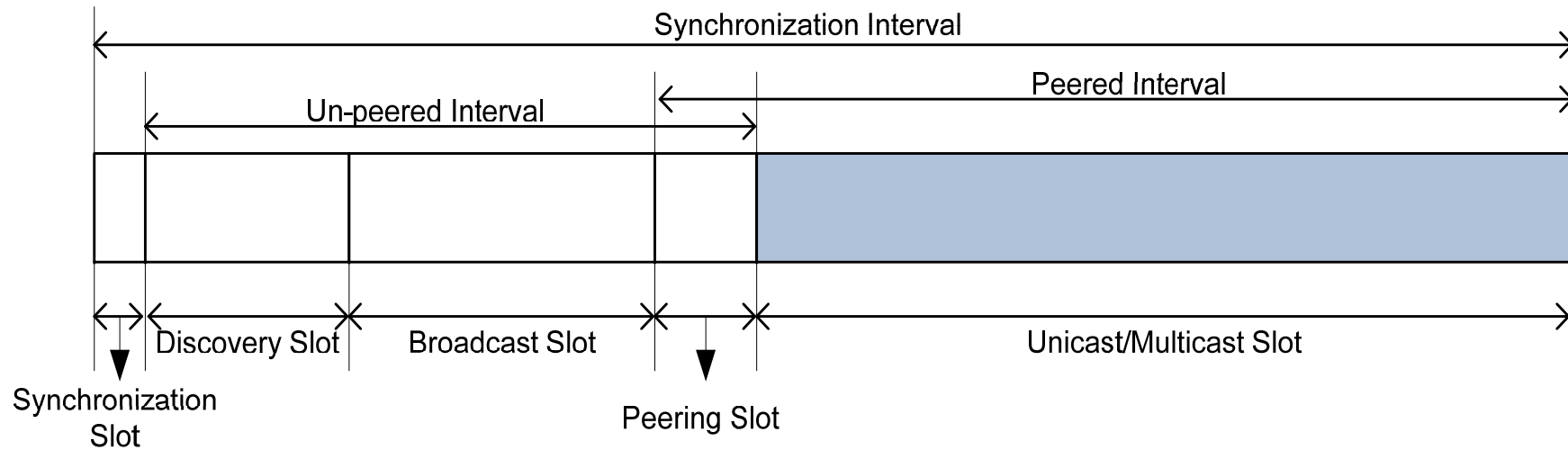
5. Peering

- Design Consideration for Peering Slot
 - Small radio resource comparing to Discovery Slot
 - Peering happens sparsely
 - Handling of multiple peering response to peering requests
 - Possible channel access scheme
 - Contention-based access



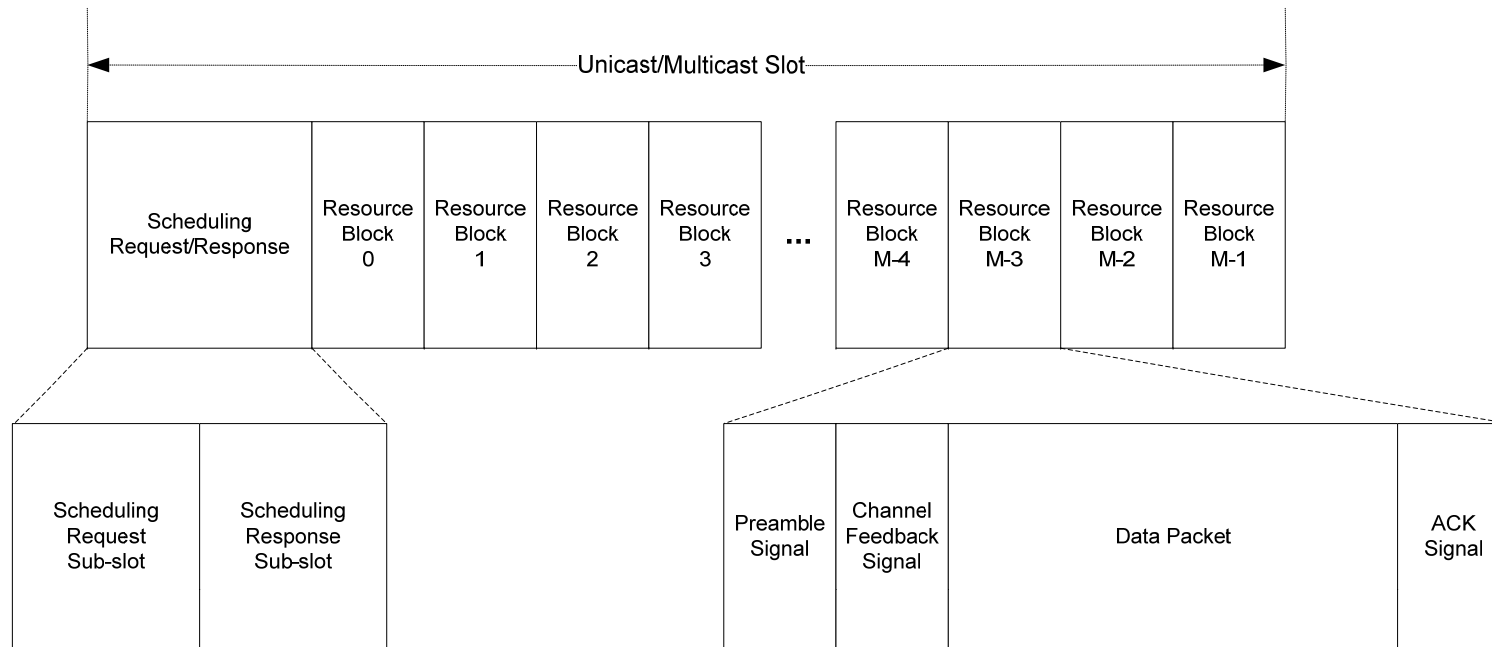
6. Unicast/Multicast

- Features of Unicast/Multicast Slot
 - Only accessed by peered PDs
 - Signaling reduction using Link ID
 - No necessity of sending multiple MAC addresses of both TX PD and RX PD(s)



6. Unicast/Multicast

- Unicast/Multicast Slot comprises
 - Scheduling Request/Response Sub-slots
 - Resource Blocks (RBs)



6. Unicast/Multicast

■ Design Approach

- Contention-free channel access
 - Low signaling overhead & high spatial reuse
- Distributed scheduling
 - Scheduling Request and Scheduling Response
 - These signaling messages contain resource information
 - Related to RB assignment
 - Broadcasted to nearby PDs

Scheduling Request :

| | | |
|---------|----------------------------|-----------------------|
| Link ID | Resource Block Start Index | Resource Block Demand |
|---------|----------------------------|-----------------------|

Scheduling Response :

| | | |
|---------|-------------------------------|--------------------------------|
| Link ID | Resource Block Adjusted Index | Resource Block Adjusted Demand |
|---------|-------------------------------|--------------------------------|

6. Unicast/Multicast

- Required Features of Distributed Scheduling
 - Resource conflict avoidance [referred from 15-13-0376-01]
 - Throughput can be increased by high spatial reuse
 - Link assignment based on the SIR of receiving PD's
 - Utilizing channel-state information
 - A RX PD gives feedback to peered TX PD
 - TX PD controls the size of resources based on feedback
 - Congestion-aware resource assignment
 - Resource is assigned according to results of congestion monitoring

7. Summary

- Key Design Considerations
 - Frame structure
 - Distributed synchronization
 - Peer discovery
 - Protocol stack
 - Discovery Block selection
 - Broadcast
 - Contention-based access within un-peered PDs
 - Peering
 - Contention-based access to establish links
 - Unicast/Multicast
 - Contention-free access within peered PDs
 - Distributed scheduling by request and response

8. References

- [1] XMPP (Extensible Messaging and Presence Protocol) <http://xmpp.org/>
- [2] UUID (Universally Unique Identifier)
http://en.wikipedia.org/wiki/Universally_unique_identifier