

Project: IEEE P802.15 Working Group for Wireless Personal Area Networks (WPANs)

Submission Title: [Samsung Proposal: PAC Operations and Frame Structures]

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Abstract: [Presentation of PAC synchronous operations and frame structure with identified features for working consensus to work on PFD]

Purpose: [Corresponding to Call for Proposal]

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Samsung Proposal: PAC Operations and Frame Structures

July, 2013

Samsung

System Design Approach

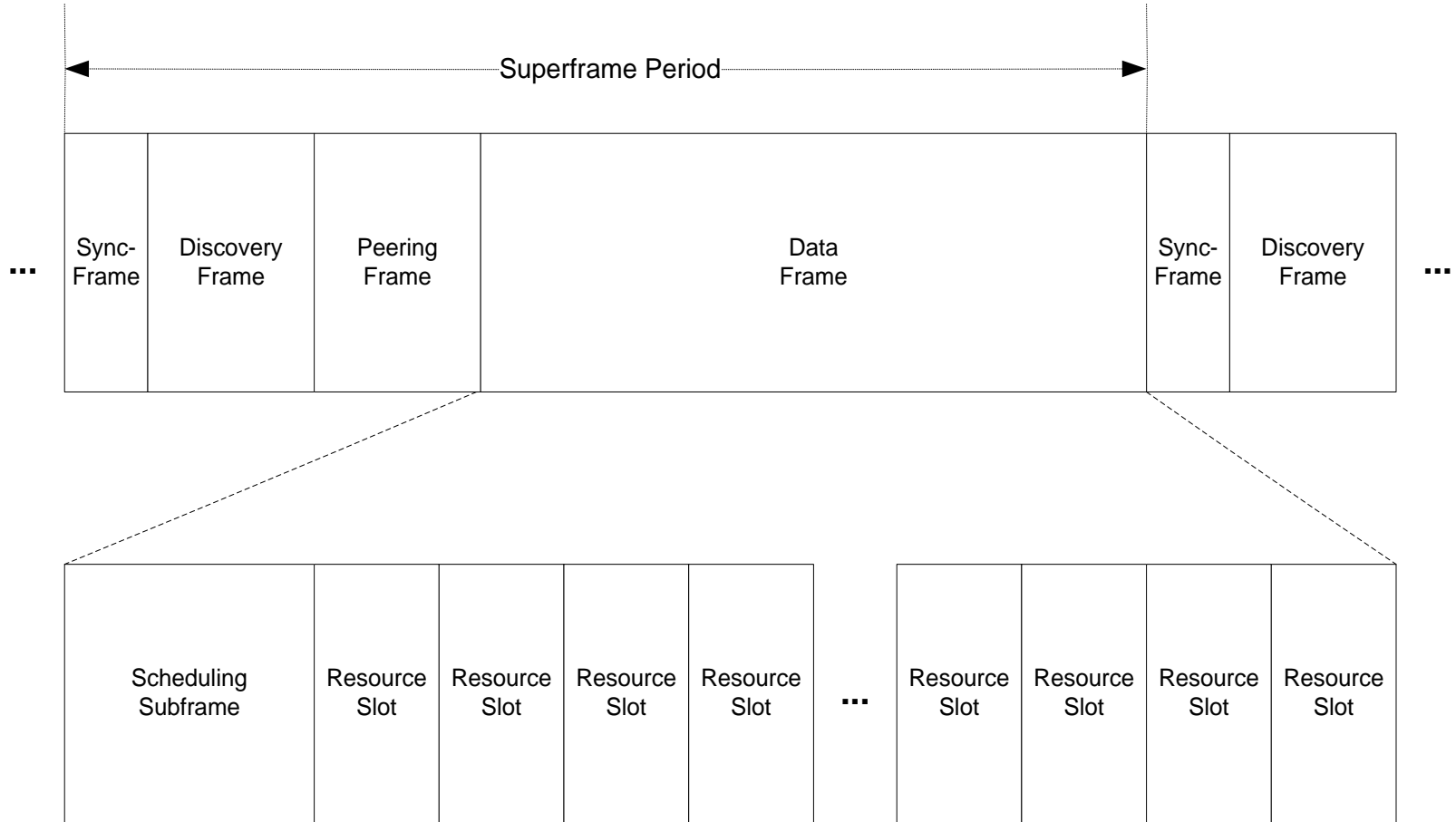
- Synchronous Operation
 - Low power consumption for peer discovery
 - Higher throughput for communication

- Frame Structure
 - Synchronization
 - Peer discovery
 - Peering
 - Data transmission
 - Scheduling for efficient slot allocation
 - Designed for unicast transmission

Fully Distributed Operations

- Synchronization
 - Based on Pulse-Coupled Oscillator (PCO) algorithm
- Peer Discovery
 - Prior to peering
 - Broadcast Peer Discovery Message via selected resource
- Peering
 - Link establishment
- Data transmission
 - Scheduling and interference management
 - Resource request and response

PAC Frame Structure



Synchronization

- Synchronization for Scalable Network
 - Distributed synchronization
 - Master-slave synchronization should be avoided
 - PDs between two different synchronized group happen
 - It is matched well to flat architecture (no hierarchy)
 - Synchronization should be done before peer discovery
 - Peer discovery prior to link connection (peering)
 - Broadcast-based synchronization mechanism is required

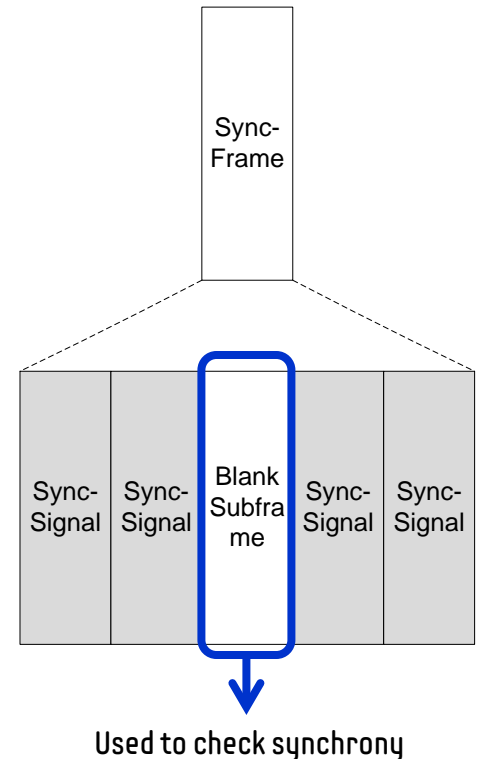
- Proposed Design Approach
 - Physical layer signaling based synchronization

Synchronization

- Initial Synchronization
 - a. Start in initial synchronization mode
 - b. PD monitors sync-signals during sync-period
 - c. If at least one sync-signal is detected during sync-period, perform according to distributed synchronization mechanism
 - d. Else, start PAC operations based on frame structure in maintaining synchronization mode

Synchronization

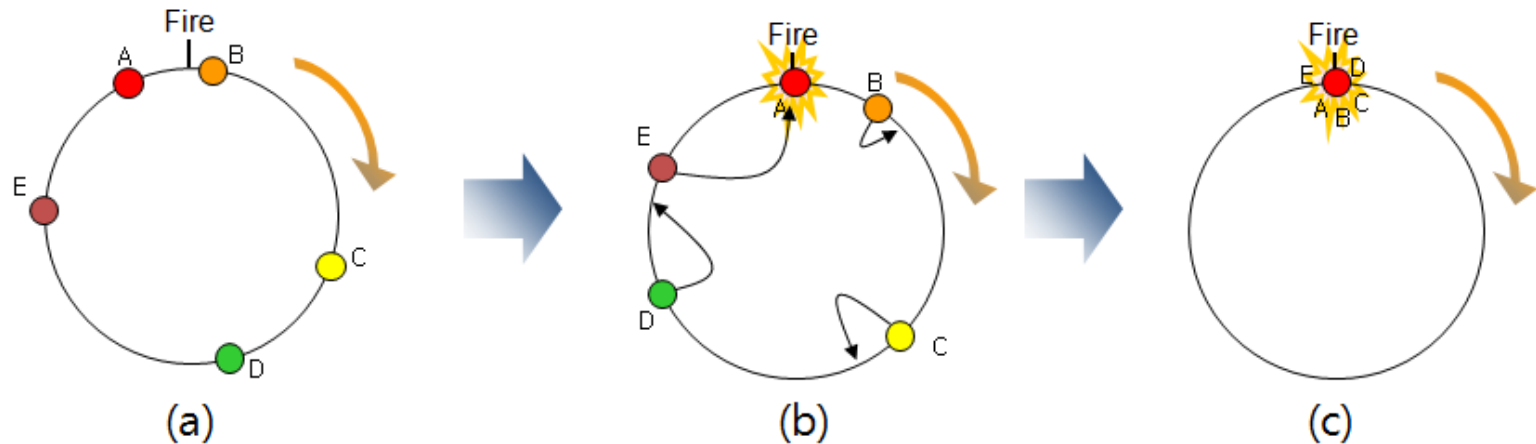
- Maintaining Synchronization
 - a. PD sends sync-signal periodically, but checks synchrony state sometimes without sending sync-signal
 - b. If in-synchrony, PD adjusts oscillator for phase drift compensation
 - c. If out-of-synchrony, go to initial synchronization mode



Synchronization

■ PCO Synchronization [1]

- In initial synchronization mode
- Pulse-based approach
 - Oscillator coupled by pulse exchange via physical layer



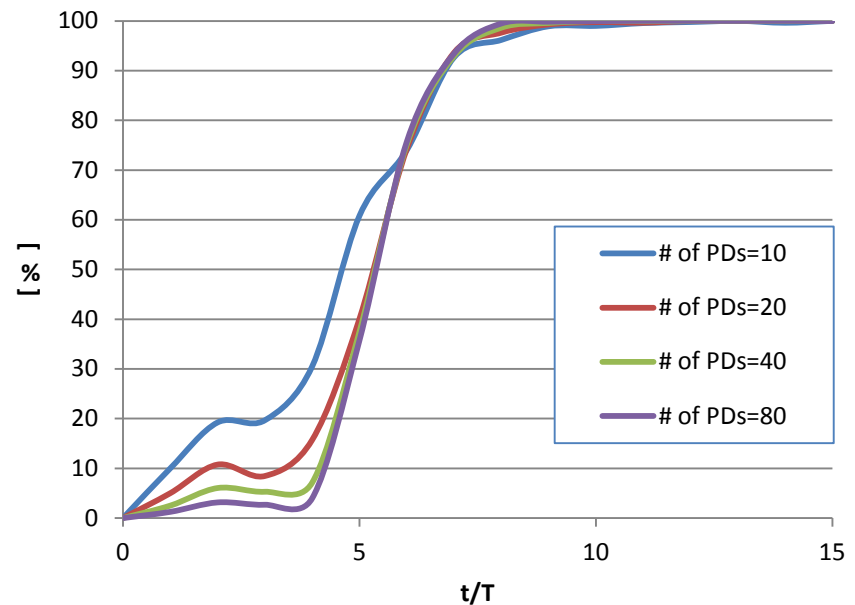
- (a) All nodes have oscillator with the same phase increment rate
- (b) One node fires, then other nodes adjust oscillator according to the predefined function without state other than its internal phase
- (c) Finally, all nodes converge to the same time base

Synchronization Performance

■ Simulation Condition

- Dimension
 - 500 x 500 m²
- Coupling factor
 - 0.05
- Dissipation factor
 - 10
- Sync-signal period
 - T = 10 msec

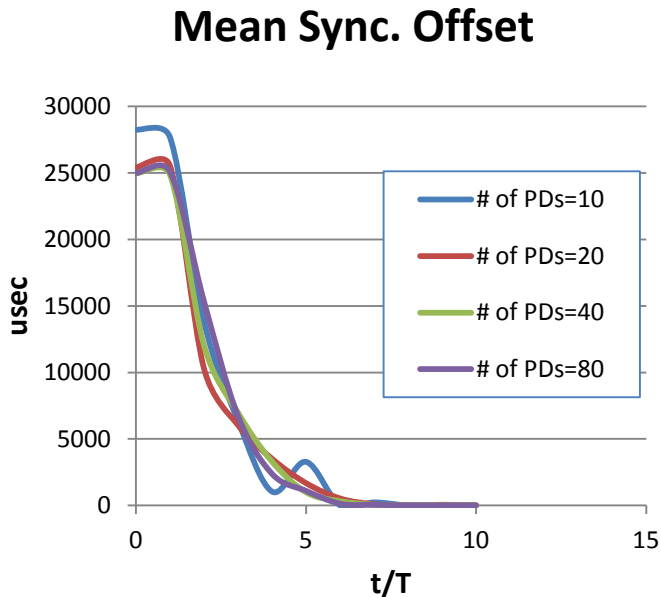
The Ratio of Synchronized PDs



* Simulation methodology refers from [2]

Synchronization Performance

■ Synchronization Accuracy



	$N_{PD}=10$	$N_{PD}=20$	$N_{PD}=40$	$N_{PD}=80$
0	28260.18	25418.36	24975.59	25000
1	27694.1	25552.14	24881.05	25101.9
2	13833.46	10297.46	12051.57	15352.84
3	6376.196	6033.111	6973.267	6729.542
4	1036.911	3505.74	3294.19	2374.513
5	3273.107	1669.398	985.3763	1094.029
6	0.248531	517.5309	312.7014	126.0405
7	223.569	78.55826	40.29059	36.55889
8	0.217967	0.239922	28.45693	14.23396
9	0.2067	39.74689	2.751351	0.258708
10	0.211183	0.253245	0.250434	0.257205

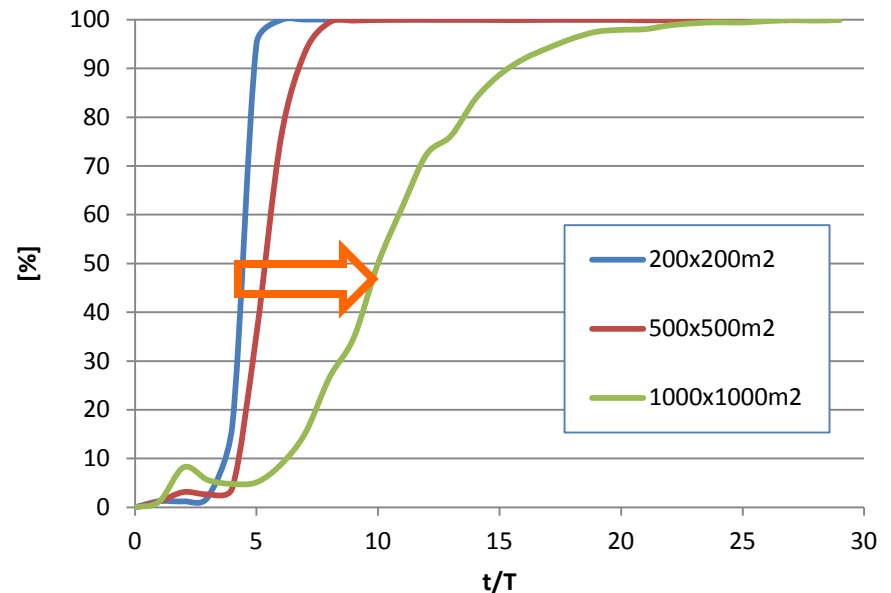
Achieves under 1 usec synchronization accuracy!

Synchronization Performance

■ Simulation Condition

- # of PDs
 - 80
- Coupling factor
 - 0.05
- Dissipation factor
 - 10
- Sync-signal period
 - $T = 10$ msec

The Ratio of Synchronized PDs

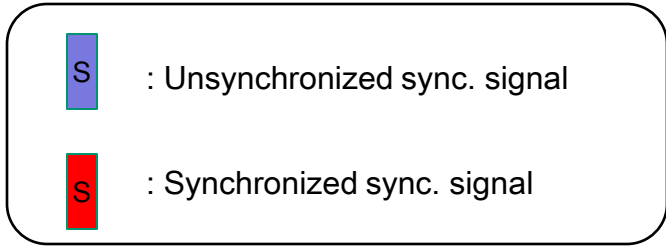
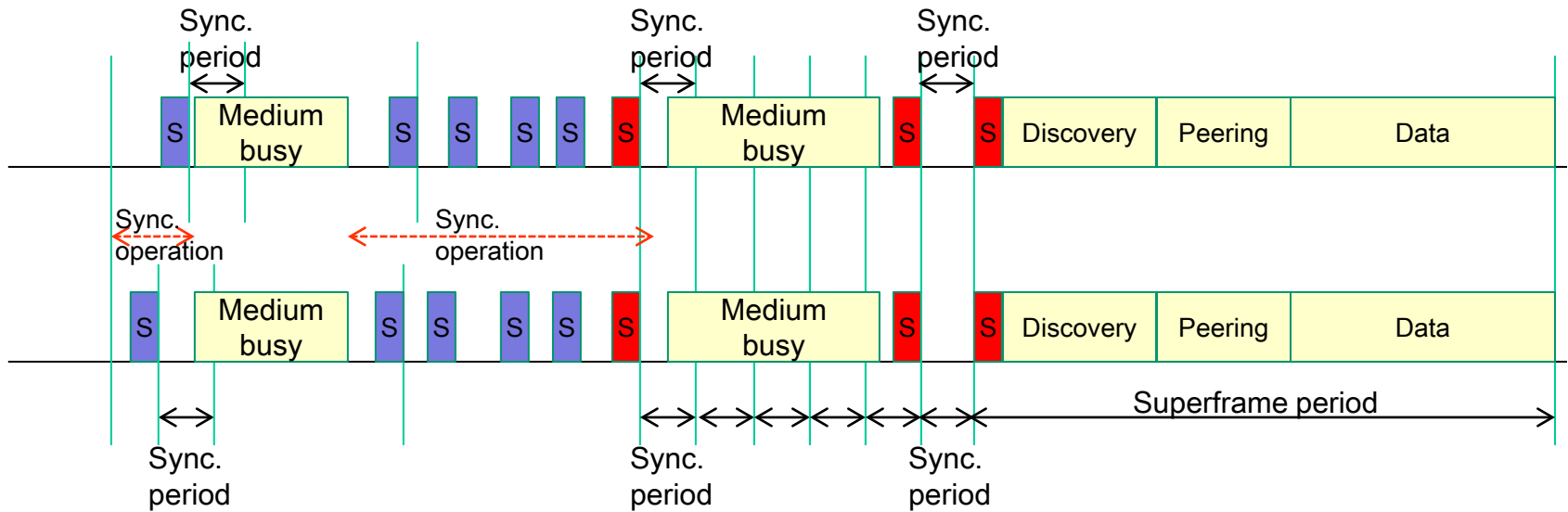


Long latency as being in low density

Synchronization for Unlicensed

- In Initial Synchronization Mode
 - a. PD senses energy level while doing sync. operation
 - b. If medium is busy,
 - Pend synchronization operation
 - else
 - Keep synchronization operation
 - c. Check synchrony state
 - Two conditions
 - No sync-signals during sync. period
 - Medium is not busy during sync. period
 - Superframe starts in maintaining synchronization mode

Synchronization for Unlicensed



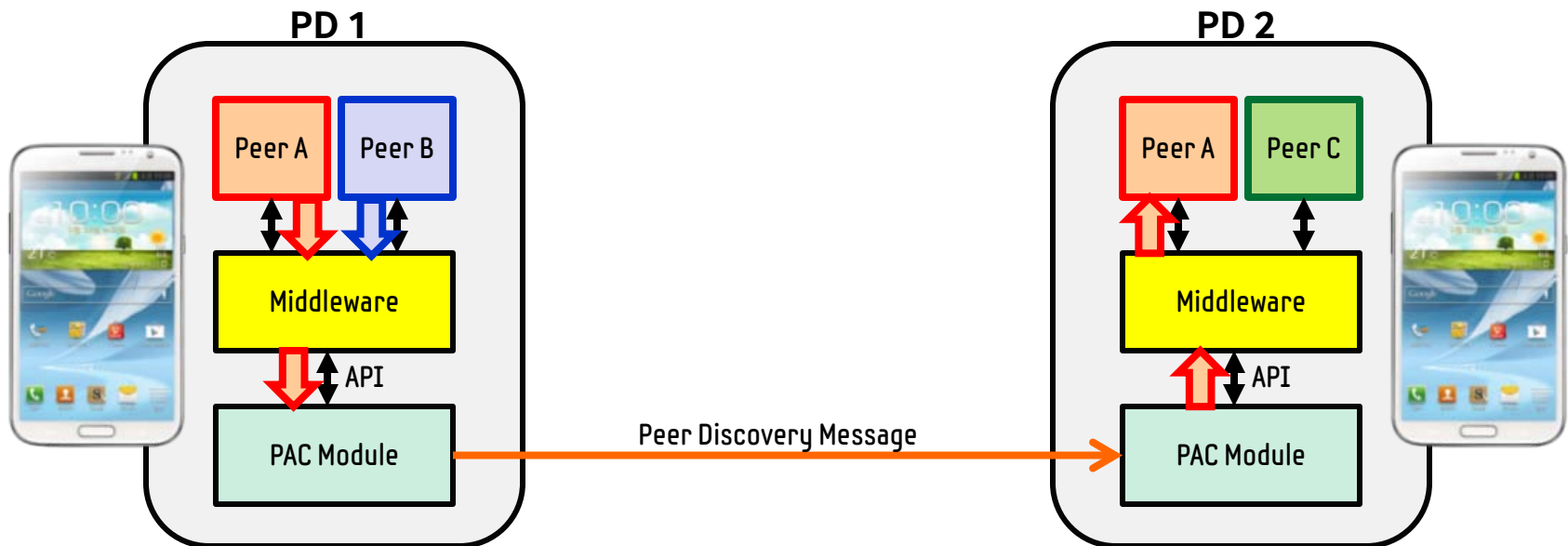
Peer Discovery

■ Design Considerations

- The meaning of Peer Discovery
 - It depends on what is discovered
 - Discovery of application → PACbook
 - Discovery of friend → bob@PACbook
 - Discovery of group → IEEE802@PACbook
- Who gives peer discovery information (PDI)
 - Application or middleware
 - They may have access to internet or may not
 - Authentication
 - Only PDs with same PDI can be discovered each other
 - PDI may be pre-installed or given from network
- Unified Peer Discovery mechanism required

Peer Discovery

- What is peer discovery?
 - A peer is an application, not a device
 - Application-centric discovery



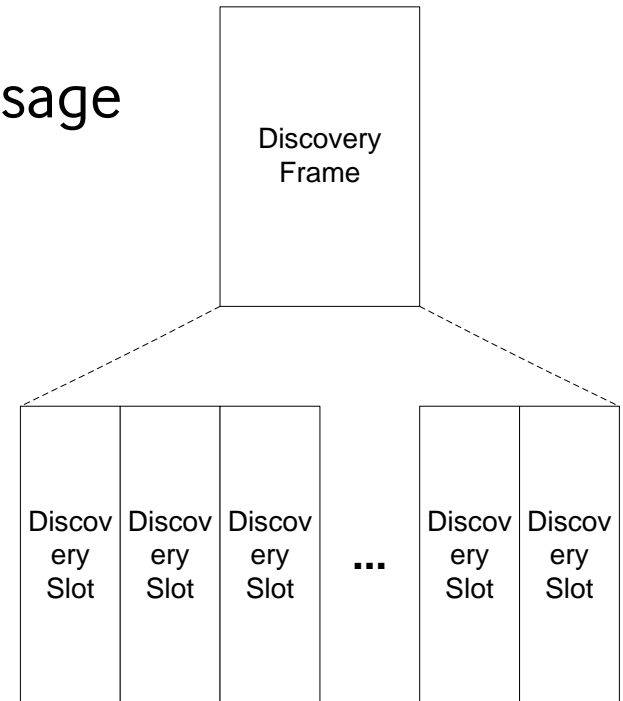
Peer Discovery

- Peer Discovery Message (PDM)
 - It is generated based on PDI including
 - Application-specific ID,
 - Application-specific user ID,
 - Application-specific group ID,
 - Or any upper layer discovery information
 - Middleware in PD1 indicates PDM to PAC module
 - At PAC module of PD1, it is broadcasted to PD2
 - At PAC module of PD2, it is delivered to middleware
 - Middleware in PD2 is aware of nearby peer
 - By matching result with pre-stored PDM

Peer Discovery

■ Discovery Frame

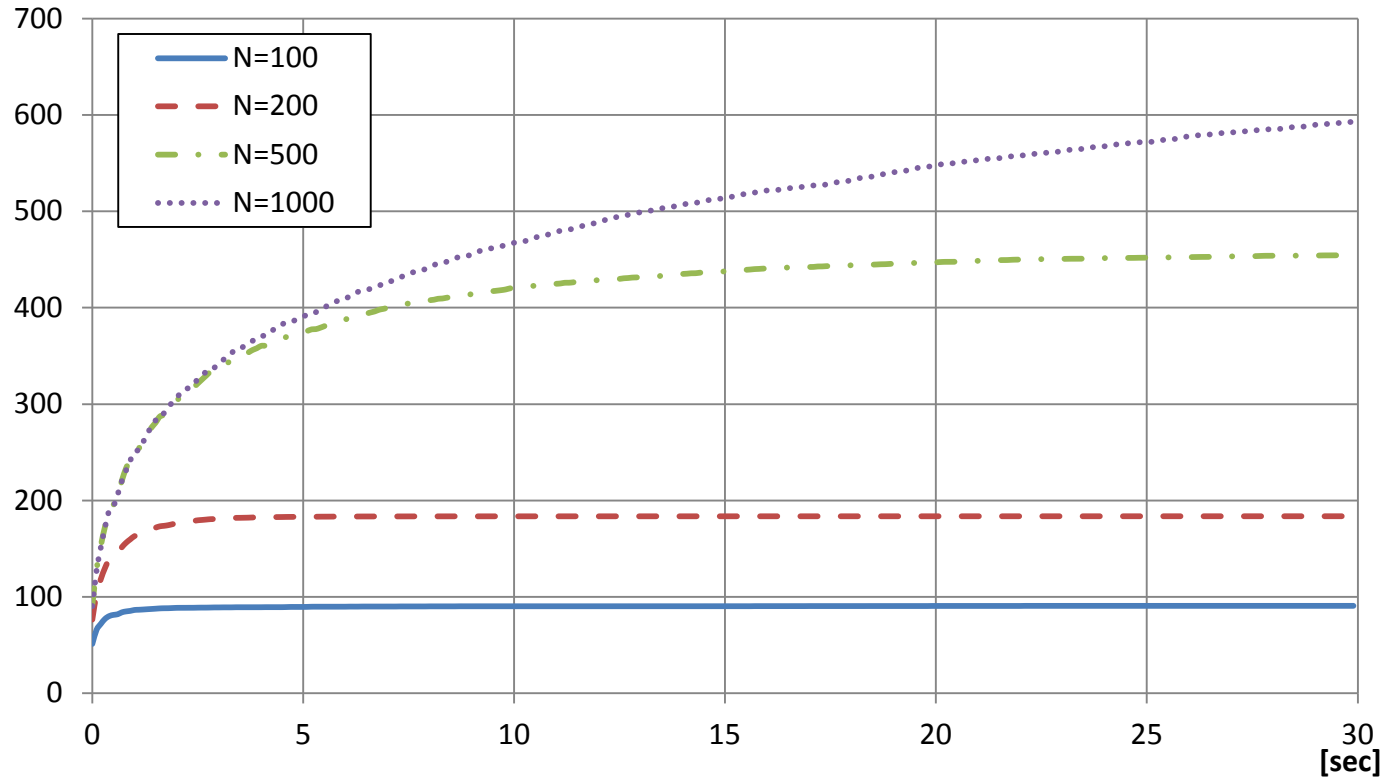
- It is comprised of multiple Discovery Slots
- PD selects one Discovery Slot
- PD broadcasts Peer Discovery Message
 - At selected Discovery Slot
- PD monitors congestion
 - If congested, PD selects different Discovery Slot for next transmission
 - If not congested, PD keeps the current Discovery Slot



Peer Discovery

■ Performance Results

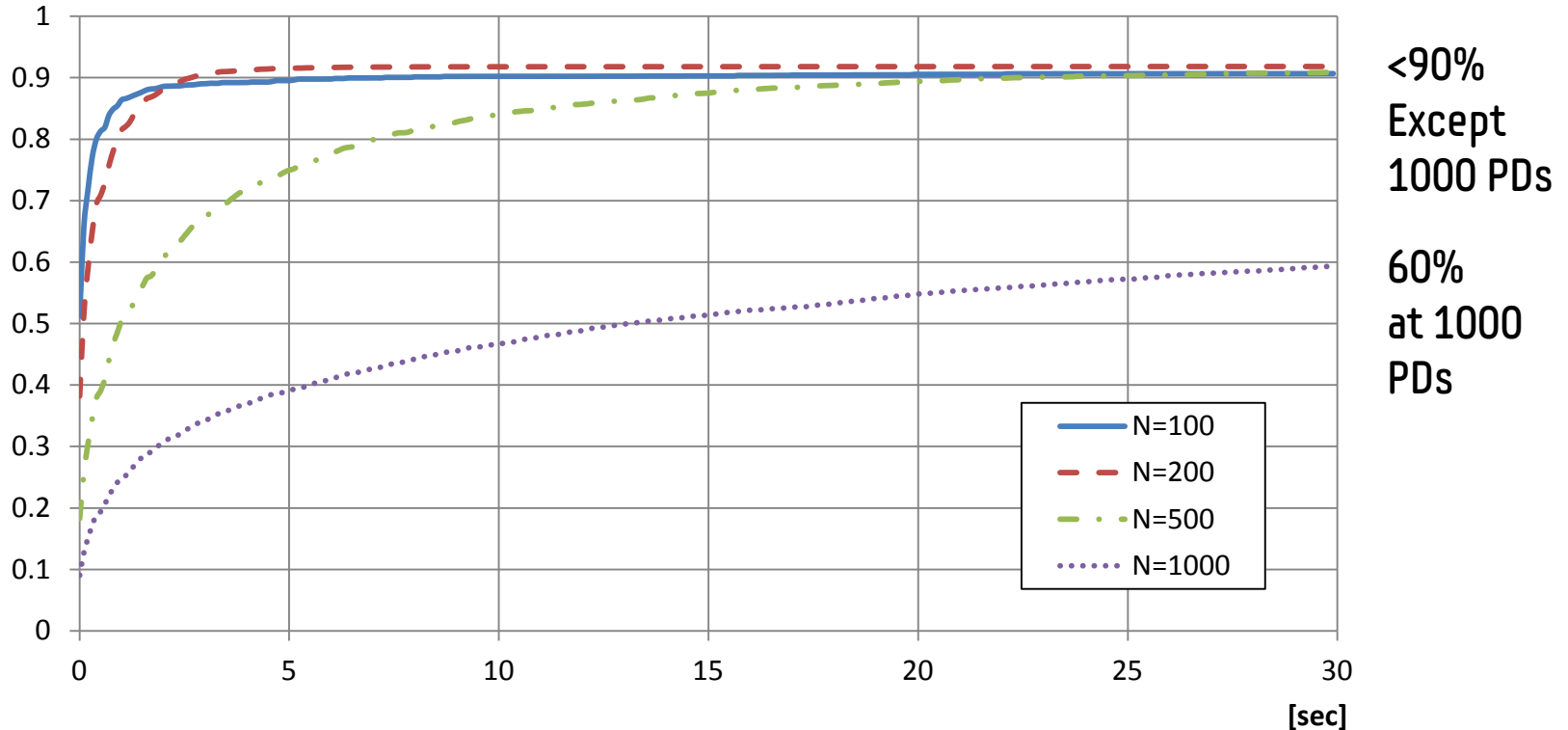
– The number of discovered PDs



Peer Discovery

■ Performance Results

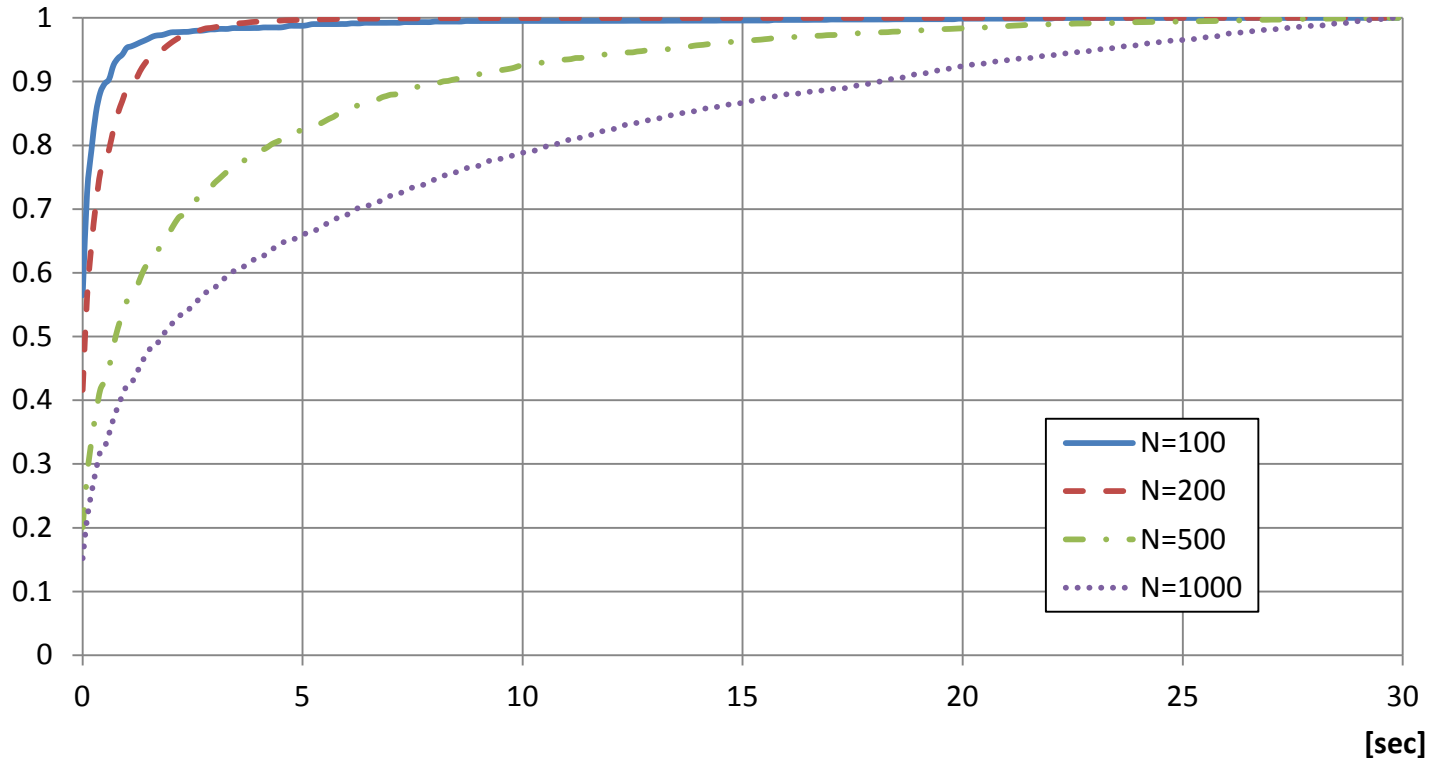
– The ratio of discovered PDs (normalized by total PDs)



Peer Discovery

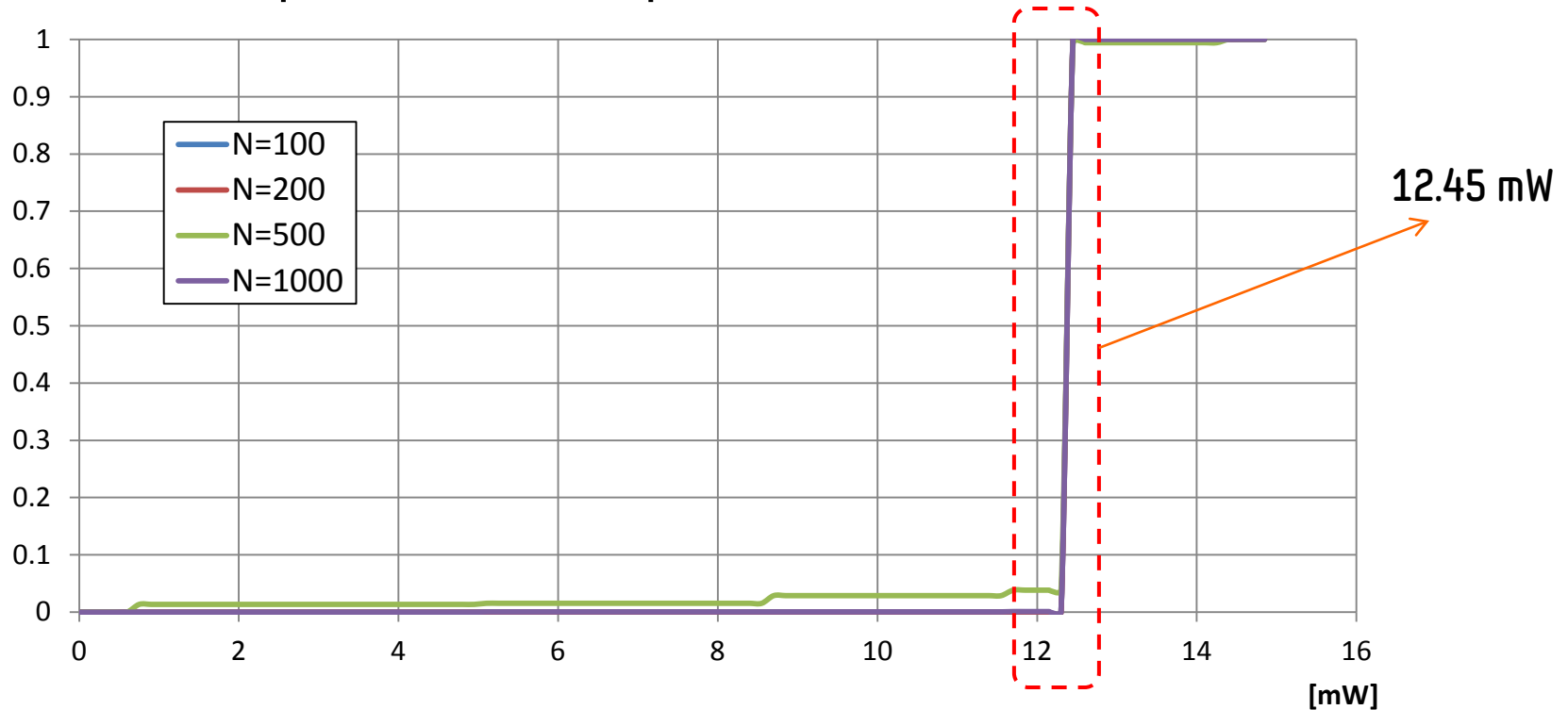
■ Performance Results

– CDF of the discovery latency



Peer Discovery

- Performance Results
 - CDF of power consumption



Peering

■ The role of Peering

- A procedure to connect to discovered peer
 - Triggered by application automatically or by user manually
- Link establishment
 - Between the discovering peer and the discovered peer
 - Exchange of information for setup
 - device capability, or etc
 - Determine link related parameters
 - Link ID, QoS class, link range, or etc
- Messages
 - Peering Request
 - Peering Response

Data Transmission

■ Data Frame

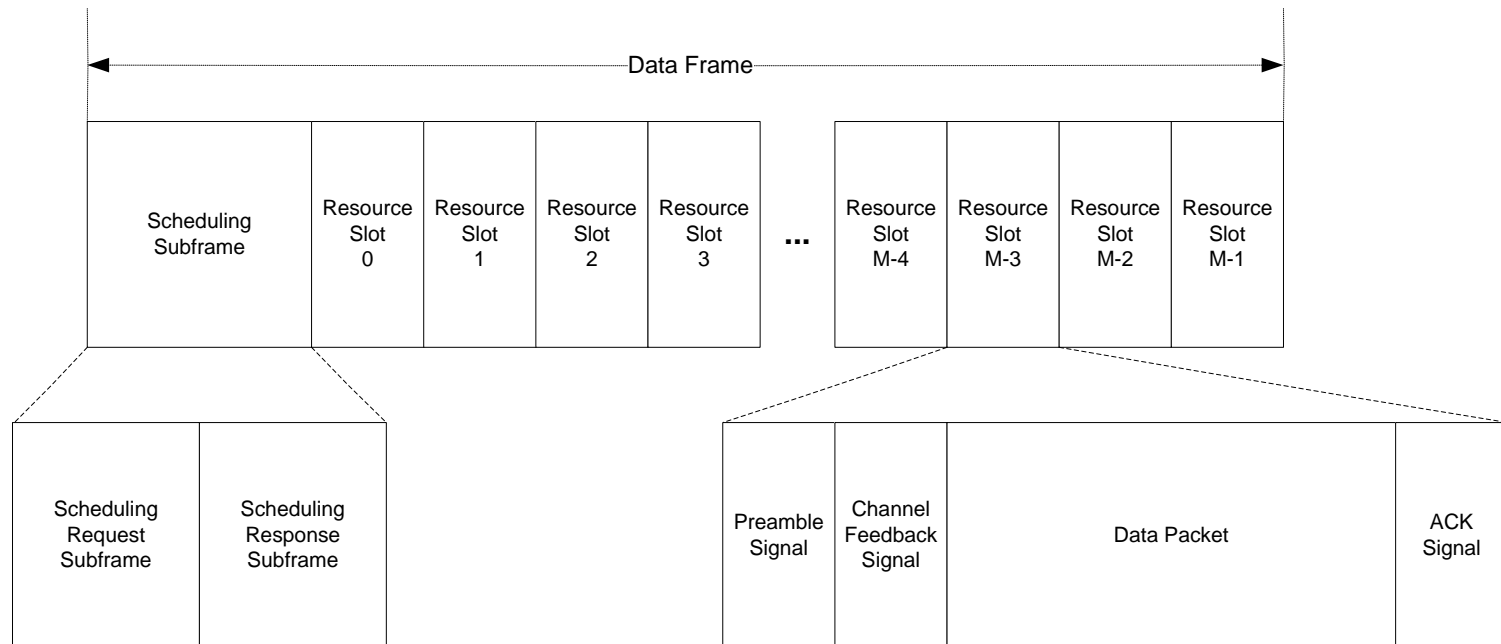
- Only accessed by peered PDs
- Signaling reduction using Link ID set up during peering
 - No necessity of keeping two IDs for both Tx PD and Rx PD

■ Design Considerations

- Connection is the result from peering
 - Unicast/multicast including single/multi-hop
 - Network protocol shall be operated only over connected links
 - E.g. routing, grouping, etc

Data Transmission

- Data Frame
 - Scheduling subframe
 - M Resource Slots



Data Transmission

■ Design Approach

- Synchronized slotted channel access
- Distributed scheduling to avoid slot confliction
 - Scheduling Request and Scheduling Response
 - These signaling messages contain resource information
 - Related to resource slot assignment
 - Broadcasted to nearby PDs

Scheduling Request :

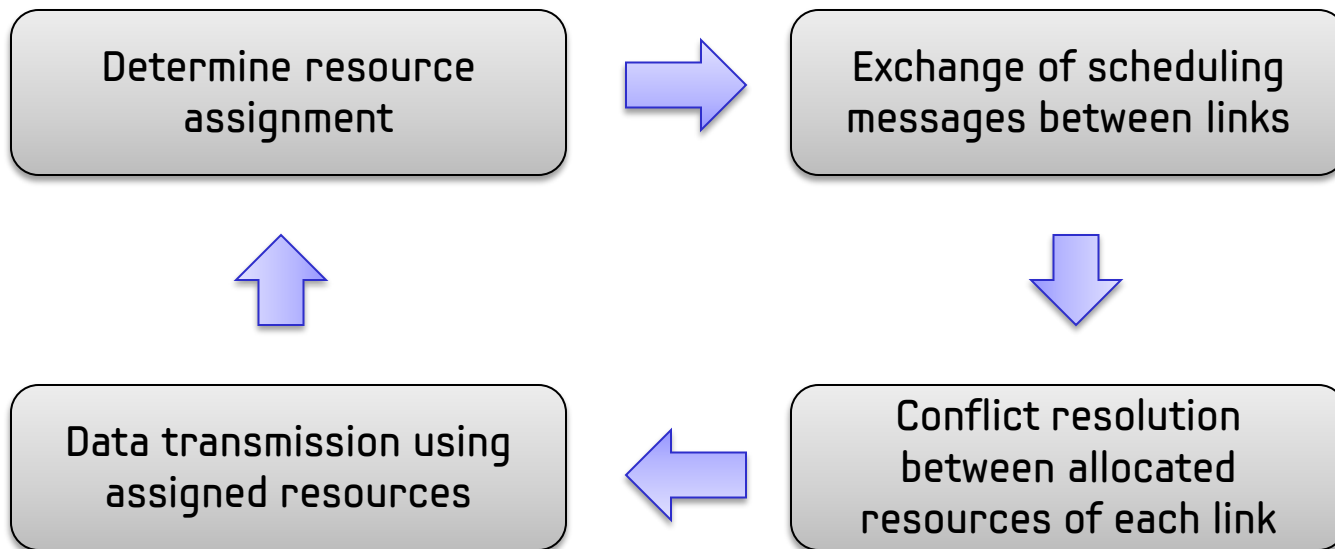
Link ID	Resource Slot Start Index	Resource Slot Length
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Scheduling Response :

Link ID	Resource Slot Adjusted Index	Resource Slot Adjusted Length
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Data Transmission

- Concept of Distributed Scheduling Procedure

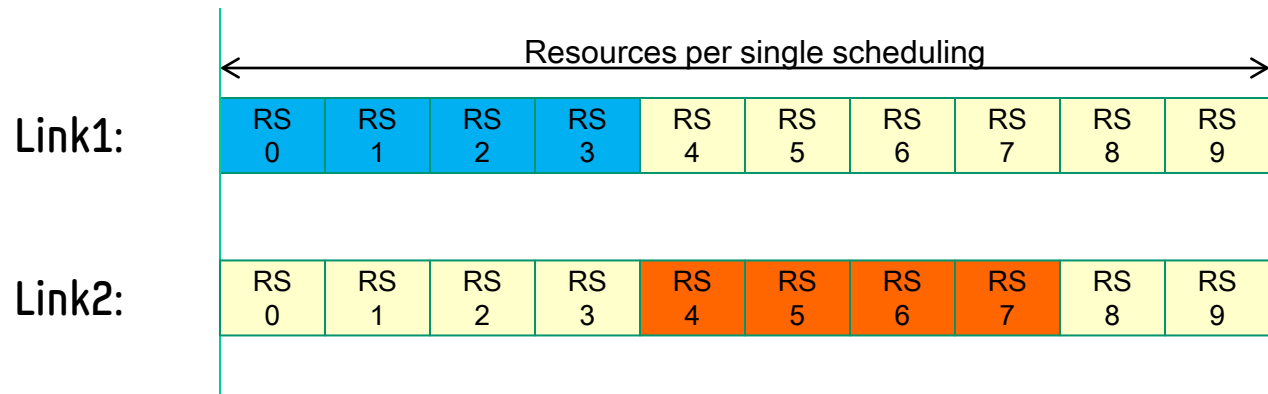
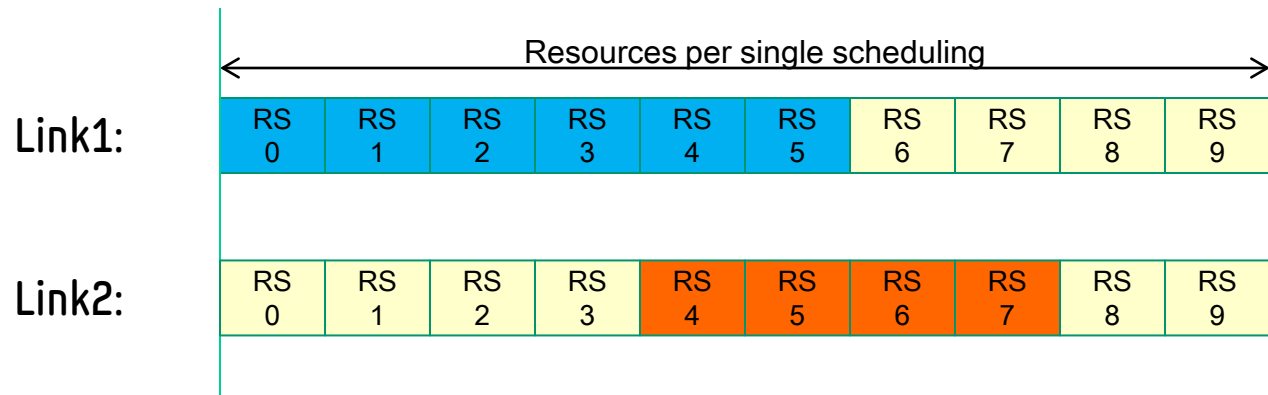


Data Transmission

Determine resource assignment

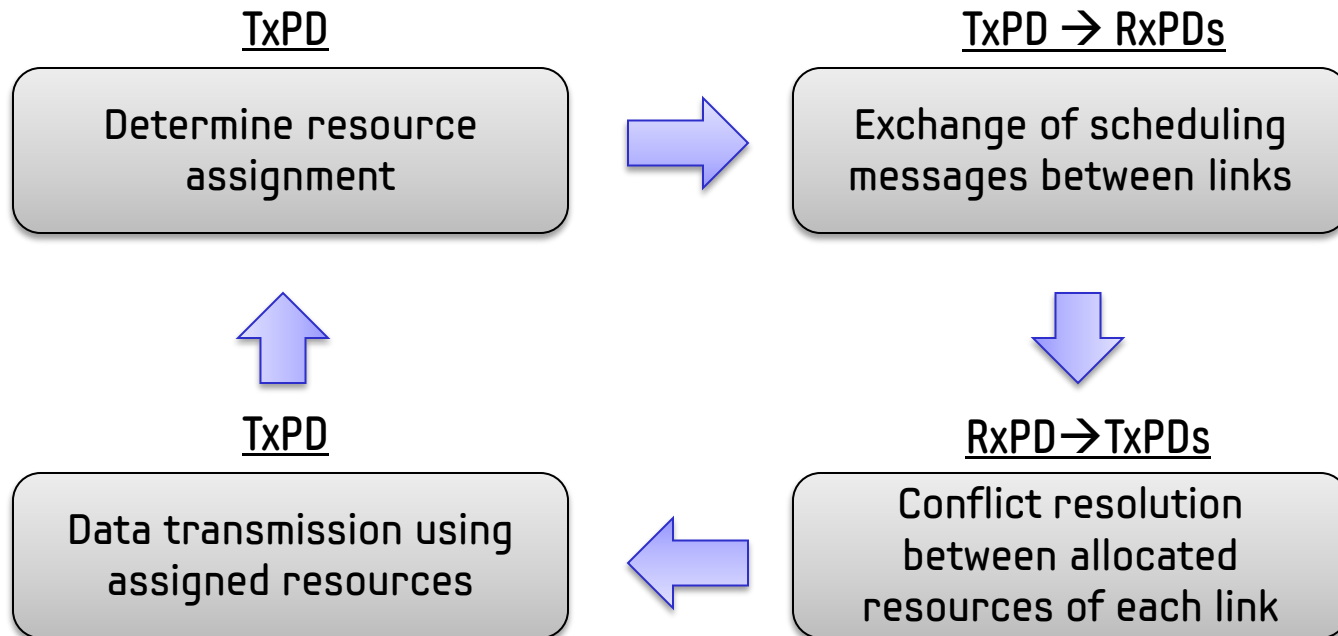


Conflict resolution between allocated resources of each link



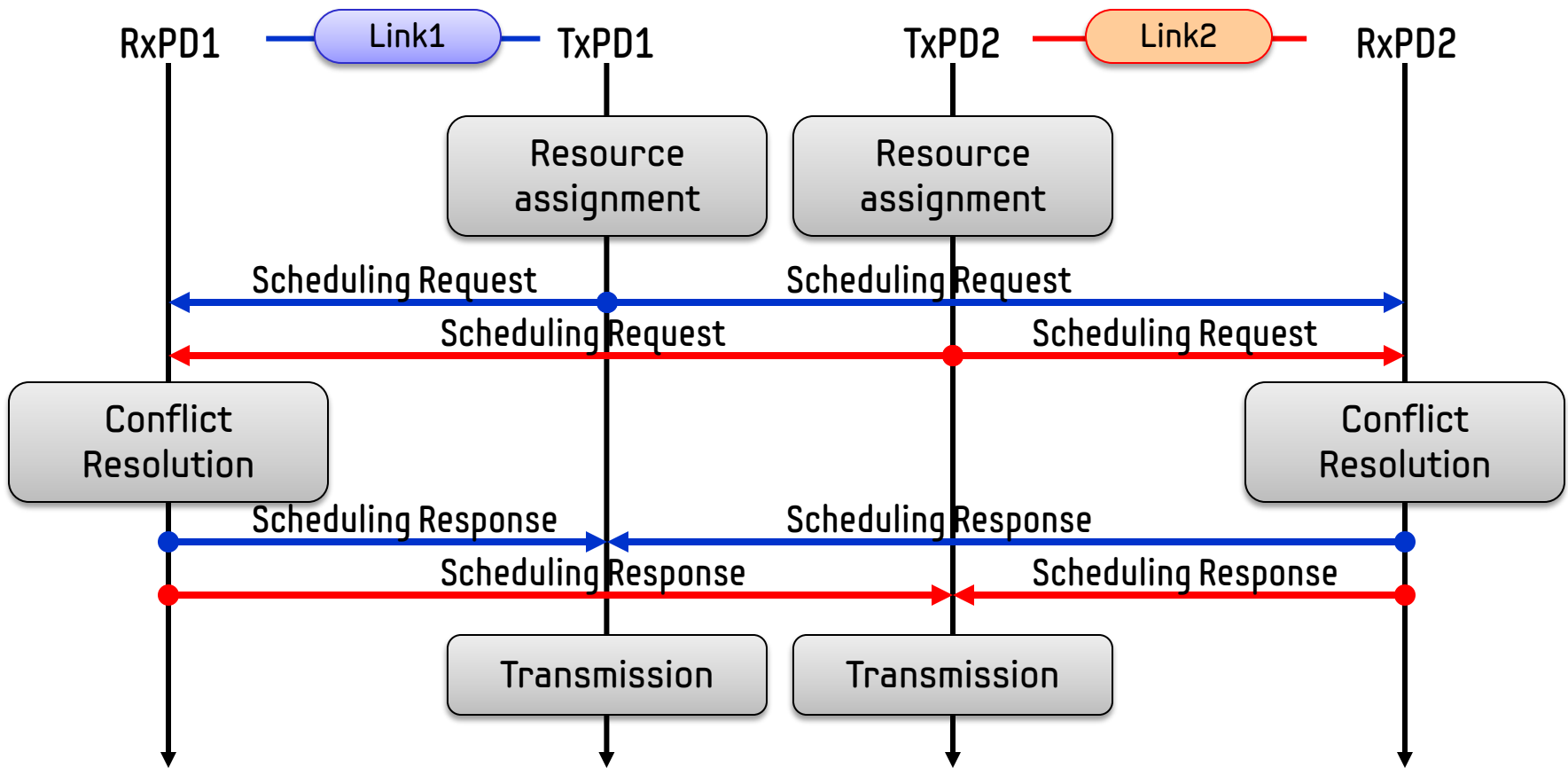
Data Transmission

- Concept of Distributed Scheduling Procedure



Data Transmission

■ Operation Flowchart



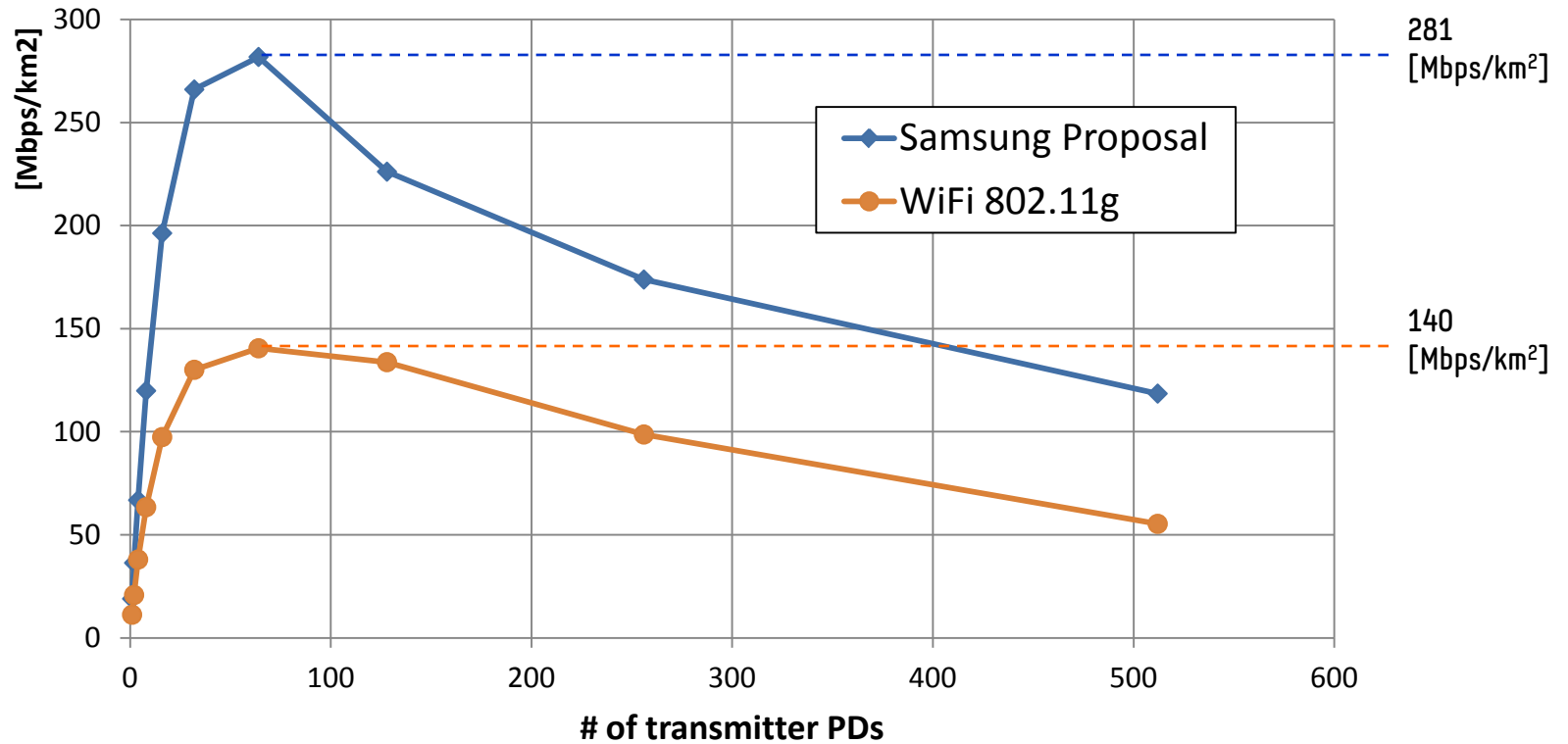
Data Transmission

- How to achieve higher throughput?
 - TDMA
 - CSMA/CA : protocol overhead due to asynchronous access
 - Low signaling overhead
 - Each scheduling message needs 16 bits ($2^{16}=2^6 \times 2^5 \times 2^5$)
 - E.g. 64 Link IDs, 32 slot start index, 32 slot length
 - 3.125% as assuming 256 FFT size
 - Spatial resource reuse
 - According to the threshold value used for conflict resolution
 - 9 dB is used for simulation
 - Maximize the number of concurrent transmission links

Data Transmission

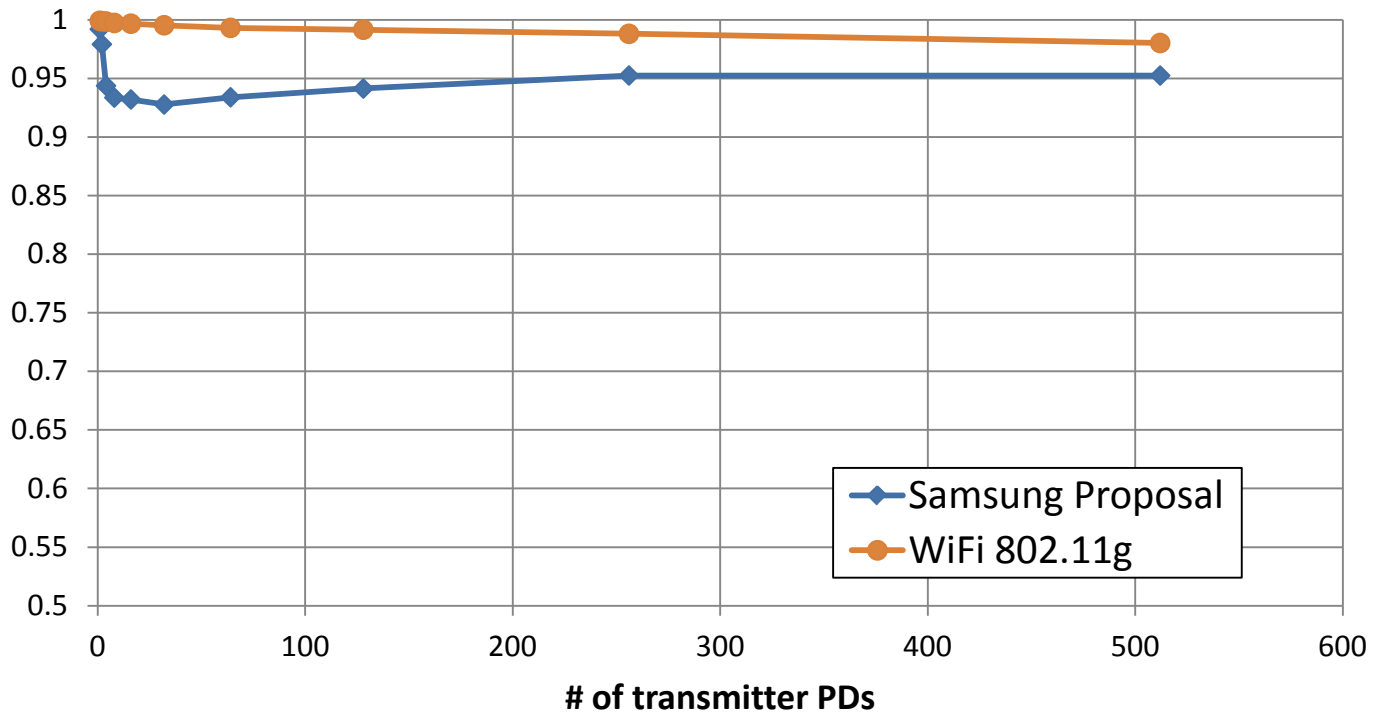
- Performance Results (full buffer)

- Areal sum goodput [Mbps/km²]



Data Transmission

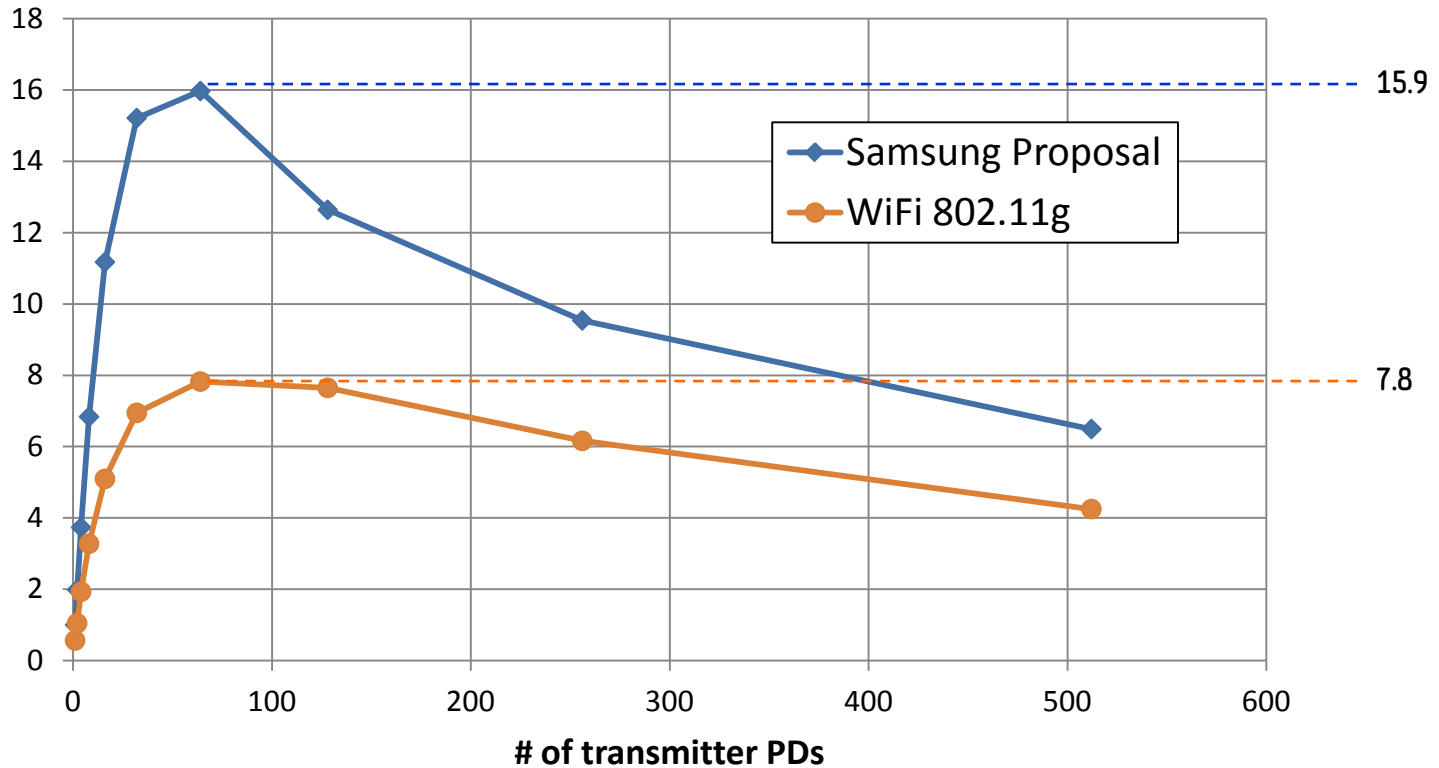
- Performance Results (full buffer)
 - Data packet reception efficiency [ratio]



Data Transmission

- Performance Results (full buffer)

- The number of concurrent transmission links

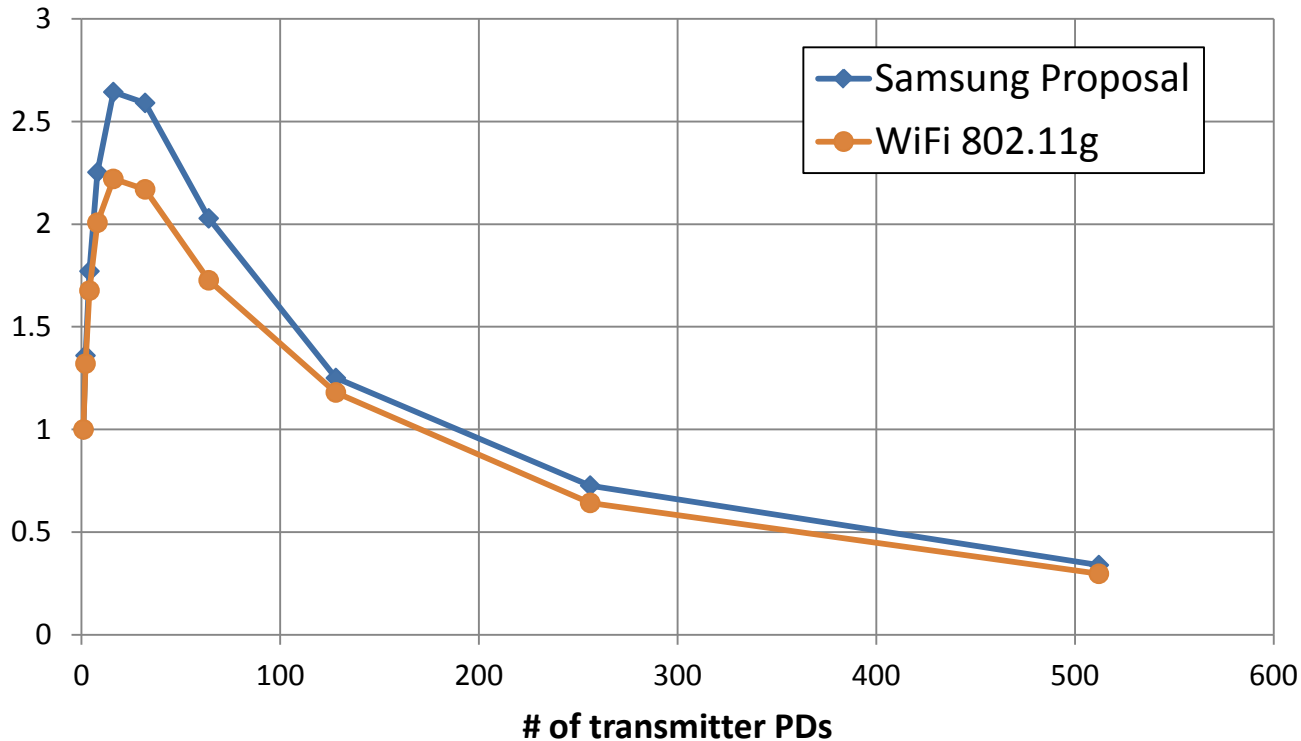


Data Transmission

■ Performance Results (full buffer)

– Jain’s fairness index (modified)

$$J(x_1, x_2, \dots, x_n) = \frac{\left(\sum_{i=1}^n x_i\right)^2}{\sqrt{n} \cdot \sum_{i=1}^n x_i^2}$$



Summary

- Key Design Considerations
 - Synchronization prior to peer discovery
 - PCO synchronization
 - Peer discovery
 - Application-centric discovery
 - Congestion-aware slotted discovery resource selection
 - Peering
 - Triggered by application or users
 - Preparation for channel access frame
 - Channel access
 - Slotted scheduling by request and response messages

Identified Working Consensus

- Channel-access
 - IEEE802.15.8 PAC shall consist of several types of frames to serve different operations, based on contention-free channel-access scheme
- Discovery
 - Only PDs with same PDI shall be discovered each other
- Peering
 - Peering Request message and Peering Response message shall be exchanged to establish a link.
 - Network protocol such as routing shall be operated only over connected links.

Identified Working Consensus

- Synchronization
 - A PD shall be in synchrony state prior to peer discovery procedure and peering procedure
 - IEEE802.15.8 PAC has two synchronization mode including Initial Synchronization mode and Maintaining Synchronization mode

Identified Working Consensus

■ Scheduling

- A PD shall determine Resource Slot based on the predetermined distributed scheduling algorithm.
- A PD shall determine one or multiple Resource Slot during Scheduling subframe as contention-free channel access scheme.

References

- [1] *Distributed synchronization in wireless networks*, Simeone, Osvaldo ; Spagnolini, Umberto ; Bar-Ness, Yesheskel ; Strogatz, Steven H., Signal Processing Magazine, IEEE Volume: 25 , Issue: 5, Digital Object Identifier: 10.1109/MSP.2008.926661, Publication Year: 2008 , Page(s): 81 – 97
- [2] *Emergent Slot Synchronization in Wireless Networks*, Tyrrell, A.; Auer, G.; Bettstetter, C., Mobile Computing, IEEE Transactions on, Volume: 9 , Issue: 5, Page(s): 719 – 732, Digital Object Identifier: 10.1109/TMC.2009.173, Publication Year: 2010 , Page(s): 719 – 732
- [3] *Selective Pulse Coupling Synchronicity for Sensor Network*, Yu Niu ; d'Auriol, B.J. ; Xiaoling Wu ; Jin Wang ; Jinsung Cho ; Sungyoung Lee, Sensor Technologies and Applications, 2008. SENSORCOMM '08. Second International Conference on, Digital Object Identifier: 10.1109/SENSORCOMM.2008.59 Publication Year: 2008 , Page(s): 123 - 128

Appendix

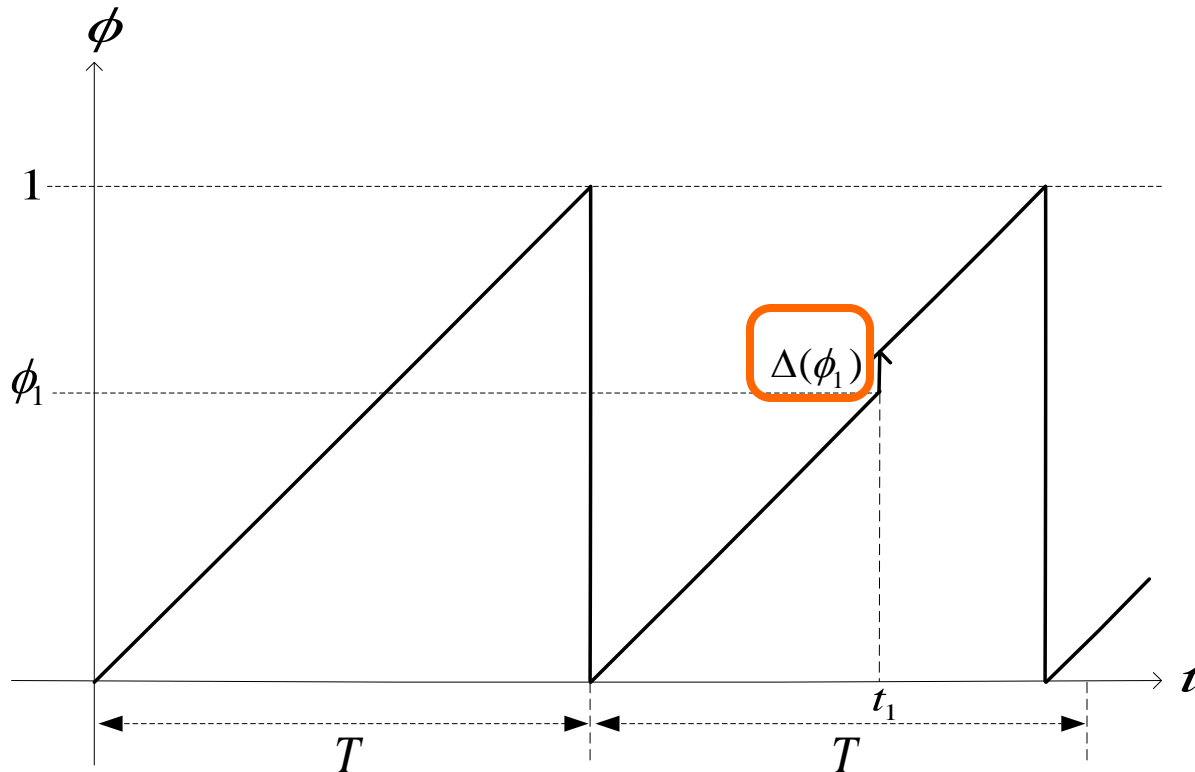
- PCO Synchronization
- Comparison of two ways for Peer Discovery

PCO Synchronization

- Pulse Coupled Oscillator (PCO) Synchronization
 - Fully distributed synchronization algorithm
 - Doesn't need other PDs timing information
 - Each PD has
 - an oscillator (or counter) with the identical increment rate
 - a same function to adjust phase of oscillator
 - Features
 - Simple
 - Scalable
 - no hierarchy (=flat)
 - Fast convergence time

PCO Synchronization

- Phase adjustment using only internal value



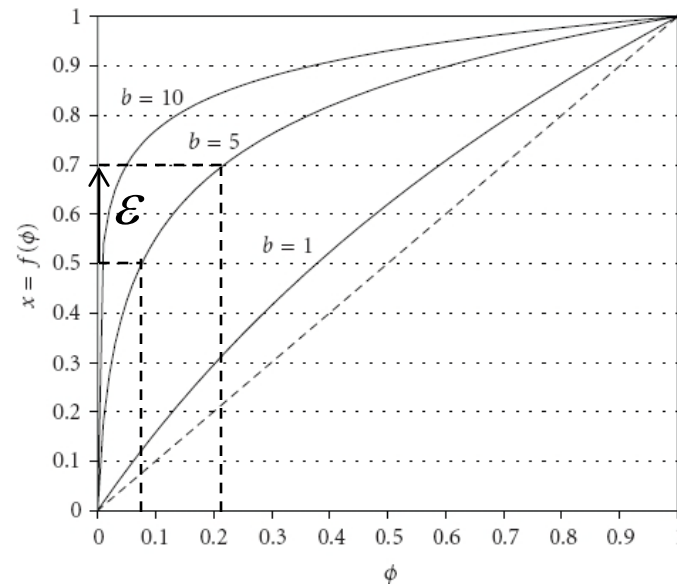
PCO Synchronization

■ Phase adjustment algorithm [2][3]

$$- \phi + \Delta(\phi) = f^{-1}(f(\phi) + \varepsilon)$$

$$- f(\phi) = \frac{1}{b} \cdot (1 + [e^b - 1] \cdot \phi)$$

- ε : coupling factor
- b : dissipation factor

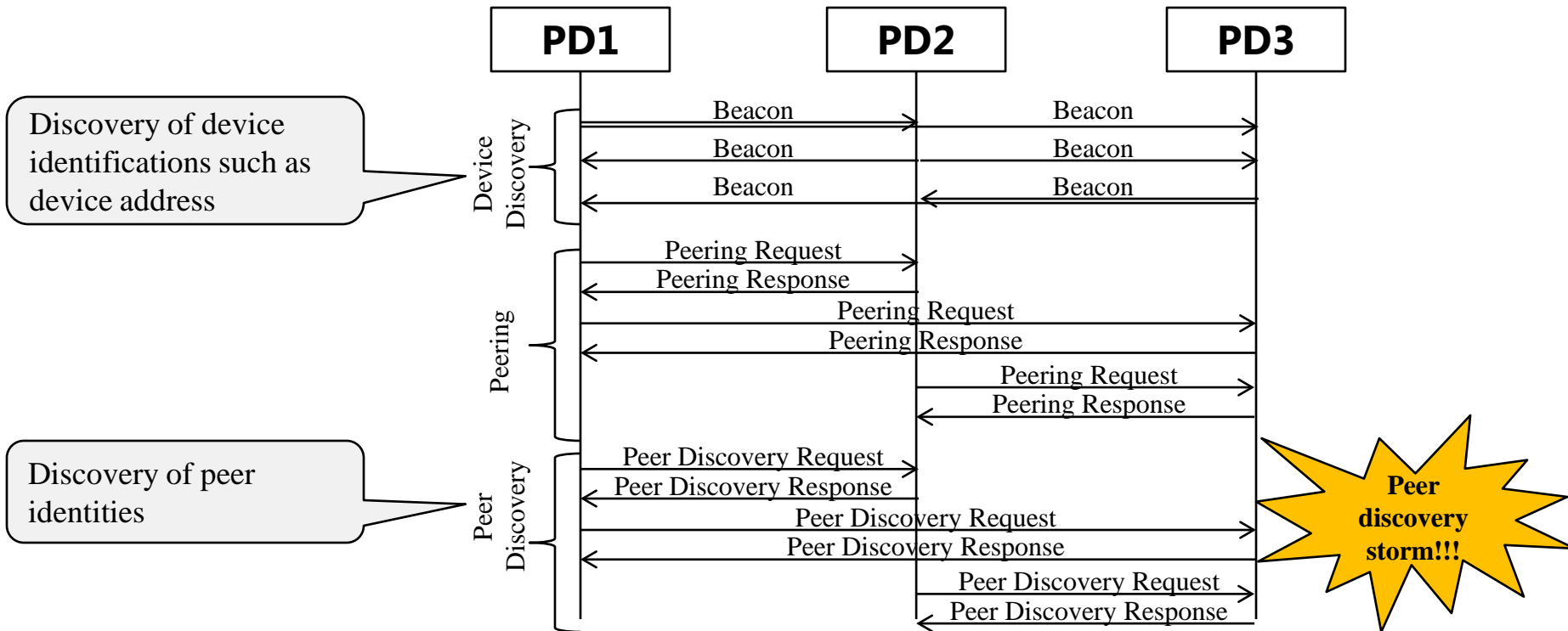


PCO Synchronization

- Phase adjustment algorithm [2][3]
 - selective update for fast convergence
 - If $(2 \times \phi + \Delta(\phi)) > 1$,
$$\hat{\phi} = \min(1, \phi + \Delta(\phi))$$
 - else no update
 - Refractory period to avoid ping-pong effect
 - No update when
$$\phi < \frac{2 \times T_{\text{max. propagation delay}}}{T}$$

Comparison of two ways for Peer Discovery

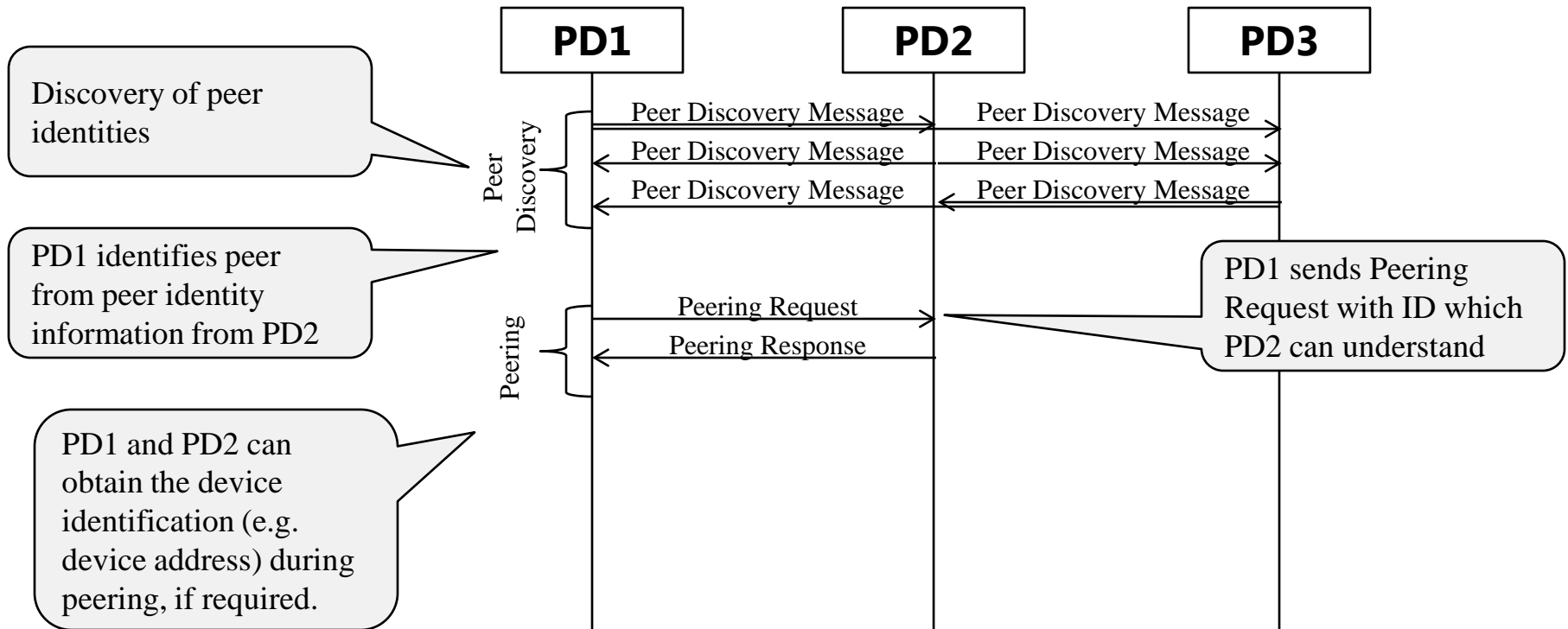
- Peer discovery after peering



* Terminologies are temporarily used for explanation

Comparison of two ways for Peer Discovery

- Peer discovery before peering



* Terminologies are temporally used for explanation