

**Project: IEEE P802.15 Working Group for Wireless Personal Area Networks (WPANs)**

**Submission Title:** [FEC simulation results]

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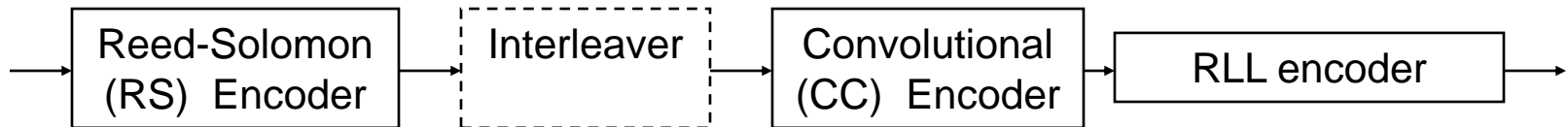
**Abstract:** Provides details on FEC simulations for VLC

**Purpose:** [Contribution to IEEE 802.15.7 VLC TG]

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# TX System block diagram



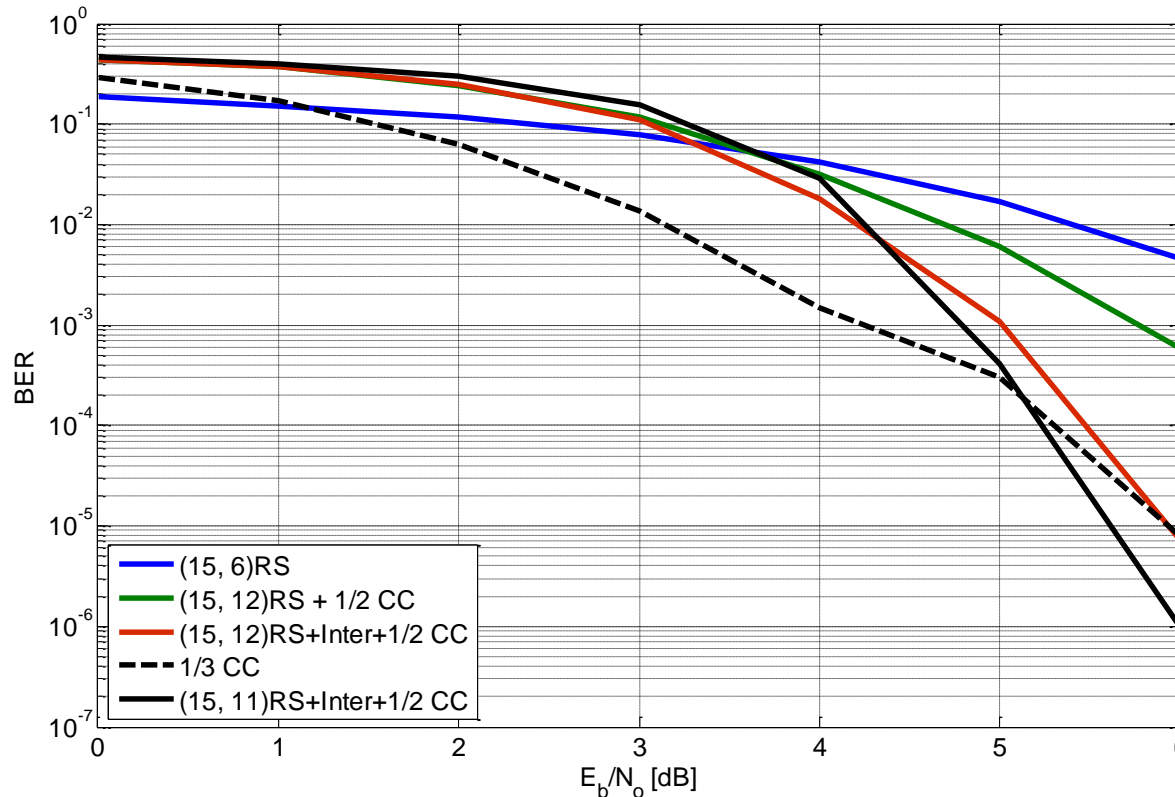
Hard-decision decoding is used in the RS and the CC decoders.

Random interleaver is used in these simulations. In practice, other interleavers such as block interleavers can also be considered.

Manchester RLL encoder is considered.

# Does the convolutional code (CC) need an interleaver?

AWGN Channel



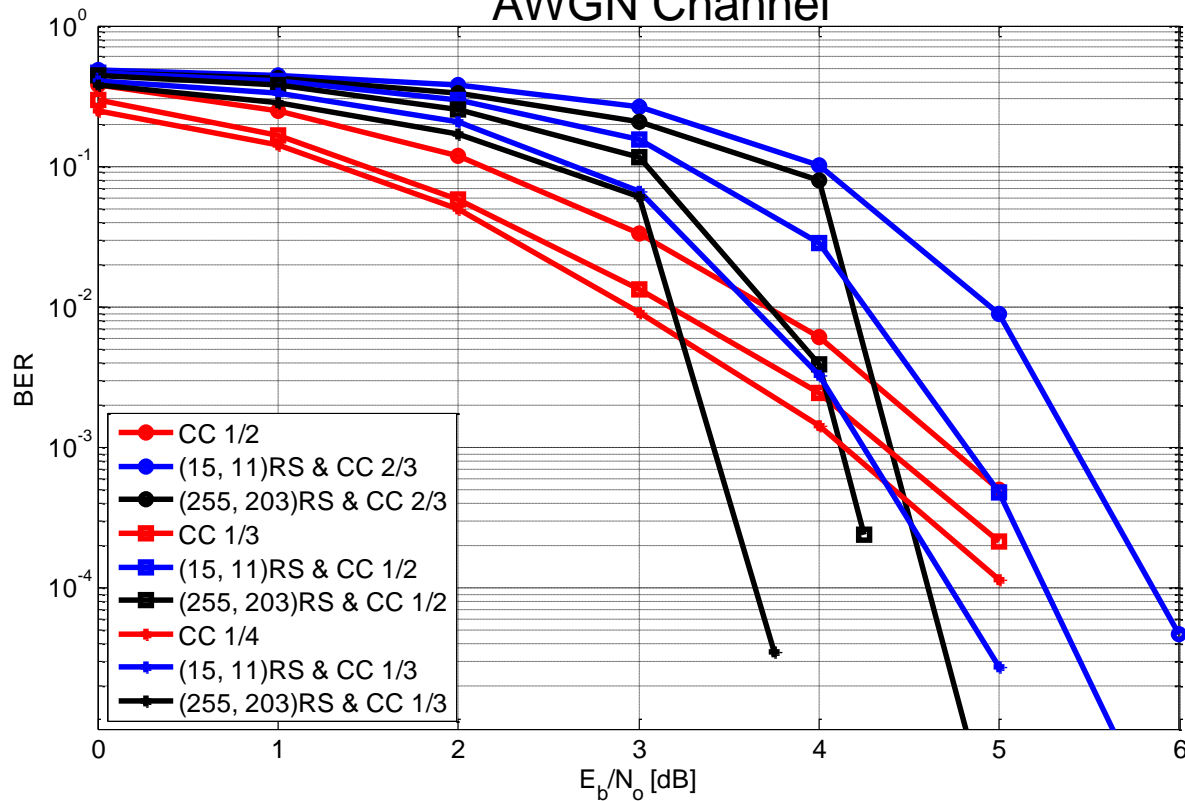
Yes, an interleaver between the RS code and the CC achieves >1dB gain.

No need for interleaver at the output of the CC.

It is better to reduce coding rate of CC instead of RS, given a choice to attain same coding rate. RS code is also better for high rate implementation from complexity standpoint.

# CC vs. RS+CC

## AWGN Channel



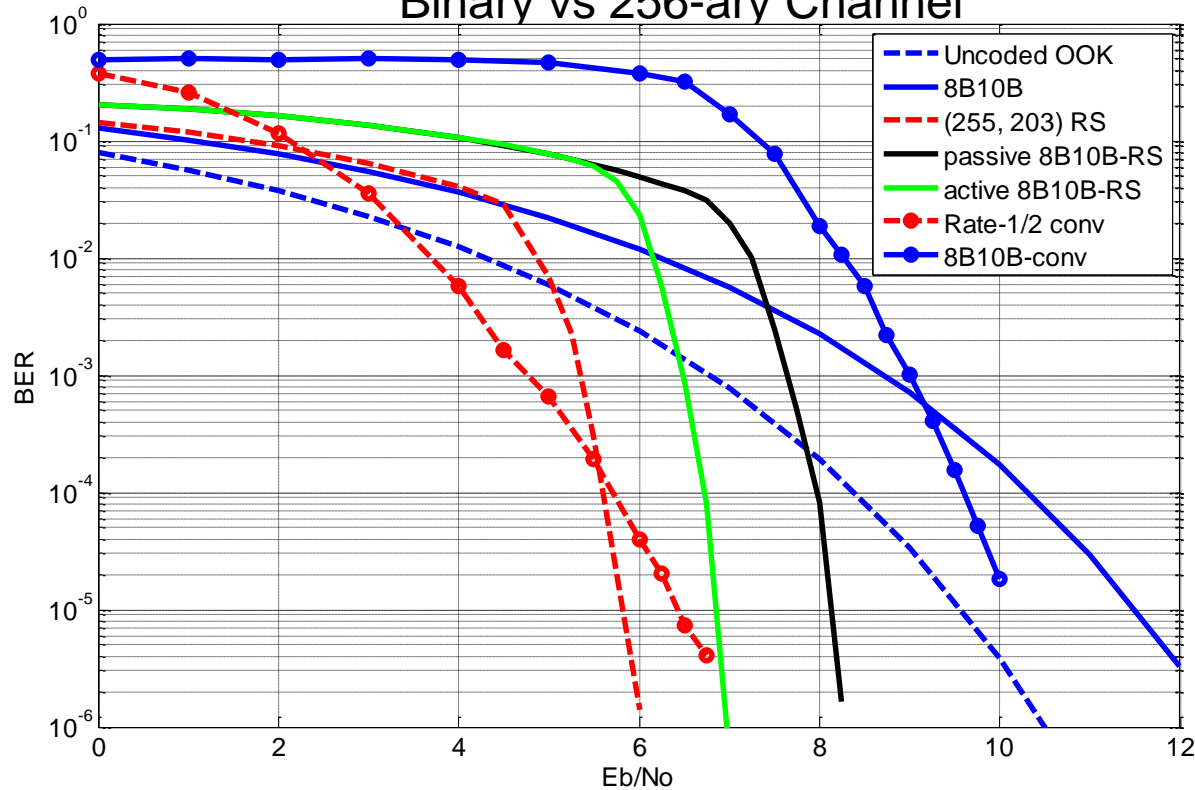
Operating in  $BER > 1e-3$ : Use CC (without RS)

Operating in  $BER < 1e-3$  (practical for  $PER < 1e-1$ ): Use RS+CC

Operating at coding rates  $< 1/4$ : Use RS+CC

# CC vs. RS

## Binary vs 256-ary Channel

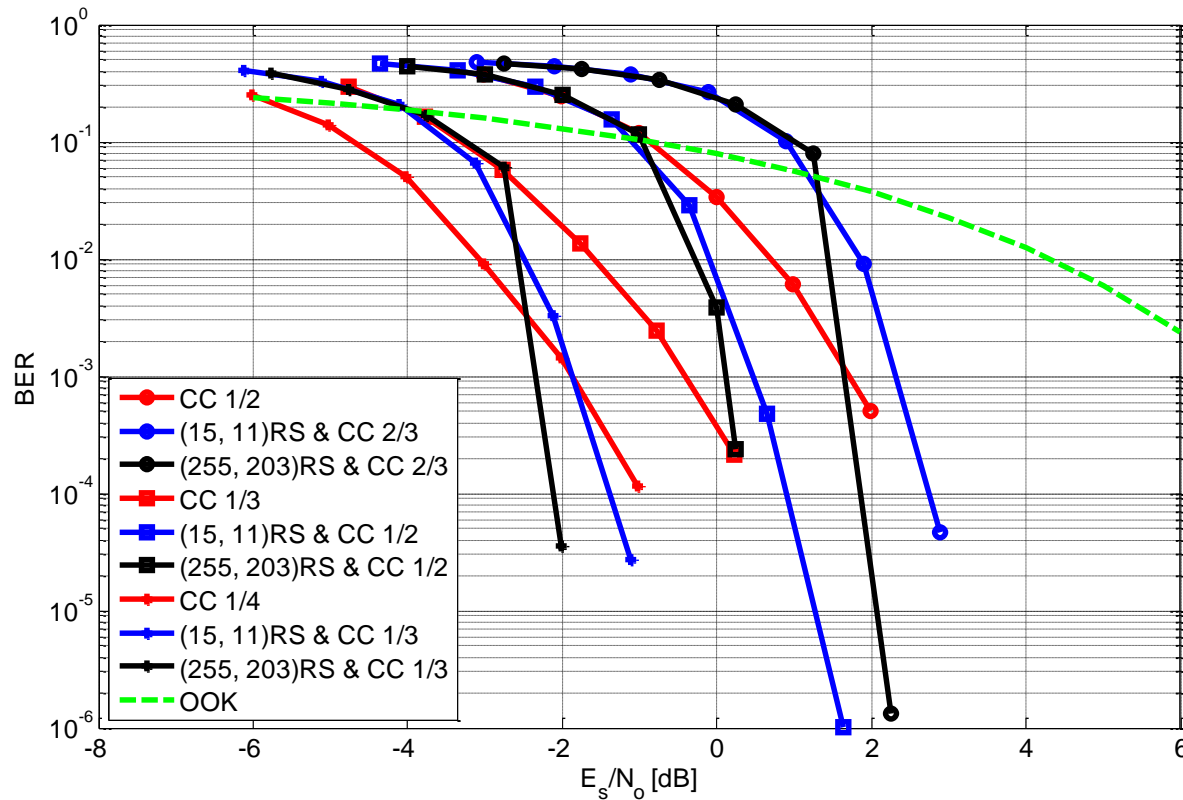


PHY I (Manchester): RS+CC

PHY I (4B6B): RS on GF(16)

PHY II : RS on GF(256)

# Dynamic SNR Range



Assuming same optical rate, 2dB between rate-1/2 and rate-1/3, 2dB between rate-1/3 and rate-1/4, and >8dB between rate-1/4 and un-encoded OOK.

More simulations are needed to verify the dynamic range for all proposed MCSs, and to guarantee that there is no overlapping MCSs.

# Recommendations

## PHY I :

- use RS + CC for Manchester coding
- Given a choice, reduce coding rate in CC over RS code.
- Avoid use of repetition coding.
- Use RS only for 4B6B code
- Use GF(16) for RS code (short packet sizes)

## PHY II

- Use RS GF(256) with 4B6B code
- Use RS GF(256) with 8B10B code

# Further simulations on RS/CC choices



# Assumptions

Random interleaver between CC and RS code.

RS decoder uses erasure information when available via RLL decoder.

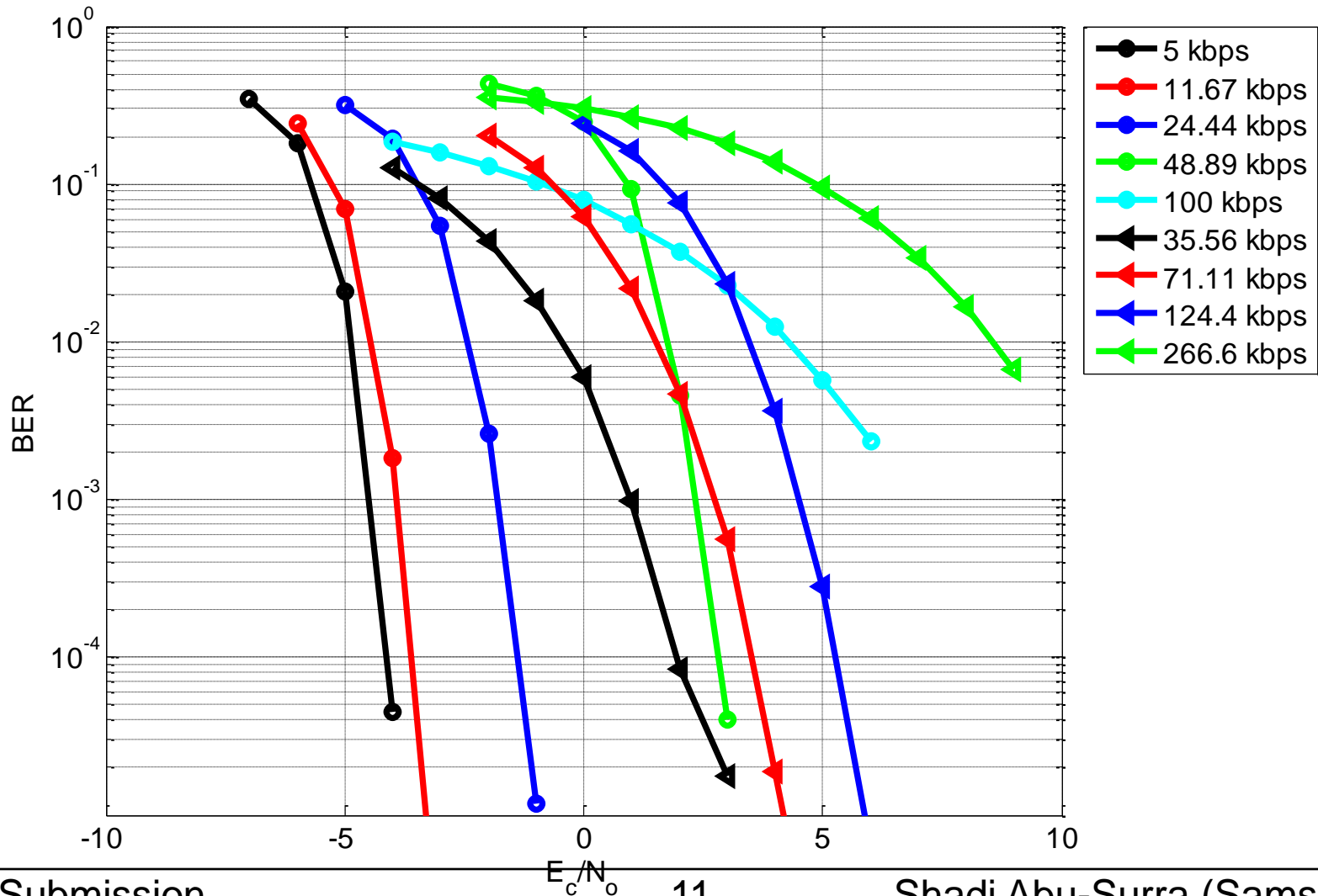
VPM and OOK have same energy on the average

- VPM – assumes 50% duty cycle. i.e. 2-PPM in time period  $T$
- OOK – assumes 100% duty cycle when '1' over time period  $T$ , 0% when '0'

# New rate tables (PHY I) – Clocks multiple of 2

	Optical rate	Modulation	Line coding	FEC	Data rate
PHY I	200 kHz	OOK	Manchester	(15, 3) RS + 1/4 CC	5 kbps
				(15, 7) RS + 1/4 CC	11.67 kbps
				(15, 11) RS + 1/3 CC	24.44 kbps
				(15, 11) RS + 2/3 CC	48.89 kbps
				1	100 kbps
	400 kHz	VPM	4B6B	(15, 2) RS	35.56 kbps
				(15, 4) RS	71.11 kbps
				(15, 7) RS	124.4 kbps
				1	266.6 kbps

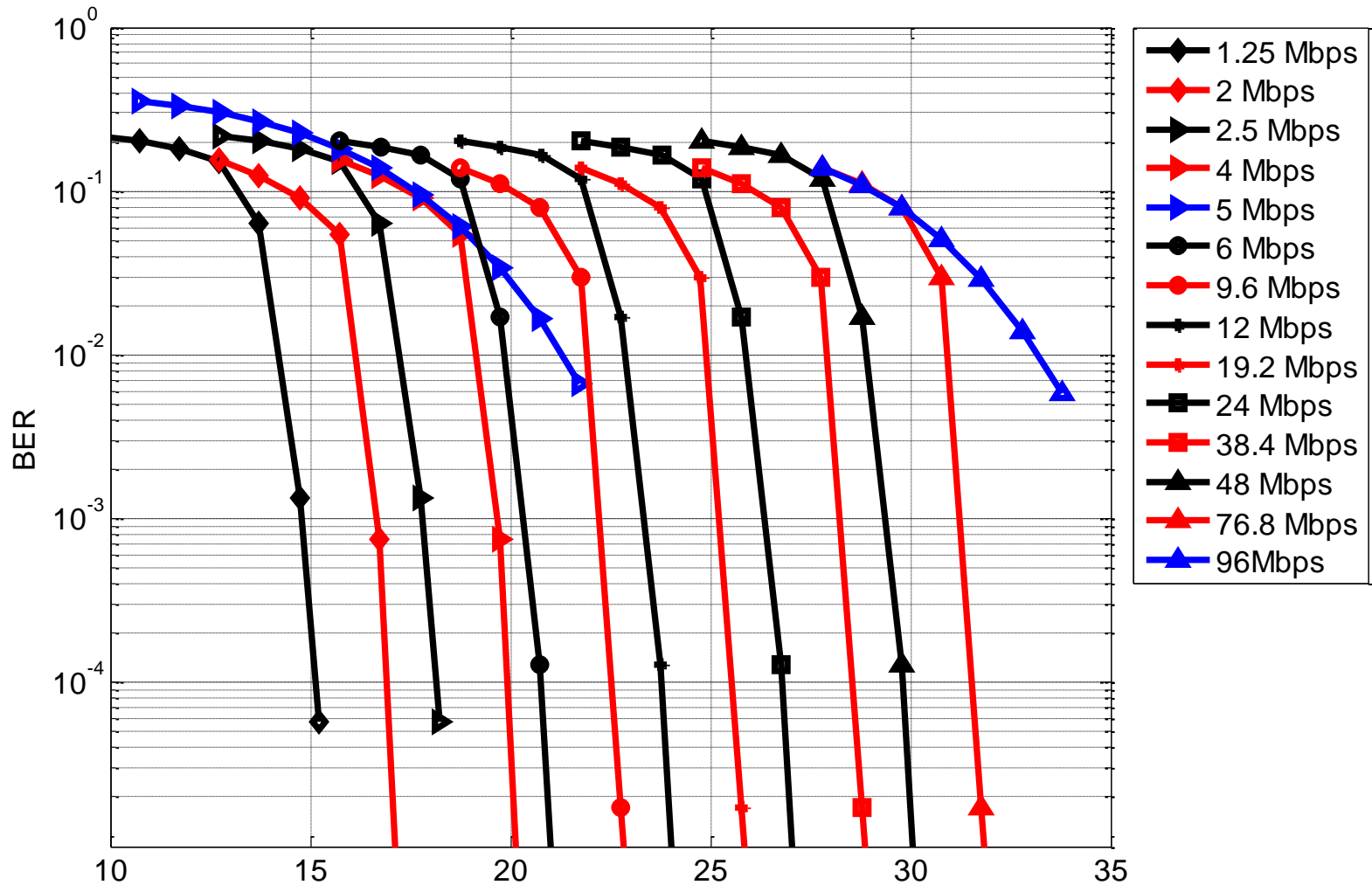
# PHY I simulation results



# PHY II

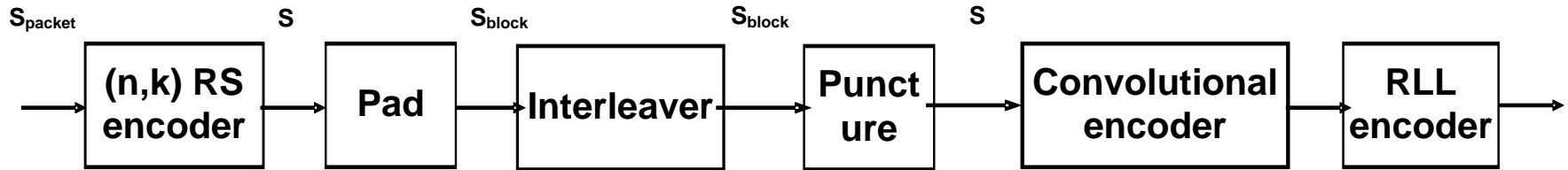
	<b>Modulation</b>	<b>Line coding</b>	<b>FEC</b>	<b>Optical rate</b>	<b>Data rate</b>
<b>PHY II</b>	<b>VPM</b>	<b>4B6B</b>	<b>(64, 32) RS</b>	<b>3.75 MHz</b>	<b>1.25 Mbps</b>
			<b>(160, 128) RS</b>		<b>2 Mbps</b>
			<b>(64, 32) RS</b>	<b>7.5 MHz</b>	<b>2.5 Mbps</b>
			<b>(160, 128) RS</b>		<b>4 Mbps</b>
			<b>1</b>		<b>5 Mbps</b>
	<b>OOK</b>	<b>8B10B</b>	<b>(64, 32) RS</b>	<b>15 MHz</b>	<b>6 Mbps</b>
			<b>(160, 128) RS</b>		<b>9.6 Mbps</b>
			<b>(64, 32) RS</b>	<b>30 MHz</b>	<b>12 Mbps</b>
			<b>(160, 128) RS</b>		<b>19.2 Mbps</b>
			<b>(64, 32) RS</b>	<b>60 MHz</b>	<b>24 Mbps</b>
			<b>(160, 128) RS</b>		<b>38.4 Mbps</b>
			<b>(64, 32) RS</b>	<b>120 MHz</b>	<b>48 Mbps</b>
			<b>(160, 128) RS</b>		<b>76.8 Mbps</b>
			<b>1</b>		<b>96 Mbps</b>

# PHY II simulation results



Interleaver design  
(Structured interleaver instead of a random  
interleaver)

# Interleaver design



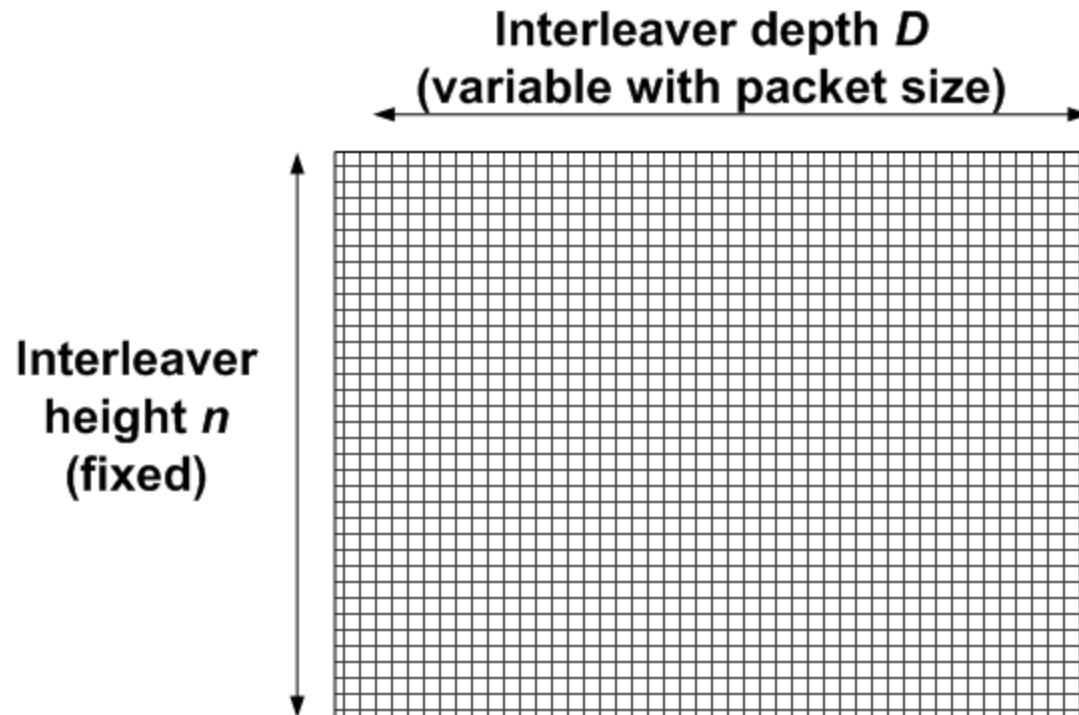
Interleaver is a block interleaver.

Fixed height  $n$

Flexible depth  $D$ , dependent on the packet size

- Optimized for short packet sizes to eliminate padding

# Block interleaver design









# Interleaver details

$n$  : RS codeword length.

$k$  : number of information data symbols in a RS codeword.

$q$  : Number of elements in the Galois field : GF( $q$ )

$L_{packet}$  : Input packet size in bytes

$S_{packet}$  : number of symbols at the input of the RS encoder

$S$  : number of symbols from the output of the shortened RS encoder.

$S_{block}$  : The size of the interleaver used

$D$  : the interleaving depth.

$i$  : ordered indices take the values  $0, 1, \dots, S_{block}-1$ .

$l(i)$  : interleaved indices.

$\rho$  : number of zero symbols

$t$  : ordered indices take the values  $0, 1, \dots, \rho$ .

$z(t)$ : locations of the bits to be punctured at the output of the interleaver before transmission

$$S_{packet} = \left\lceil \frac{L_{packet} * 8}{\log_2(q)} \right\rceil$$

$$S = n * \left\lceil \frac{S_{packet}}{k} \right\rceil - (k - (S_{packet} \bmod k))$$

When  $S$  is a multiple of  $n$ , there is no padding or puncturing, The RS encoder produces  $n$  symbols for every codeword, which are then sent to the interleaver. When  $S$  is not multiple of  $n$ ,  $p$  remaining symbols for the last codeword are padded with zero symbols and encoded to produce  $n$  symbols. After interleaving, the zero symbols are punctured out and are not transmitted.

$$D = \left\lfloor \frac{S}{n} \right\rfloor$$

$$S_{block} = n * D.$$

$$p = n - (S \bmod n)$$

Interleaver:

$$l(i) = (i \bmod D) * n + \left\lfloor \frac{i}{D} \right\rfloor ; \text{ for } i = 0, 1, \dots, (S_{block} - 1)$$

Locations to be punctured:

$$z(t) = (n - p + 1) * D + t * D - 1; \text{ for } t = 0, 1, \dots, p - 1$$

# Simulation assumptions

(15, 12) RS +  $\frac{1}{2}$  CC

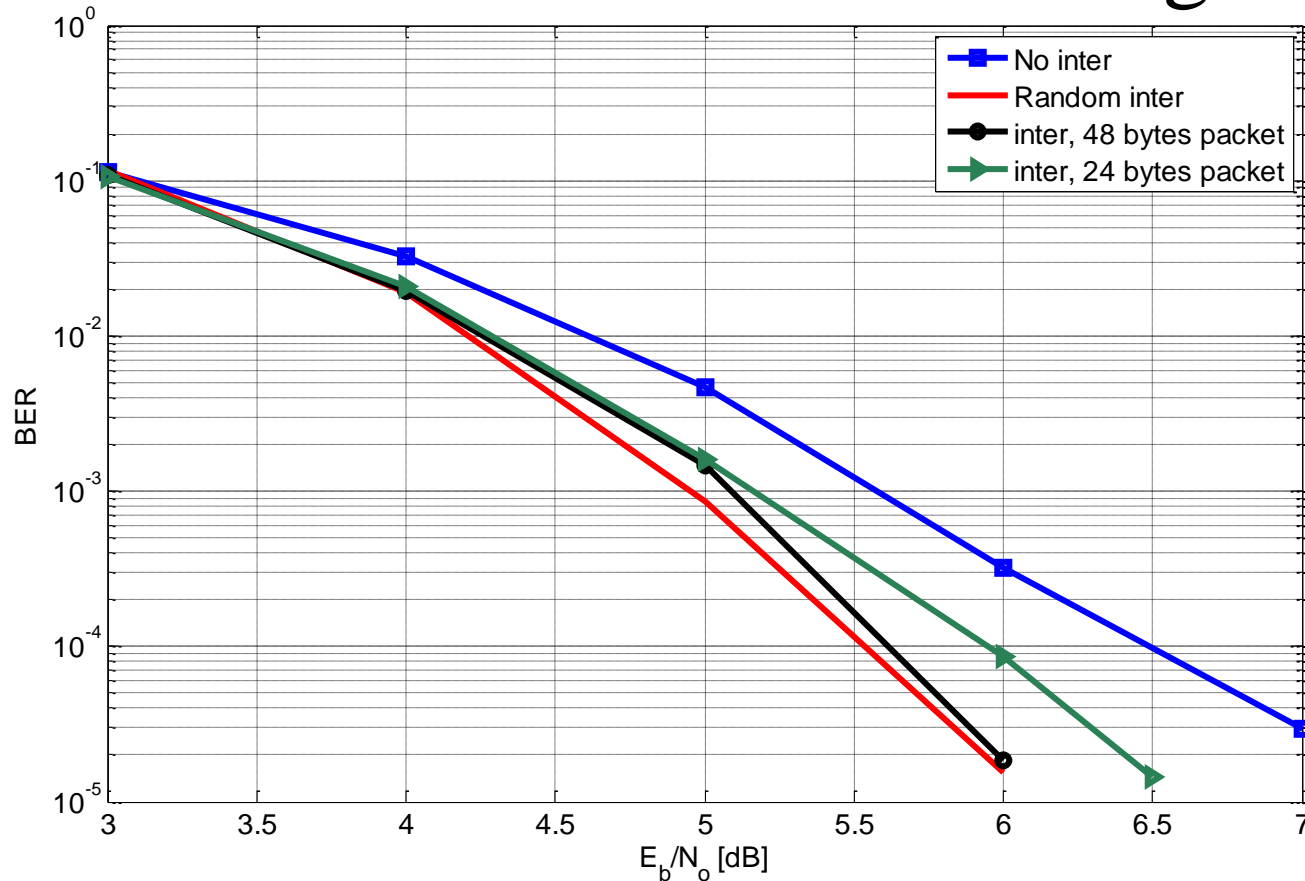
Hard-decision decoding is used in the RS and the CC decoders.

Manchester channel DC-balance encoder is considered.

AWGN channel

Small packet sizes

# Benefits of Interleaving



Further gains can be expected for fading channels if increased burst errors over multiple RS symbols

# Usage

The length of the packet is communicated to the receiver in the header so that the receiver can adaptively adjust the interleaver based on the packet sizes.

When the data rates corresponding to the robust transmissions using the concatenated codes are used, the header shall also be interleaved according to the above procedure.

Since the length of the header is fixed, the receiver can deinterleave the header without explicit transmission of the header length.

# Summary

Symbol interleaver between RS code and CC can give performance advantages, providing increased reliability for applications such as vehicular communication

Interleaver designed for short packet sizes – no overhead due to padding

Simple, block interleaver structure

- Feasible for implementation

Interleaver works on very low symbol speed (50 KHz : 200 KHz clock, 4 bits/symbol)



# Recommendations

Adopt the proposed symbol interleaver between the RS and CC code for applicable data rates