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ABSTRACT

The steps, approaches, and tools of the collective inquiry process posited by a Regional Environmental Learning System is the subject of this volume. Approaches discussed include: (1) Charlotte approach, (2) AR&F/Battelle approach, and (3) the Washington State approach. Tools for collective inquiry are described and field tests are discussed. References and four topical bibliographies are provided. An appendix presents a computer program for modeling technique proposed for use in collective inquiry. (RE)

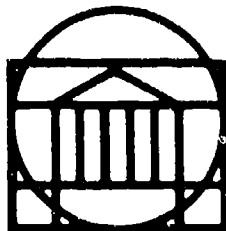
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RESEARCH LABORATORIES FOR THE ENGINEERING SCIENCES



SCHOOL OF ENGINEERING AND APPLIED SCIENCE

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DEVELOPMENT OF AN INTERPRETIVE STRUCTURAL MODEL
AND STRATEGIES FOR IMPLEMENTATION BASED ON
DESCRIPTIVE AND PRESCRIPTIVE ANALYSIS OF
RESOURCES FOR ENVIRONMENTAL
EDUCATION/STUDIES

VOLUME IV

CONDUCTING COLLECTIVE INQUIRY

Submitted to:

Office of Environmental Education
Department of Health, Education and Welfare
400 Maryland Avenue, S.W.
FOB #6, Room 2025
Washington, D. C. 20202

Submitted by:

John N. Warfield

Report No. UVA/522032/EE79/124

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RESEARCH LABORATORIES FOR THE ENGINEERING SCIENCES

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A SOURCEBOOK FOR THE DESIGN OF A
REGIONAL ENVIRONMENTAL LEARNING SYSTEM
VOLUME IV

CONDUCTING COLLECTIVE INQUIRY

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A SOURCEBOOK FOR THE DESIGN
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VOLUME IV: CONDUCTING COLLECTIVE INQUIRY

PREFACE

This is one of six Volumes of a report which, collectively, is intended to be a Sourcebook for the Design of a Regional Environmental Learning System. The report was prepared under Contract 300-700-4028 with the Office of Environmental Education.

This six-volume report presumes some background concerning the concept of a Regional Environmental Learning System, and with environmental education as a whole. Considerable relevant background was supplied in Volume 9 of the 4th Quarterly Report (A Descriptive Analysis of Environmental Education) and in the 5th Quarterly Report (Conceptual Basis for the Design of Regional Environmental Learning Systems), both of which are available from the Office of Environmental Education.

Volume 1 contains an Overview of the Sourcebook, with short summaries of the other Volumes.

VOLUME IV
CONDUCTING COLLECTIVE INQUIRY

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August 1979

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VOLUME IV

CONDUCTING COLLECTIVE INQUIRY

EXECUTIVE SUMMARY

As indicated in previous volumes, especially Volumes I and III, learning about one's environment encompasses a large number of factors, factors like resources, conservation, pollution, economics, demography, urban and regional planning, technology, and transportation. Understanding all these factors, their interrelations, and their implications for planning and managing one's future cannot be done alone. Involvement in collective inquiries with other participants with a variety of skills, knowledge, experiences, perceptions, and values is required. For such inquiries to be efficient and productive, means for managing the generation, structuring, communication, and documentation of ideas are needed, in other words, tools for idea management are needed. Through the use of such tools, the process of collective inquiry may enable individuals and groups to understand and improve their environments.

In addition to helping individuals, the process of collective inquiry is useful to teachers and others concerned with managing or facilitating learning. Collective inquiries can be exceedingly useful for achieving environmental education "that synthesizes and integrates pertinent subject matter across and between a variety of disciplines". In addition to idea management tools, approaches to planning and managing the conduct of collective inquiries are needed.

This volume provides some tools and approaches for conducting collective inquiries. Although an attempt is made to provide the reader with enough information to carry out an inquiry, the references provide more information and may be referred to if the information provided in this volume needs to be augmented.

As used here, "approaches to collective inquiries" mean suggestions and ideas for planning and managing the process of an inquiry. "Tools" mean particular ways to facilitate a particular step in the process.

COLLECTIVE INQUIRY STEPS

Why might one want to conduct a collective inquiry? The answer depends on the stage at which a collection of participants is dealing with an environmental issue. For different stages, different purposes are served. Some examples are:

- to raise awareness of environmental issues and problems or opportunities,
- to define or explain problems and issues concerning the environment,
- to share perceptions or facts,
- to discover values, beliefs, or assumptions,
- to set goals,
- to develop candidate policies, plans, or programs to achieve goals,
- to assess the worth of candidate policies, plans, or programs,
- to choose policies, plans, or programs,
- to develop approaches for implementation of choices,
- to identify who have what tasks to accomplish for implementation,
- to evaluate performance of policies, plans, or programs.

To achieve any of these purposes, it is helpful to give careful consideration to the following steps for conducting an inquiry.

1. At the outset it is important to define clearly the purpose(s) of the collective inquiry. This is needed to insure that the participants work together toward a common goal. If it is not done, some participants may try to engage in activities that are incompatible with the activities of other participants during an inquiry. As an example, some may try to develop candidate policies addressed to an environmental issue while other participants at the same time are trying to assess the worth of policies. The former group is engaged in a process requiring conceptualization, while the latter is engaged in analysis. Mixing these kinds of activities at the same time often leads to conflicts and/or failure. An approach more likely to succeed would separate the activities in time by first generating candidate policies and then assessing the worth of each candidate. Defining the purpose clearly will help the group of participants to focus on a single task at any given point in time and, thereby, improve their chances for success. [1] [2]

Another reason for defining the purpose clearly is to help insure that appropriate participants, appropriate arrangements, and appropriate tools are obtained and involved in the inquiry. More is said on these below.

2. It is important to plan well for obtaining good participants and arrangements for an inquiry. Too often collective inquiries fail because not enough attention is given to one or both of these. Table I indicates some of the main points requiring attention.

TABLE I
PLANNING COLLECTIVE INQUIRIES

- o GET GOOD PARTICIPANTS
- o INSURE GOOD PHYSICAL ARRANGEMENTS
- o INSURE GOOD STAFFING
- o PREPARE SOURCE MATERIAL FOR THE INQUIRY
- o INSURE GOOD RECEPTION AND ORIENTATION FOR PARTICIPANTS
- o INSURE CONTRIBUTIONS ARE RECORDED
- o PREVENT DISRUPTIONS TO ORDERLY THINKING
- o SUMMARIZE AND FEED BACK RESULTS
- o DOCUMENT AND ANALYZE INQUIRY PROCESS AND RESULTS
- o DISTRIBUTE FINAL REPORT

Good participants have the right qualifications, are willing to contribute, and have an interest in the issue under consideration,

Good physical arrangements have adequate blackboards, flip charts, sound-proofing, temperature control, audio-visual equipment, sleeping accommodations, and everything else needed to facilitate the productivity of the effort.

Good staffing insures adequate secretarial help, receptionists, etc., as well as resource persons with knowledge of the content of the issue in the inquiry.

In addition to resource persons, the content of the inquiry may be aided by providing "issue" or "background" papers for the participants to be read before the meeting and/or referred to at the meeting. Relevant "facts books" may be prepared also for use at the meeting.

A good reception for the participants reinforces in their minds that their contributions are important for the success of the meeting. Good orientation focuses their thinking on the purposes for the meeting, the processes to be employed to achieve those purposes, and the schedule to be followed. Good orientation also helps insure that no participants' contributions are lost because the participants are lost.

Recording the contributions of the participants insures that after the meeting one can refer to the record for whatever reason.

Preventing disruptions of the participants' work is important whether from telephone calls, the need for minor decisions, lobbyists for special interests, or whatever.

Summarizing and feeding back results is useful during and after the meeting. Errors can be corrected. Results can be improved through iteration. The amount of progress can be assessed.

Documenting the process and the results of the inquiry serves several purposes. It allows analysis both by the participants and by others not involved. It provides a springboard for later stages of inquiry. And it provides valuable information for improving the tools for collective inquiry. It allows others not involved in the process to share in the learning achieved.

None of this can be achieved without careful and persistent attention to detailed planning.

3. After the purpose for the inquiry is defined and the participants and arrangements are fixed, attention can turn to the inquiry process as such. In

most cases, the participants first need to identify the elements of the issue being addressed. For example, if the purpose of the collective inquiry is to define an environmental problem, different participants can probably see different contributing factors. These factors put together make up the environmental problem. Each is an element of the problem; none is the whole problem. Through the process of generating the elements, the participants gain an improved and shared understanding of the whole problem.

4. After the elements of the issue under consideration have been generated, it is natural in most circumstances to begin to consider what the relations are among the elements. Transitive relations are useful for developing chains of relations. For example, suppose the elements of the issue are the contributing factors of an environmental problem. Consider the relation "contributes to". Then if factor A contributes to factor B, and if factor B contributes to factor C, then, by transitivity, factor A contributes to factor C.

Choosing the relation to use is an important decision in most inquiries. Consider for a moment the differences in the following: "contributes to"; "is more important than"; "is more easily managed than"; and "occurs before". Depending on which one of these is chosen, in most cases, different structures of related elements will result. It is sometimes useful to consider more than one relation with a set of elements, but not at the same time.

5. After a contextual relation is chosen, the participants may begin to consider pairs of elements to decide whether the relation holds between them. In the process of deciding whether the relation holds, two situations can arise. The first occurs when the group decides the relation holds in both directions. For example, the group might decide the chicken is more important than the egg and that the egg is more important than the chicken, in which case, it seems natural to agree they are of equal importance. Such situations can occur with

many relations, e.g. "contributes to". The second situation occurs when the group finds it impossible to decide that the relation holds in either direction, e.g. A is not more important than B, and B is not more important than A. There is no requirement that the relation hold in either direction, and no one should be concerned if it does not.

In the process of deciding whether the relation holds between two elements, the group often develops an improved definition or understanding of the elements or the relation. They also often gain a better understanding of other participants' views about the elements or the values, beliefs, or perceptions of other participants. These improved understandings are among the main, beneficial outcomes of collective inquiries.

This step ends when all the pairs of elements have been tested with the chosen relation either directly, or by inference using transitivity. By using the latter, the group usually does not have to decide a large number of cases. That is, if A relates to B and B relates to C, then it can be inferred immediately that A relates to C without asking the group to decide. The use of a computer greatly facilitates making such inferences.

6. Once the relation has been tested with all pairs of elements and the results have been displayed, then the group can analyze the results to determine if they seem intuitively reasonable or, failing that, determine an explanation for the resulting structure of related elements. If this is possible, new insights about the issue under consideration are often gained.

If the results do not seem reasonable and no explanation can be found, then the group may choose to amend the structure to achieve one that does seem intuitively reasonable or that can be explained. One of the ways to amend the structure is to reconsider decisions where the group was almost evenly divided on whether the

relation held or not. Through additional discussion, the group may decide to reverse the decision. This will alter the structure. Another way to amend the structure is to recognize and decide that a major portion of the structure belongs in a different place within the overall structure. In these cases and others the amendment process is facilitated by the use of a computer.

Several different structures may be achieved by using different elements, or different relations, or different groups of participants. In these cases and in others, there may be a need to rank or assess the worth of the various structures.

7. When the collective inquiry is completed, documentation of the results and the process is useful for several reasons. It allows others who were not participants to gain some understanding of how the results were achieved. If they see weak points, these can be strengthened. Another benefit is to enable other participants to build on the results. A third use is for reference by the participants or others who may want to look back at the process or the results.

After documentation, the results should be disseminated to all of those who need to be aware of or to use the results for follow-up actions.

APPROACHES TO COLLECTIVE INQUIRIES

Approaches to collective inquiries considers the questions of who, where, what, and when. Who will participate? Who will serve as resource persons? Who will the staff be to support the effort? Where will the meetings be held? Where will those involved stay? What equipment will be needed? What resource materials will be needed? When will the meetings begin? When will they end?

Answers to such questions will depend on a number of factors. Chief among them will be how soon the results must be obtained and how many of the objectives given in the previous section are to be sought. Two important constraints on the

approach to be taken are the funding and staffing available for the inquiry. These kinds of factors and constraints are important for deciding whether to undertake a collective inquiry at all and, if that decision is in the affirmative, what approach to take.

In this section, three approaches are described. All of them have been used successfully in actual inquiries. They differ mainly along one dimension, the time required to conduct the inquiry. The first requires about forty-eight hours; the second, about four to six weeks; and the third may take more than a year. It is not intended to imply that inquiries could not take more or less time. These are examples.

For a particular inquiry, combinations of the three might be appropriate. For example, the third approach might include a sequence of inquiries using the first approach. It is intended that the descriptions of these three approaches may help to suggest to the reader how to design her/his own approach to satisfy the need at hand.

1. CHARETTE APPROACH

The term "charette" comes from the early French schools of architecture and refers to the approach used by students for design projects. They worked around the clock during a short period of time to complete their designs.

This approach was adopted by the Forest Service for land management planning for the Shawnee National Forest in Illinois. [3] Simply stated this was a joint effort of Forest Service personnel and concerned citizens to develop the management plan for the Shawnee National Forest through the year 1985. By working together, collectively, they exchanged ideas, analyzed the available and potential natural resources, considered the constraints and demands, and developed what they considered to be the best management plan they could. The work began on Friday at 5:30 pm and closed the following Sunday at 5:00 pm. The work continued past midnight for some of the participants.

Although the inquiry only lasted about forty-eight hours, the time required to prepare for it and to process the results was considerably longer. Interested and capable citizens were identified and recruited to participate in the meeting. A study guide was prepared that included salient facts on the issues that would be covered at the meeting. The study guide was mailed to the participants followed by a reminder letter to study the material, attend the meeting, and submit comments on the issues. The comments were analyzed, summarized, and fed back to the participants. Arrangements were made to pre-register the participants by mail. Motel arrangements were made and checked to be sure they were adequate. Audio-visual equipment was obtained and inspected to insure it was working. Arrangements were made and checked to insure there were enough and good meeting rooms, blackboards, chalk, flip-charts, markers, etc. Secretarial services were obtained for the meeting and for processing the results. Resource persons and group leaders were identified, selected, recruited, confirmed, and informed. The group leaders were chosen on the basis they were known to be interested, open minded, and able to encourage broad coverage and participation. Work assignments and schedules were prepared for supporting staff including scribes for each work group. A briefing was given to the supporting staff. Major issues or projects that were to be covered in the Charette were selected, and briefing papers prepared for each of them. Arrangements for registering and welcoming the participants at the meeting were made including assignments to work groups. A speech was prepared for the forest supervisor to give to welcome the participants, and to provide motivation and orientation for their work.

Following the meeting, reference and other materials were retrieved from the participants. Reports from the work groups were collected. The scribes and other relevant staff remained at the site of the Charette to analyze the written reports

and to clarify or elaborate on the language when it was deemed desirable. Still at the site, a draft of the summary report was prepared. Later the final report was prepared and disseminated.

Additional arrangements were made for conducting the meetings. Scribes were asked to report on the process of the meetings and to record agreed-upon definitions or constraints, rejected ideas, accepted ideas, and conclusions. Plans were made to minimize interruptions to insure the maximum benefit from the participants' contributions. A plenary session was planned for Sunday afternoon where the group leaders delivered oral reports on the results of their meetings and the forest supervisor thanked the participants for their help.

A summary of the main steps in the Charette approach is given in Table II. Helpful tools for use during such meetings are given in the section on Tools for Collective Inquiries.

TABLE II
CHARETTE PROCESS

1. Prepare a brochure to serve as a study guide for the participants. Include the salient facts on the issues they will address.
2. Identify and recruit relevant participants and mail them the brochure.
3. Send follow up letter to remind participants to study material, attend sessions, and provide comments.
4. Make arrangements to process comments i.e. analyze, summarize, and disseminate.
5. Make arrangements for registering participants by mail.
6. Insure that motel arrangements are in order including audio-visual equipment, auditorium, conference rooms, flip charts, blackboards, secretarial services, computer services and telephone lines, if needed.
7. Select, recruit, confirm, and inform resource persons.

8. Develop list of potential projects to be discussed.
9. Select, recruit, confirm, and inform Group Leaders. Select on basis of personal knowledge of abilities, interest in project area, open-mindedness, and ability to encourage broad coverage and participation.
10. Arrange for registration and reception at meeting. Prepare name tags.
11. Prepare work assignments and schedules for support staff, including Scribes for each work group.
12. Prepare motivation and orientation speech.
13. Give briefing to support staff.
14. Prepare reference materials to be available at meeting.
15. Register participants and assign them to work groups.
16. Deliver motivation and orientation presentation.
17. Introduce Work Group Leaders, Scribes, and resource persons, and break into Work Groups.
18. Insure that Scribes record process, rejected ideas, and accepted ideas.
19. Prevent special-interest "floaters" from traveling from group to group to carry their proposals.
20. Convene final plenary session and make opening remarks on its purpose.
21. Have Group Leaders deliver oral and written reports to top manager.
22. Have top manager thank participants and all others involved for their work.
23. Request participants and others to return reference and other materials and adjourn meeting.
24. Collect registration list and Work Group reports.
25. Hold Scribes and other relevant staff at Charette site to analyze written reports and elaborate on language if needed.
26. At site, prepare draft of summary report.
27. Prepare final report.
28. Distribute final report to participants and other appropriate persons.

2. AT&T/BATTELLE APPROACH

The approach described in this section was first used in Columbus, Ohio, to assess that community's expectations for a public school curriculum. The development of the approach was sponsored by AT&T and Battelle Memorial Institute, Battelle Columbus Laboratory (BCL). [4]

The approach is iterative. Information is requested from participants. The information received is analyzed, categorized, and summarized. The results are then returned to the participants with a request for additional information that elaborates or refines the previous information. A summary report is prepared and distributed.

In the implementation of this approach in Columbus, one of the first steps taken was to obtain the support of the Board of Education for the assessment and to have it appoint a coordinator for the assessment. This was important for conducting the assessment and for the acceptance of the results. Concern for "legitimacy" is important in most collective inquiries.

In developing the approach, a plan was established for recruiting participants for the assessment who represented a cross-section of those who were concerned about the learning outcomes achieved by students. Among those concerned about learning outcomes are students, parents, employers, teachers, and other members of the community. It was important that the plan for obtaining participants emphasized this need for representative participants. There were about 1700 participants (including group leaders) from 14 school areas in the Columbus assessment. The first question used to generate information by the participants was "What does a person leaving high school need to know or be able to do to 'make it' in life?" (This terminology was employed because it was recognized that a number of students leave school before they obtain diplomas.)

Two types of supporting staff were involved, trainers and group leaders. The trainers were graduate students at Ohio State University and the group leaders were teachers, students, parents, businessmen, and other members of the community. The means for selecting group leaders in the 14 school areas is shown in Figure 1. There were 8 trainers and 257 group leaders. The trainers provided training for the group leaders, gave them advice when it was needed, monitored their progress, and helped insure the assessment proceeded on schedule. The group leaders recruited the participants, made arrangements for meeting locations, notified and reminded participants of meeting times, distributed necessary materials at the meeting, led the meetings, collected results, and fed them back to the Battelle staff coordinating the assessment.

| | STUDENTS | TEACHERS | COMMUNITY MEMBERS | BUSINESSMEN |
|--------------------------------------|--------------------------------------------------------------------------------------------------------------------------------------------------------------------|------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| DISTRIBUTION OF LEADERS | 3 LEADERS REPRESENTING EACH ATTENDANCE AREA | 5 LEADERS REPRESENTING EACH ATTENDANCE AREA | 2 LEADERS REPRESENTING EACH PUBLIC SCHOOL | 12 TO 20 LEADERS REPRESENTING THE ENTIRE COMMUNITY |
| RESPONSIBILITY FOR SELECTING LEADERS | EACH HIGH SCHOOL PRINCIPAL | EACH HIGH SCHOOL PRINCIPAL | EACH SCHOOL PRINCIPAL | PROJECT DIRECTOR |
| SELECTION CRITERIA | <ul style="list-style-type: none"> o JUNIOR OR SENIOR o BOTH SEXES o INTEREST IN TOPIC AND WILLINGNESS TO ORGANIZE OTHER STUDENTS | <ul style="list-style-type: none"> o FROM FEEDER SCHOOLS AS WELL AS HIGH SCHOOL o BOTH SEXES o INTEREST IN CURRICULUM o FROM VARIETY OF FIELDS | <ul style="list-style-type: none"> o PARENTS AND NON-PARENTS o BOTH SEXES o NOT EMPLOYED BY CITY SCHOOLS o WILLINGNESS TO SEE AN ACTIVITY THROUGH TO COMPLETION o BASIC READING, WRITING, AND SPEAKING SKILLS | <ul style="list-style-type: none"> o FAMILIARITY WITH JOB REQUIREMENTS IN HIS BUSINESS o DIRECTLY INVOLVED IN EMPLOYEE TRAINING o REPRESENTATIVES FROM A VARIETY OF TRADES, PROFESSIONS, BUSINESSES, AND SERVICES |

FIGURE 1. SELECTION OF GROUP LEADERS

TABLE III

TOPICS COVERED IN THE GROUP LEADER'S GUIDE

- (1) WHY IS THIS PROJECT IMPORTANT?
- (2) HOW IS THE PROJECT ORGANIZED?
- (3) WHAT ARE THE RESPONSIBILITIES OF A GROUP LEADER?
- (4) WHAT ARE THE DETAILS ABOUT THE MEETINGS THE GROUP LEADER MUST CONDUCT?
- (5) HOW SHOULD THE GROUP LEADER SELECT HER/HIS GROUP MEMBERS? WHAT IS THE "INFORMATION SHEET" THAT IS TO BE DISTRIBUTED? WHAT IS THE "PROJECT OVERVIEW"?
- (6) IN ADDITION TO SELECTING AND CONTACTING GROUP MEMBERS, HOW SHOULD THE GROUP LEADER PREPARE FOR THE FIRST MEETING? WHAT ARE SOME OF THE SKILLS THAT AN EFFECTIVE GROUP LEADER POSSESSES? WHAT IS A GOOD WAY TO ARRANGE THE MEETING ROOMS FOR THE COMMUNITY GROUP MEETINGS?
- (7) HOW DOES THE GROUP LEADER CONDUCT THE FIRST COMMUNITY MEETING?
- (8) WHAT SHOULD A GROUP LEADER DO BEFORE THE SECOND TRAINING SESSION?
- (9) WHAT MATERIALS SHOULD A GROUP-LEADER'S TEAM RECEIVE IN THE EQUIPMENT KIT?
- (10) HOW SHOULD THE GROUP LEADER PREPARE FOR THE SECOND COMMUNITY MEETING?
- (11) HOW DOES THE GROUP LEADER CONDUCT THE SECOND COMMUNITY MEETING? WHAT SHOULD BE REVIEWED FROM SESSION I? HOW DETAILED SHOULD THE SKILL STATEMENTS BE?
- (12) WHAT SHOULD THE GROUP LEADER DO AFTER THE TRAINING SESSIONS ARE COMPLETED?

TABLE IV

AT&T/BATTELLE APPROACH FOR ASSESSING COMMUNITY

EXPECTATIONS FOR PUBLIC SCHOOL CURRICULUM

1. Obtain the support of the Board of Education and a Coordinator for it.
2. Recruit Trainers and Group Leaders. Develop a plan for recruiting participants.
3. Identify and secure appropriate locations for meetings and schedule them.
4. Prepare materials explaining the goals of the project and the processes to be employed to generate information.
5. Prepare audio tapes to be used by Trainers in training Group Leaders.
6. Develop Group Leaders' Guide.
7. Instruct Trainers on use of audio tapes and Group Leaders' Guide. Discuss anticipated problems Group Leaders might encounter in their community meetings.
8. Prepare instruction for Group Leaders including what their responsibilities are, explanation of how the project is organized, how to recruit community participants, how to arrange location for community meeting, how to prepare for community meeting, what to distribute (posters, hand-outs, index cards, badges, magic markers, etc.), and the results needed.
9. Insure that Trainers have recruited, confirmed, and reminded Group Leaders and have prepared for and scheduled training sessions.
10. Prepare plan for analysis of meeting results and the scheduling of succeeding meetings.
11. Insure that Group Leaders have recruited representative participants from the community and arranged for meetings.
12. Send postcards to Group Leaders and participants to remind them before each meeting.
13. Insure that Group Leaders report results of meeting to Project Director.
14. Analyze results and prepare feedback to participants and request more refined input.
15. Repeat 12 and 13 and summarize the results.
16. Have participants review results and feed back.
17. Prepare and distribute final report.

The Battelle staff prepared materials for orienting the supporting staff and others. These materials explained the goals of the project and the processes to be employed for generating information. The Battelle staff also prepared audio tapes for the trainers to use in training group leaders and developed a group leader's guide. The topics covered in the guide are shown in Table III. The trainers were instructed on the use of audio tape equipment and the use of the tapes and the group leader's guide. Anticipated problems that the group leaders might encounter were discussed along with possible ways the trainers might suggest for handling them.

The results of the assessment were sent to the 1700 participants, the Board of Education, and the school administration.

A summary of this approach is given in Table IV. Application kits for this approach are available from Communication Technology Corporation in Marlton, New Jersey.

3. WASHINGTON STATE APPROACH

The approach described in this section was used in the State of Washington to help define alternative futures for the state and to formulate proposed legislation to help achieve specific goals. The effort was initiated by the Governor to guide the state government's program planning and to provide criteria for making budget decisions.

There were several categories of participants. First, there was a task force of 165 members. These were chosen from over 4000 nominees as a representative sample of the state's population. Second, there were 1500 participants from ten geographic areas. Third there were 500 from each of these ten geographic areas making a total of 5000. And finally there was a group of 1000.

The initial efforts consisted of four, 3-day workshops lasting 16 hours each

day. The participants in these were the 165 members of the task force. The workshops consisted of classroom sessions, in-depth exchanges of information and perceptions, and futures-creating exercises. The 165 produced a number of candidate goals and analyzed them for the impact each would have on the other goals.

These workshops were augmented by areawide meetings in ten geographic areas of the state involving 1500 participants. These meetings were briefer futures-creating sessions.

Following the ten area meetings, the 165 reconvened and finalized eleven alternative futures and a large number of specific goals.

Next an intensive media campaign was mounted to inform the population at large about the results. Weekly television programs were given in an effort to educate citizens. News releases were prepared and distributed to newspapers throughout the state.

Following the media campaign, citizens were surveyed for their preferences. The questions were developed by researchers working with the 165. A telephone survey of 1000 citizens was done to get feedback as quickly as possible. A mail survey was made of 500 citizens in each of the ten geographical areas to determine geographical variations in preferences. The 165 were surveyed by mail to obtain the views of those who had struggled with the issues. And, finally, the 1500 participants in the area conferences were surveyed by mail to get "informed" perspectives.

The results of the surveys were cross-tabulated to provide views of subgroups, e.g. high and low income groups. The results were presented to the 165 task force members who then prepared final recommendations which were strongly influenced

by the survey results. There were seven issue areas included: economic growth and population settlement; environmental protection and land use; natural resources and energy; transportation and communication; human development; education and training; and government.

The recommendations were incorporated into the governor's 1975 state of the state address and integrated into 20 of the 30 pieces of legislation he recommended to the legislature.

As may have been noted by the reader, this approach incorporates features of the two approaches described previously.

More detail can be obtained from references [5] and [6].

TOOLS FOR COLLECTIVE INQUIRY

In the first section of this volume, seven steps of a collective inquiry are described beginning with defining the purpose of the inquiry and ending with documenting the process and the results.

In the second section, three approaches to planning the arrangements for a collective inquiry are given. These approaches are especially useful for step two of the first section, i.e. planning the process for the inquiry.

In this section some tools are presented that are intended to be useful in carrying out the other steps. After the tools are described, a report is given on a field test of the tools.

DESCRIPTIONS OF TOOLS

In the pages that follow, brief descriptions of five recommended tools for conducting collective inquiries are given. These are brainwriting, nominal group technique, interpretive structural modeling, voting procedures, and worth assessment. (The options profile is described in Volume 2.)

The applicability of these tools and the approaches described earlier are indicated in Figure 2. The descriptions of the tools are intended to give the reader enough information to see their usefulness. On the other hand, those who want to apply the tools will probably need additional information. That can be obtained from the references provided at the end of this volume.

TOOLS AND APPROACHES

| STEPS | Shawnee Forest Charette | AT&T/BCL School Curricula | Goals for Washington State | Options Profile | Nominal Group Technique | Brain-writing | ISM Program | ISM Amendment | Worth Assessment | Voting Procedures |
|--------------------------------------------|-------------------------|---------------------------|----------------------------|-----------------|-------------------------|---------------|-------------|---------------|------------------|-------------------|
| Define the Purpose(s) | | | | | X | X | | | | |
| Plan the Inquiry Process | X | X | X | | | | | | | |
| Generate the Elements of the Issue | | | | | X | X | | | | |
| Choose a Contextual Relation | | | | | X | X | | | | |
| Determine the Relations Among the Elements | | | | | | | X | X | | X |
| Analyze, Amend, and Rank the Results | | | | X | | | | X | X | X |
| Document the Process and the Results | | | | X | | | | | | |

Figure 2. COLLECTIVE INQUIRY PROCEDURES*

*An X in a cell indicates that the Tool or Approach at the top is applicable to the Step in the row.
A blank cell does not mean that the Method is not applicable.

BRAINWRITING (IDEAWRITING)

A collective inquiry tool useful for generating ideas by small groups. Ideas are generated by individuals stimulated by a carefully prepared, trigger question. Ideas are written on sheets of paper. The sheets are passed to other participants who elaborate or clarify the ideas already recorded or add new ideas. This tool is helpful in defining problems, conceptualizing approaches to solutions, and formulating policies. It is especially useful when uncertainty or controversy exists about an issue or problem and its possible resolution. It works well even when it is important to neutralize the effects of dominant individuals.

APPROPRIATE CONDITIONS FOR USE

- o A need to collect ideas relevant to some issue.
- o A carefully prepared, focused trigger question.
- o Qualified persons willing to participate.
- o Qualified group leader willing to facilitate.

APPLICATION AREAS

- o Generally applicable when there is a need for collective idea generation. Especially useful for defining problems, requirements, or objectives to be used as inputs for other tools, e.g. ISM or worth assessment.
- o Useful for involving stakeholders in planning.

RESULTS OF APPLICATION

- o A list of 50 to 100 ideas about some issue.
- o Increased understanding by the participants about the issue, the ideas, and each other.
- o An opportunity for all to contribute ideas.

RESOURCES REQUIRED FOR APPLICATION

- o No more than six persons for a single group. Multiple groups may work simultaneously.
- o Each group needs a quiet place to work, a table and chairs, paper, and pencils.
- o From 15 minutes to two hours for process.
- o Funds, if required, for participants or leader.

HOW THE PROCESS WORKS

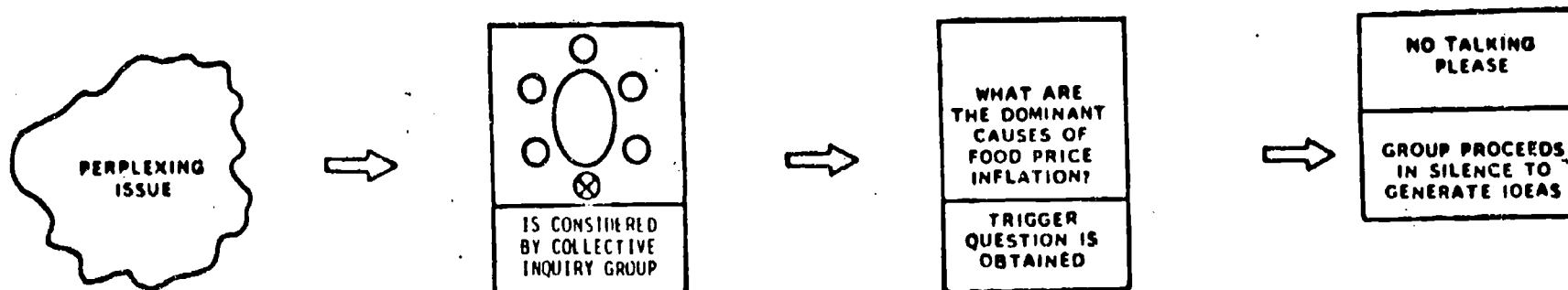
- o First, silent generation of written ideas by individual participants stimulated by carefully prepared, focused trigger question written at the top of a sheet of paper.
- o After about 5 to 10 minutes, participants exchange sheets of paper and build on the ideas already recorded or add new ideas that occur.
- o Process continues with further exchanges until all or most of the participants have seen all or most of the sheets.
- o Participants or facilitator edits the lists by clarifying and coalescing similar ideas.

IMPORTANT ATTRIBUTES OR FEATURES

- o Potential for generating many ideas concerning organizational or behavioral issues.
- o Potential for encouraging contributions from reticent participants or moderating dominant ones.
- o Opportunity for all stakeholders in issue to provide inputs to the process.

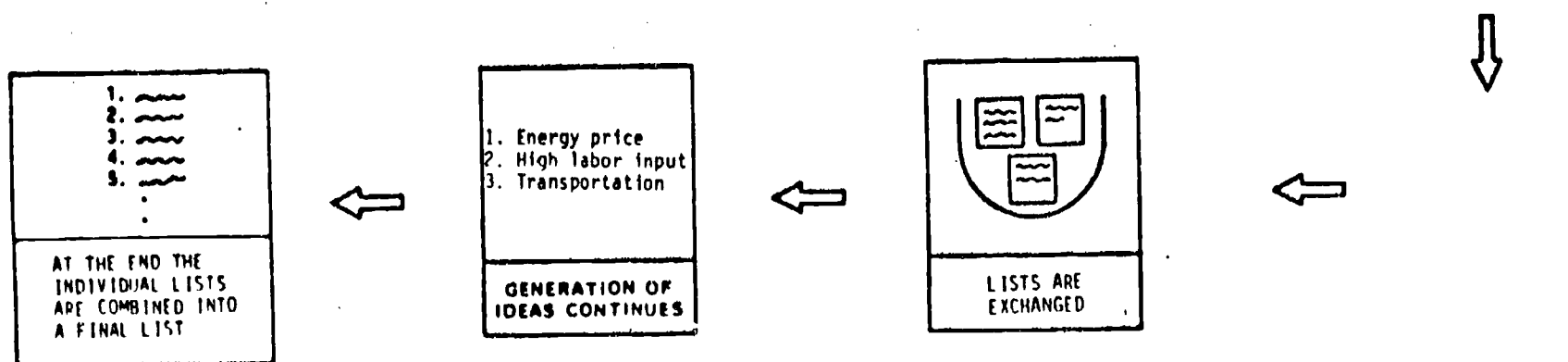
RELATED TOOLS

- o No other tools are required to use brainwriting.
- o Alternate tools are nominal group technique, Delphi, and content analysis.
- o Results may be used as input to ISM or worth assessment.



BRAINWRITING

Faced with the need to generate ideas related to an issue or problem, a facilitator is obtained, facilities are obtained, and a group knowledgeable about the issue is convened. A carefully prepared, focused trigger question is phrased by or explained to the group. Stimulated by the question, the members write their ideas on sheets of paper. After 5 to 10 minutes, the lists of ideas are exchanged among the members and each tries to elaborate or clarify the ideas on the sheet or adds new ideas that occur to him. The process continues until a set time is reached or the group has no new ideas or the sheets have been passed to everyone. Then the ideas are clarified and/or coalesced to produce a final list.



NOMINAL GROUP TECHNIQUE (NGT)

A collective inquiry tool useful for generating ideas by small groups. Ideas are generated by individuals, then discussed, clarified, elaborated, and combined by the group. This tool is helpful in defining problems, conceptualizing approaches to solutions, and formulating policies. It is especially useful when uncertainty or controversy exists about an issue or problem and its possible resolution. It works well even when it is important to neutralize the effects of dominant individuals to allow all participants to contribute ideas.

APPROPRIATE CONDITIONS FOR USE

- o A need to collect ideas relevant to some issue.
- o A carefully prepared, focused trigger question.
- o Qualified persons willing to participate.
- o Qualified group leader willing to act as a facilitator; preferably not an expert on the issue.

APPLICATION AREAS

- o Generally applicable when there is a need for collective idea generation. Especially useful for defining problems, requirements, or objectives to be used as inputs for other tools, e.g. ISM or worth assessment.
- o Useful for involving stakeholders in planning.

RESULTS OF APPLICATION

- o A list of 20 to 100 ideas about an issue.
- o A preliminary ranking of the ideas according to a chosen criterion.
- o An opportunity for all to contribute ideas.

RESOURCES REQUIRED FOR APPLICATION

- o 6 to 10 persons for a single group. Multiple groups may work simultaneously.
- o Each group needs a quiet place to work, a table and chairs, paper and pencils. Leader needs a flip chart and felt-tip pen plus a place to tape up sheets.
- o One to two hours time for process.
- o Funds, if required, for participants or leader.

HOW THE PROCESS WORKS

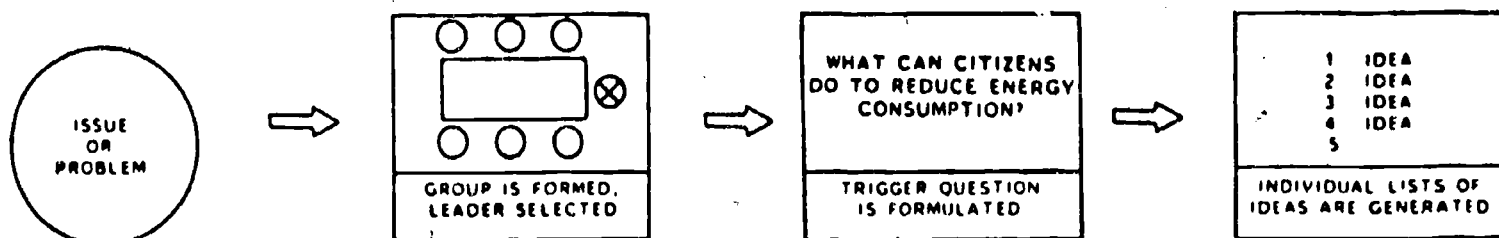
- o First, silent generation of written ideas by individual participants stimulated by an oral presentation of a carefully prepared, focused trigger question.
- o Individuals present ideas one at a time round robin.
- o Spontaneous elaboration or clarification of ideas is encouraged, but no discussion or criticism.
- o Ideas are recorded on flip chart to form group list.
- o Round robin discussion of the resulting list.
- o Voting on the ranking of the ideas generated.

IMPORTANT ATTRIBUTES AND FEATURES

- o Potential for generating many ideas concerning organizational or behavioral issues.
- o Potential for encouraging contributions from reticent participants and moderating dominant ones.
- o Opportunity for all stakeholders in issue to provide inputs to the process.

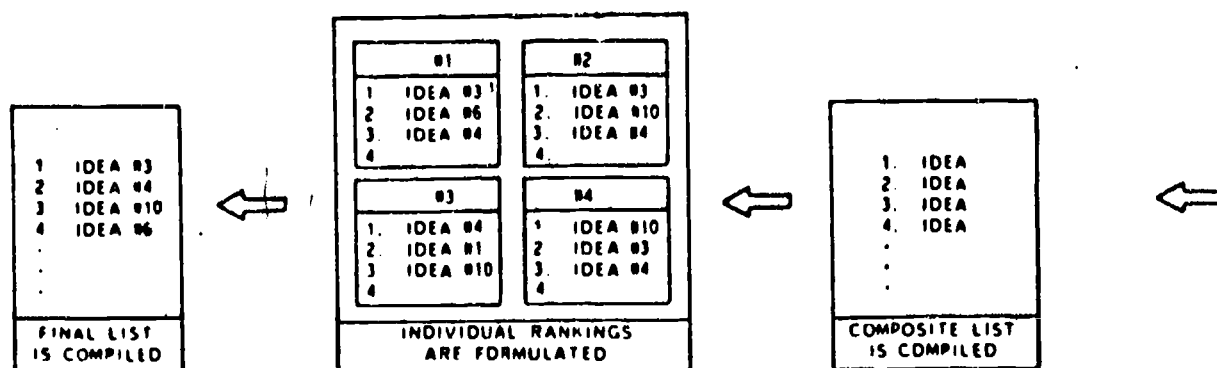
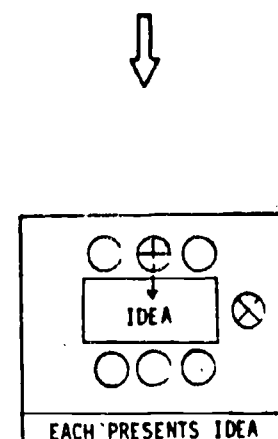
RELATED TOOLS

- o No other tools are required to use NGT.
- o Alternative tools are brainwriting, Delphi, and content analysis.
- o Any of several voting schemes may be used with NGT.
- o Results may be used as input to ISM or worth assessment.



NOMINAL GROUP TECHNIQUE

Faced with a need to generate ideas related to an issue or problem, a group leader is chosen, facilities are obtained, and a group familiar with the issue is convened. A carefully prepared, focused trigger question is explained to the group. Stimulated by the question, the members write their ideas on sheets of paper. Then, in turn, each presents one idea to the others for elaboration or clarification. After all the ideas are presented and recorded on a flip chart, they are ranked by each participant using some agreed-upon-criterion as usefulness, relevance, etc. Then these rankings are combined by some voting scheme to produce a final list.



A computer assisted learning process than enables an individual or a group to develop a structure showing inter-relations among a set of given elements.

APPROPRIATE CONDITIONS FOR USE

- o A set of elements related to the issue.
- o A relationship which is appropriate to inter-relate the elements.
- o A complex issue requiring understanding of interactions among many elements.
- o Persons who have knowledge of the issue and are willing to participate.

APPLICATION AREAS

- o Generally applicable where there is a need to relate elements of an issue, problem, system, etc. to obtain a holistic view.
- o Useful as a step toward the development of quantitative models
- o Useful for setting priorities or precedences.

RESULTS OF APPLICATION

- o A structured model of related elements.
- o A carefully refined language to describe or discuss an issue or system.
- o A shared and improved understanding and definition of elements and relations used.
- o Enhanced understanding of the whole.

RESOURCES REQUIRED FOR APPLICATION

- o A responsible client for the study who wants to benefit from knowledge of available participants concerning a particular issue.
- o Up to 8 willing and able participants, a group leader familiar with interpretive structural modeling, and a computer operator.
- o A time-shared computer with programs for structuring. A large screen display is helpful, but an exercise may be conducted using a teletype-writer-type terminal with auxiliary small screen television display.

RESOURCES REQUIRED (continued)

- o Funds for participants' and leader's time; equipment, approximately \$50 per hour; and telephone lines.
- o The time required for an exercise depends upon the number of elements in the model and their complexity. A ten-element exercise might take one to two hours.

HOW THE PROCESS WORKS

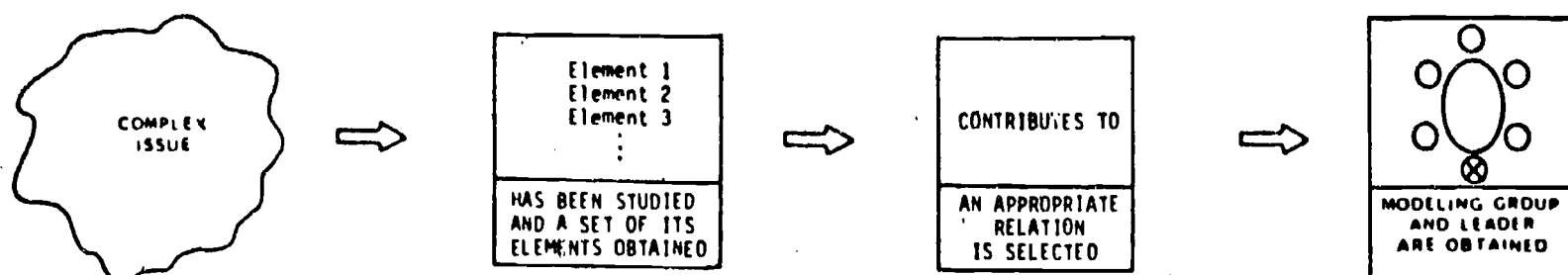
- o An issue and structuring theme are identified.
- o A group and a process leader are chosen.
- o Elements and contextual relation are obtained.
- o The group responds to computer-posed questions relating elements.
- o The computer displays the structure developed.
- o The group amends the structure until it is satisfactory.

IMPORTANT ATTRIBUTES AND FEATURES

- o Appropriate relation must be chosen carefully.
- o Elements and relation are clarified by reasoning and discussion stimulated by the process.
- o The quality of the results obtained is strongly dependent upon the leader. He should facilitate and not impose his knowledge of the issue.
- o Overemphasis of the mechanistic and technical aspects of the process should be avoided.

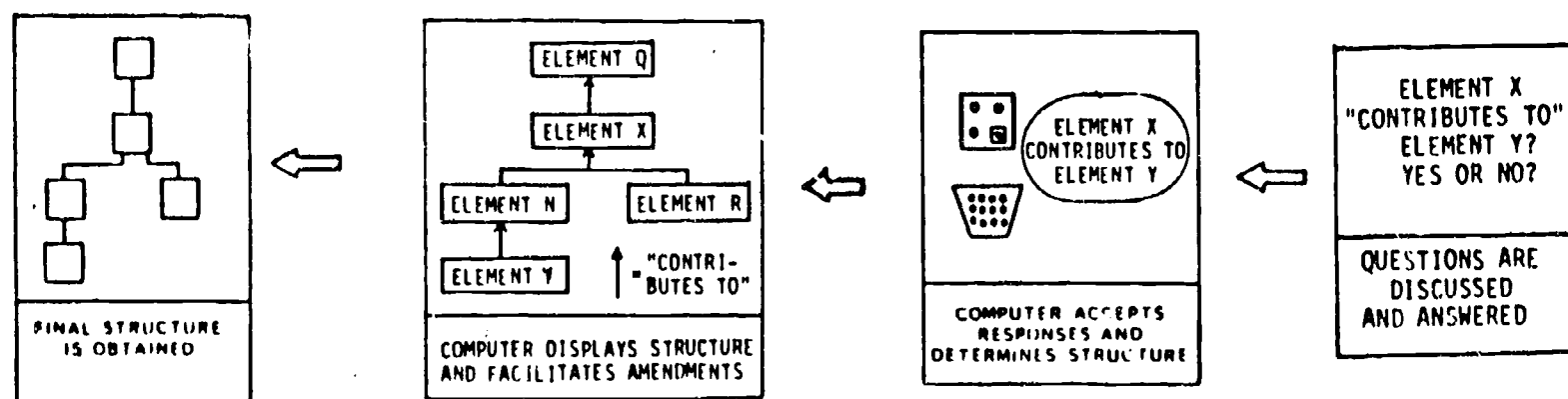
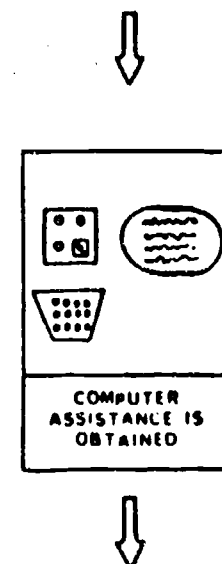
RELATED TOOLS

- o Nominal group technique, brainwriting, Delphi, literature search, or a combination of these can generate elements and contextual relations for the process.
- o In simple situations, "rearrange and tape" and other heuristic, non-computer-assisted methods may be used. Great care must be taken to avoid overly simple, erroneous results using these approaches.
- o The process may be used as a step to quantitative modeling or to structuring attributes or assessing worth for making choices.



INTERPRETIVE STRUCTURAL MODELING

This process systematically relates the elements of an issue or problem using a computer program with three kinds of inputs: 1) the elements; 2) the relation; and 3) participants' yes or no answers to questions posed by the program. Through the discussion preceding the answers, considerable learning is shared by the participants about the elements, the relation, and the structure of the model relating the elements. To decrease the number of answers required, the program uses logical inference. This requires that the relation used be transitive, i.e. if X relates to Y and Y relates to Z, then X will relate to Z. From the accumulated answers the program develops and displays a structural model of the relations among the elements and allows the modeling group to amend the model to their satisfaction.



Voting is a prevalent method of democratic group decision making. Any voting system can lead to unintended or paradoxical results where preferred candidates lose. In certain situations certain voting algorithms are better than others. It is helpful to be able to identify which voting systems are appropriate for given situations.

APPROPRIATE CONDITIONS FOR USE

- o Candidates and their impacts have been determined and a democratic method is needed to select a single alternative.
- o A group needs to make a decision concerning selection of an alternative.
- o Group consensus to select an alternative is not easily obtainable.
- o A legal mandate for voting exists.

APPLICATION AREAS

- o Generally applicable in all democratic situations when a group must choose among competing candidates.

RESULTS OF APPLICATION

- o A winner is determined as well as an indication concerning the strength of preference of the group for the winner.

RESOURCES REQUIRED

- o A set of candidates and a set of voters.
- o A mandate or desire to decide democratically.
- o A leader to explain the procedure and tally the votes.

HOW THE PROCESS WORKS

- o There are many voting systems. Some useful ones are:
 - Plurality: Voters vote for one candidate and the one receiving the most votes wins.
 - Majority Rule: Voters vote for one candidate. A candidate must receive more than 50% of the votes to win. If no candidate wins, runoff elections are held among candidates whose combined votes in the previous election constitute a plurality.
 - Weighted Voting: Voters are given a fixed number of votes. These are assigned by voters to competing candidates according to their strength of preference. The winner is the one

with the greatest point total. In one such system with N candidates the most preferred is given N-1 votes, the second most preferred N-2 votes, etc. Alternately, voters may assign their own preference votes on a fixed end point scale. This method of voting, also called Borda voting, is particularly vulnerable to "strategic" voting, in which people do not vote for candidates according to their preferences in order to insure victory for their preferred candidate.

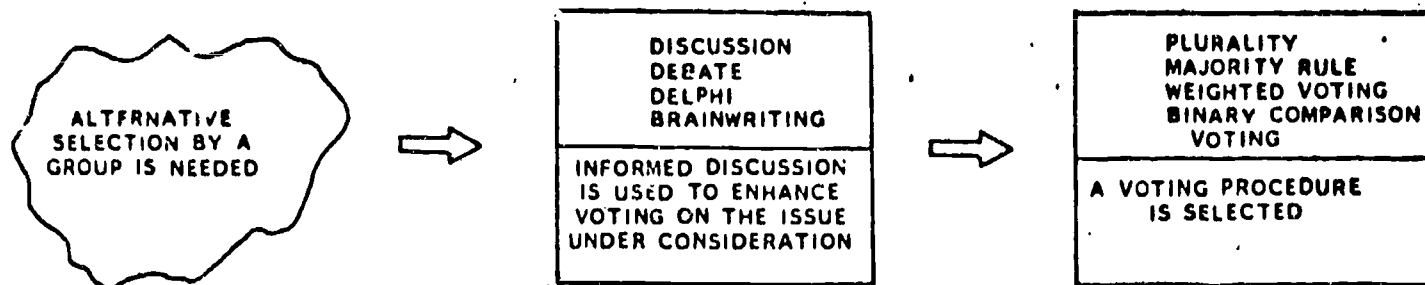
- Binary Comparison Voting: Voters note their preference ordering among all candidates or vote in a binary fashion for all possible paired alternative combinations. In one method, known as Condorcet voting, the candidate that beats all others in pairwise contests is the winner. In another method, known as Copeland voting, a candidate's score equals the number of candidates which it defeats by simple majority voting minus the number that defeat it. The candidate with the highest point total wins. A Condorcet winner will also be a Copeland winner. If there is no Condorcet winner, the Copeland method selects the candidate that has the greatest win-minus-loss score.

IMPORTANT ATTRIBUTES AND FEATURES

- o Potential for developing recommendations of preferred candidates.
- o Potential for voting paradoxes.

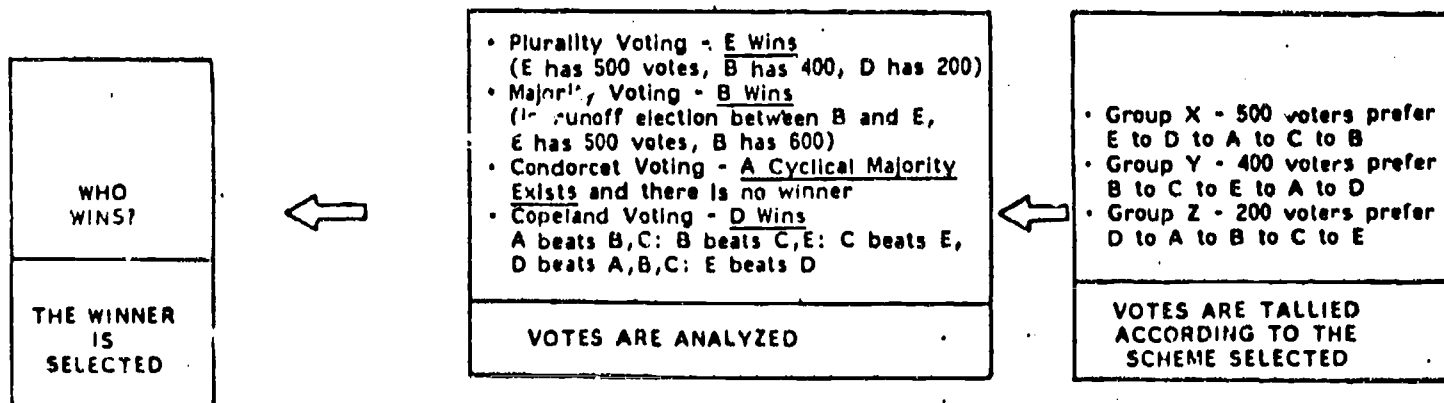
RELATED TOOLS

- o Various tools which require elements for inputs use voting. It is especially appropriate to precondition voters for informed voting.
- o Worth Assessment is an alternative approach.
- o Informed discussion which may be assisted by many of the collective inquiry tools presented in this volume may lead to informed voting and perhaps even consensus.



VOTING

A procedure in which a group expresses its preference for an alternative from among competing alternatives. Many voting procedures exist and the competing concerns involving trade-offs between simplicity and ability to detect voting paradoxes should be considered in selecting a voting procedure.



WORTH ASSESSMENT

Evaluation or selection of activities or alternatives is often based on subjective estimates of their worth. Worth assessment is a collective inquiry tool for translating qualitative impressions of value into understandable, consistent, and meaningful quantitative evaluations. This is accomplished by identifying and organizing the attributes of the activities or alternatives into a worth structure. The attributes are assigned weights according to their importance to the final result. Each alternative is then valued using the weighted worth structure to provide a quantitative basis for decision making.

APPROPRIATE CONDITIONS FOR USE

- o A need to decide among alternatives.
- o A need to evaluate alternatives quantitatively.
- o A need to predict individual's decisions.
- o A need to make individual or group values explicit.
- o A need to communicate the implications of alternatives.
- o A need to consider and organize multiple objectives and assessment criteria.

APPLICATION AREAS

- o Generally applicable for decision making by helping to specify, interpret, or value the impacts of alternatives on individuals or groups.

RESULTS OF APPLICATION

- o A tree structure showing lower level criteria (attributes or objectives) which compose higher level criteria.
- o An easily communicated display of the value or importance individuals or groups place upon criteria for proposed alternatives.
- o An explicit valuation of proposed alternatives according to the importance placed on their attributes.

RESOURCES REQUIRED FOR APPLICATION

- o A set of alternatives for a decision.
- o An individual or group that wants them evaluated.
- o Qualified persons to identify and weight criteria.
- o One to ten hours of participants' time.
- o A qualified leader.

HOW THE PROCESS WORKS

- o The individual or group whose subjective values are to be assessed is identified.
- o The alternatives to be evaluated are identified.
- o A set of performance criteria (attributes or objectives) are arranged in a tree structure such that lower criteria are parts of the criteria above them.
- o The criteria are assigned weights according to their relative importance (worth) to give a worth structure.
- o Each alternative is evaluated using the worth structure.
- o The values of the alternatives are compared for making decisions.

IMPORTANT ATTRIBUTES AND FEATURES

- o An explicit procedure that incorporates human judgment and values into a quantitative assessment.
- o Multiple worth connections among assessment criteria and their consequences on performance can be shown.
- o Risk considerations can only be approximated.
- o Sensitivity analysis can be done to show the effects of the weights assigned.

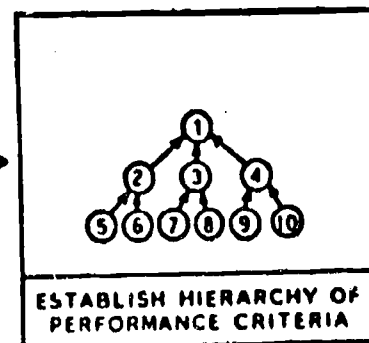
RELATED TOOLS

- o Brainwriting, NGT, or Delphi may be used to identify criteria (attributes or objectives).
- o Interpretive structural modeling is useful for developing the worth structure.
- o Decision analysis is an alternative tool, especially when probabilistic considerations are important.

COMPLEX DECISION
REQUIRED
(SELECT A POWER
PLANT SITE)

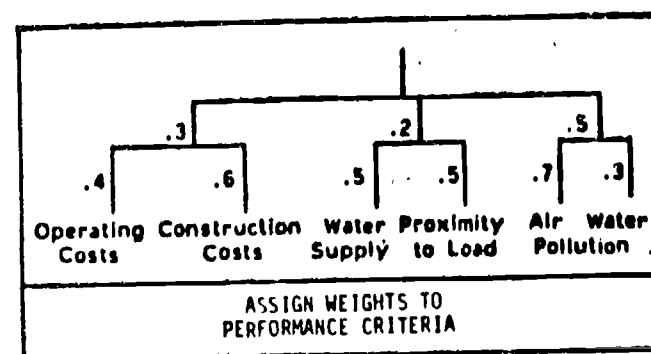


1. Attributes of New Power Plant
 2. Economic Factors
 3. Engineering Feasibility
 4. Environmental Conditions
 5. Operating Costs
 6. Construction Costs
 7. Water Supply
 8. Proximity to Load Centers
 9. Air Pollution
 10. Water Pollution
- IDENTIFY PERFORMANCE CRITERIA



WORTH ASSESSMENT

There is a need to choose among alternatives proposed to resolve an issue or problem. A set of performance criteria is established. The criteria are structured into a worth hierarchy and weights are assigned to the criteria according to their importance to the overall issue or problem. The worth of each alternative is valued using the weights and its particular attributes. The values for all alternatives are compared for making the final decision.



$W(x) = 0.677$
 $W(y) = 0.422$
 $W(z) = 0.352$
 x is "best" alternative

DETERMINE FINAL
WORTH SCORES AND
SELECT BEST
ALTERNATIVE



$W(x)$ = function of the
criteria weights and
the attributes of the
particular alternative

ESTABLISH WORTH SCORING
FUNCTION AND SCORE EACH
OF THE ALTERNATIVES



FIELD TEST OF TOOLS

After the tools described above had been chosen as candidates for assisting the conduct of collective inquiries, a field test was conducted. The purposes of the test were to determine if the descriptions of the tools were adequate and whether the tools were useful.

To accomplish the test, it was desirable to find panelists who would be interested in participating and who would be qualified to test the tools. The decision was made to ask the Tennessee Valley Authority to participate. It has a large environmental education program. It is responsible for regional development. And, it is very much interested in and involved in conducting collective inquiries. TVA consented and ten of its staff participated in the testing: 4 from environmental education; 2 from community development; 2 from the director's office; 1 from the citizen action office; and 1 from natural resources. All were involved with collective inquiries in one way or another.

The tools that were tested were: brainwriting; nominal group technique; worth assessment; voting procedures; and interpretive structural modeling. For each of these tools, the panelists were asked to judge the pictorial and written information describing the tool. These evaluations of the materials were used to make improvements in the descriptions.

For brainwriting, nominal group technique, and interpretive structural modeling, the participants made trial applications of each on real issues of importance to TVA. These applications were evaluated by the participants' responses (using a five point scale) to the following questions.

Q1. Was the method difficult to use? 1: very difficult; 5: easy.

Q2. Did the results achieved justify the resources used? 1: too expensive; 5: well worthwhile.

Q3. Did the results achieved justify the time of the participants?

1: wasteful; 5: well worthwhile.

Q4. As a participant, did you like the method? 1: disliked it;

5: liked it very much.

Q5. Was the application of the method effective? 1: not effective;

5: very effective.

For brainwriting, the averages of the evaluations for each question were as follows. Q1:3.6; Q2:3.7; Q3:3.9; Q4:3.9; Q5:3.4.

For the nominal group technique, the averages were: Q1:4.5; Q2:4.4; Q3:4.4; Q4:4.6; Q5:4.7.

For interpretive structural modeling, the averages were: Q1:4.3; Q2:3.7; Q3:3.8; Q4:4.0 Q5:4.1.

In addition to providing the numerical ratings, the participants also made a number of useful suggestions for improvements.

Although no statistical significance can be given to the results (lack of time, money, and appropriate panelists prevented such results), they were useful for a number of reasons. The suggestions made for improving the pictorial and written materials were very helpful. The panelists were well qualified, interested, and openly skeptical when they started. Their evaluations, especially of the applications, supported the judgments of the project team that these were useful tools for conducting collective inquiries. Moreover, applications of combinations of tools proved useful and indicated some improvements in ways to combine tools.

In summary, the field test supported the choice of tools and provided useful information for improving them.

In addition to the field test of the tools described above, another field test was conducted on the options field described in Volume 2. Briefly, this tool is designed to help in collective inquiries aimed at the design of environmental learning systems.

The participants in this field test were four TVA staff and twelve staff from colleges and universities that work with TVA in environmental education.

This test was conducted to obtain the judgements of experienced environmental educators on the options field as a design tool. The participants were given two questions at the beginning that they would be asked to answer at the end. These questions were:

1. What are the pros and cons of the design process (using the options field)?
2. What are the pros and cons of the dimensions and the options?

To answer these questions, the participants needed a "hands-on" experience using the design process. To that end they were asked to participate in a design exercise. The first part of the exercise was to develop a restriction structure. Detailed information is provided in Volume 2, but stated briefly this means the following. In the design of an environmental learning system, choices of options in one dimension, say, Presumed Learner Skills Base can restrict the choice of options in another dimension, say, Basic Learning Outcomes Sought. To develop the restriction structure, the participants used an interpretive structural modeling (ISM) approach. The elements were:

- A. PRESUMED LEARNER SKILLS BASE
- B. PRESUMED LEARNING STYLE
- C. SOURCE OF INFORMATION
- D. MODE OF ENVIRONMENTAL EDUCATION
- E. BASIC LEARNING OUTCOMES DESIRED
- F. TYPE OF ENVIRONMENTAL EDUCATION
- G. CURRICULUM DELIVERY CONCEPT
- H. MEDIATOR MODEL
- I. LEARNER INTERACTION RESOURCES
- J. ORIGIN OF FINANCING

The relation used was: "A choice of option(s) in Dimension X can restrict the choice of options in Dimension Y".

Through their participation in this exercise to design the restriction structure, the participants gained familiarity with the dimensions and the options within the dimensions.

After the restriction structure was developed, the participants participated in a second ISM exercise to develop a precedence structure. This structure ordered the dimensions of the options field for selection of options. The elements of this exercise were dimensions, and the relation used was: "A choice of option(s) in Dimension X should take priority over a choice in Dimension Y".

After the precedence structure was completed, the participants began (but did not finish because of time) to design a regional environmental learning system. The purpose was to involve the participants in the use of the

design process so that they would be able to evaluate it.

The first question listed above was addressed first. To answer it a nominal group technique was used with the following questions: "What are the positive features of the design process?" and "What are the limitations of the design process?" Items given in response to the first question (with no priority implied) were:

- o Flexible; adaptable to region; can be modified
- o Comprehensive
- o Gives direction, structure for eliciting decisions; systematic
- o Offers opportunity for different views and disciplines to be included in the design
- o Leads to consensus decisions which should lead to lasting effects. Includes realities of compromise
- o Engenders thinking and analysis
- o Involves individuals in the process
- o Illustrates design and dimensions of the design - enables review by the group
- o Permits computer-aided design.
- o Forces distinguishing between synthesis and analysis
- o Useful for capturing random ideas into useable form
- o Builds on experience and, hopefully, encourages ingenuity

Items given in response to the question on limitations (with no priority implied) were:

- o Limited population with ability to use process, limited ability to think through process
- o Restriction concept needs to be presented more clearly
- o Requires individuals who are interested and willing to give their time
- o Time to use right limits its use. Needs to be refined
- o Parameters (scenario) needs to be defined at start
- o Understanding dimensions' and options' denotations and connotations is very hard
- o Paramount to use well-organized facilitator training procedure
- o Lack of understanding of systems concepts affects people's feelings about process. Creates hostility, opting out of process.
- o Needs to be made clear that process is iterative
- o Outcome of process is affected by group interaction skills of participants
- o Facilitator needs to lay out work, pace work, and take stock of progress periodically

Time did not permit the second question regarding dimensions and options to be addressed in the same way as the first. Instead, the participants individually submitted lists of pros and cons regarding the dimensions and the options including some suggestions for additional options. Many of these have been incorporated in the material presented in Volume 2.

In summary, the options field is judged to be a potentially useful tool in the design of environmental learning systems but must be used with qualified participants and facilitators, and the context should be clearly defined.

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APPENDIX

COMPUTER IMPLEMENTATION OF ISM

PART ONE

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PART TWO

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Appendix to Volume 4:
Conducting Collective Inquiry
COMPUTER IMPLEMENTATION OF
INTERPRETIVE STRUCTURAL MODELING
(Part One)

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COMPUTER IMPLEMENTATION OF INTERPRETIVE STRUCTURAL MODELING

Introduction

One of the chief obstacles to solving the environmental problems that confront the world today is the inability of both average citizens and policy decision-makers to deal with the complexity of these problems. Environmental problems, such as air and water pollution, are comprised of many factors and involve interrelationships that may be difficult to understand. As a result, solutions to environmental problems have often been too simplistic, and have not always improved the quality of the environment.

This appendix presents information about interpretive structural modeling, a method that can be useful in dealing with complex environmental issues in a collective inquiry setting. The first section discusses interpretive structural modeling as it can be used for collective inquiry and presents an overview of an interpretive structural modeling exercise. Later sections of the appendix discuss the computer equipment and the interpretive structural modeling (ISM) software package.

What is Interpretive Structural Modeling?

Interpretive structural modeling (ISM) is a computer-aided method to assist a group of people in studying and analyzing complex problems. It is appropriate to use ISM when the issue or problem under study can be broken down into the component parts that describe the situation. Participants in an ISM exercise define the structure of a complex system by focusing on the relationships between the elements of the system.

Some ISM applications include:

- a study of children with learning disabilities;
- work with neighborhood groups to identify factors involved in neighborhood crime;
- an analysis of the Goals for Dallas;
- a long-range planning study of the Sahel region of Africa;
- the identification of goals for a state-planning effort;
- defining the objectives of PLANALSUCAR, an agency involved in the Brazilian National Alcohol Fuels Program;
- establishing priorities for teacher in-service education programs; and
- a study of environmental issues by high school biology and social studies students.

Sometimes the model is developed to reflect the group's existing knowledge; at other times the model is structured after additional study about the issue. The results of an ISM exercise are:

- a greater understanding of the complex issue or problem through focused debate and clearly defined terms;
- an easily understood model or diagram that shows the structure of the issue; and
- some degree of order to the problem so that solutions can be more readily identified.

Therefore, the ISM method can be a useful tool for environmental education.

Encouraging Participation in Collective Inquiry

Many environmental educators are concerned with the question of how to encourage broad participation in collective inquiry. Increased citizen involvement in urban issues serves as a good example. Collective inquiry is more likely to lead to effective issue resolution when it:

1. takes into account the complexity of the problem or issue;
2. results from a wide base of influence and supporting resources; and
3. shows an awareness of the high interdependency that characterizes urban society.

When confronted with a complex environmental issue, people may be highly uncertain and ignorant about the dynamics involved. There is a need, therefore, for a learning process that can help people (citizens as well as political and organizational leaders) to create a shared understanding of the issue.

In Volume 3, Creating a Regional Environmental Learning System, we presented a discussion of the collective inquiry and action processes -- dialogue, decision, action, and evaluation -- and explained how collective inquiry can be used to address and resolve regional environmental issues. Interpretive structural modeling can contribute to more effective collective inquiry, by encouraging participation in the discussions about the environmental issue.

Some of the features of ISM that have appealed to collective inquiry groups make it an effective method for environmental education. These features include:

- provides a clearer understanding of the issue under study;
- makes maximum use of the group's previous research;
- organizes their collective knowledge;
- leads to consensus; and
- provides a way to communicate the results of the group's work.

An Overview of an Interpretive Structural Modeling Exercise

To do an ISM exercise, a group first decides on a meaningful set of elements and an appropriate relationship to study. The next step is to use the chosen relationship to discuss how the elements are related to one another. This discussion should lead the group to a new understanding of the total system of elements and relationships. The graphical representation, or structural model, which is produced with the aid of a computer, serves as the basis for developing this understanding.

During an ISM exercise the computer functions as a bookkeeper, displaying the questions for the group to discuss, maintaining logical consistency, and recording the

responses. The computer is programmed to do this. This appendix includes detailed information about the computer programs for ISM.

During a typical ISM session the participants are seated at tables, with TV monitors located so that everyone can see one. A telephone receiver connects a small portable terminal with the computer. The TV monitors are wired to the terminal to provide a visual display for the participants. It is desirable to have a printer available to provide a typewritten copy of the session. A blackboard or flipchart is useful, and a special ISM Notetaking Form (Figure 1) should be provided for the participants.

An interpretive structural modeling exercise has six steps:

1. Orienting the exercise
2. Generating the elements and choosing the relationship
3. Structuring the elements with the relationship
4. Amending the model
5. Exploring the implications of the model
6. Documenting, evaluating, and communicating

The main purpose of the orientation is to establish clear expectations for the ISM exercise. This step includes such tasks as choosing an environmental topic to model, setting the schedule for the sessions, acquainting the participants with the ISM process, and clarifying the roles during the ISM exercise. The topic can be divided into specific areas for the participants to research individually or in task forces. Resource speakers can provide background information through lectures and additional resource materials.

In the second step of the ISM exercise, the group generates a list of elements that contribute to the environmental problem or issue and chooses the relationship that will be used for the modeling. During the discussions to produce this element list the participants can share their research on the environmental topic and develop an initial definition for each element.

7.3

NOTETAKING OF ISM SESSION

GROUP _____

PAGE _____ OF _____ PAGES

TOPIC _____

DATE OF SESSION _____

| Relationship | Y/N | Notes on discussion, definitions, voting, and other comments | Relationship | Y/N | Notes on discussion, definitions, voting, and other comments |
|--------------|-----|--------------------------------------------------------------|--------------|-----|--------------------------------------------------------------|
| | | | | | |
| | | | | | |
| | | | | | |
| | | | | | |
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The third step of the ISM process involves the use of the computer to structure the elements with the relationship. For the structuring sessions, you will need access to the computer by way of a terminal, telephone connection, and television monitors. You will also need someone to set up the equipment and manage the terminal during the structuring sessions. In addition to the physical setup, you must also arrange to use the ISM computer software package. Detailed information about computer equipment and the ISMS-UD software package is presented in later sections of this appendix.

With the terminal, telephone connection, and television monitors in place, you are ready to begin structuring. During the structuring sessions, the computer prompts the group to systematically examine the relationships between the elements, by displaying questions on the monitors. As each question appears on the TV monitor, the participants discuss their views and arrive at a response. A yes or no vote is recorded at the end of the discussion for each question. The leader's role during structuring is to guide the discussion and call for a vote to answer the questions, and the terminal manager's role is to operate the computer equipment. At the end of the structuring, the computer prints out the information needed to draw the model that the group has produced.

Experience has shown that it always takes longer to model than a beginner might expect. The time required to produce a model is a function of the number of questions that the group must answer, and the number of questions is a function of the number of elements to be modeled. As the number of elements increases, the number of questions to be answered rises exponentially. The computer is programmed to infer some of the relationships, thereby reducing the number of questions to be answered. In addition, when a new element is introduced the computer calculates the best sequence of questions to ask so that the group can examine all the interconnections in the most efficient way.

Notetaking is also an important part of the modeling sessions. Notes on each relationship can be helpful in later review of the model, and may point to areas that

require further study. A special form (Figure 1) for notetaking during an ISM session is available.

At certain points, you may stop the sequence of questions and use the computer to display the results of the modeling, from which you may draw a diagram of the model. The computer does not actually draw the model, but it prints out the information you need to make the drawing. The group can examine the drawing and review its understanding of the issue.

Sometimes the model includes a cycle among several elements. Within a cycle, every element is related to every other element. When a model includes a cycle, the group can give structure to that part of the model by discussing the strength of relationships between the cycle elements. By weighting the elements, it is possible to identify the strongest relationships within the cycle, thus further clarifying the model's structure.

In interpretive structural modeling, the final model emerges slowly, as much as a result of amending as of the initial structuring process. Often amending will be done both during modeling sessions and through outside staff work. Periodically, the group should examine the structural model for internal consistency and logic, and then amend as necessary. The first model may not be satisfactory -- it may contain relationships that do not exist, or perhaps an absence of relationships that should be there. The ISM software includes the capability to make any refinements in the model that the group feels necessary.

Sometimes the easiest way to amend the model is to remove one of the elements from the model and rework it through the usual series of questions. In other words, prompted by the computer, the group reworks the element back into the model. The amending process calls for a word of warning: the leader should be careful that the model the group has spent several hours to construct is not destroyed by haphazard amending.

To achieve a useful model the group should take time to explore the implications of the model. This might include, for example, discussing the assumptions that are implicit in the model. Another way to approach the model would be to identify key points (for example, points where the model branches). If these points can be considered strategic places to intervene, the group could spend some time discussing what persons or institutions might be most effective in resolving the issue under study.

Consider, too, how the model can be helpful in evaluating possible outcomes of alternative actions. Does the model suggest anything about what is right or wrong with our present approaches to the environmental issue or problem?

To summarize these possible discussion points, we offer this list of suggested questions:

1. How can the model be interpreted?
2. What was learned?
3. What can be said about the assumptions behind the model?
4. Does the model suggest any course of action that could help resolve the environmental issue?
5. Who should be involved in the solution?
6. Other ideas?

The final step in an ISM exercise is to document the process and the model for later reference, although the extent to which a model is formally documented will depend on how you intend to use it. The documentation might include a summary of how the model was developed, comments from the discussion, the element list and definitions, conclusions drawn from the model, and so on. Usually the documentation of the ISM exercise is only adequate as a summary for the participants. The group must use other methods to communicate the results to others who did not participate in the ISM sessions.

Since ISM exercises vary considerably (leaders, participants, issues, time spent), each exercise should be evaluated in terms of the purposes and goals of that particular

exercise. Avoid comparisons with other ISM exercises, using comparison only to assist planners and leaders in improving their skills with the ISM process.

It is clear that the steps in an interpretive structural modeling exercise closely correspond with the processes of collective inquiry and action discussed in Volume 3. Field testing, described in Volume 4, Conducting Collective Inquiry, and previous applications of ISM have shown it to be an effective method for collective inquiry.

Subsequent sections of this appendix present information about the computer equipment and the ISM software package needed to use interpretive structural modeling for collective inquiry and action.

Computer Equipment for Interpretive Structural Modeling

This section provides a description of computer equipment that is suitable for implementing ISM. First, there is a discussion of the modular approach to the computer equipment. The modular approach allows freedom to select equipment that meets budget and computer capability requirements. Next, we discuss alternative equipment for each module. Finally, the alternative equipment is combined to provide some alternative system configurations. Most of these system configurations have already been used and proven in during interpretive structural modeling exercises.

The equipment and the systems recommended in this section are representative of the current technology available. The reader should be aware of the fact that prices fluctuate with the state of the art in computers and computer peripherals. When it comes time to acquire equipment, be sure to check various models and different manufacturers' equipment to determine how to get the most computer capability for your dollar.

Modular Approach to Equipment

In this section, we discuss a modular approach to configuring computer equipment for ISM. A modular approach facilitates the design of a system to maximize capability and minimize cost. We use this approach because of the varying budgets of possible

purchasers. The computer equipment for ISM has four distinct subsystems. These subsystems are: the computer system, display units, graphics units, and voting systems. The first two subsystems are required; the last two are optional.

The computer system is the heart of the equipment. It consists of the central processing unit, memory, and peripherals. The computer is the device that executes software and drives the display and graphics units. We will describe the alternative computer systems later in this section.

The display units (more commonly called terminals) provide the interface between the users and the computer equipment. A display unit looks similar to a typewriter keyboard, but it usually has a CRT (cathode ray tube), i.e., TV-type screen or monitor, and/or a printer mechanism. The CRT-type units usually have a plug available so that a limited number of extra monitors can be attached, which increases the number of people able to see the computer's output. You will need at least one display unit.

A graphic unit is also a display unit, but it is considered separately because of the additional functions it can perform. The graphic units (graphics terminals) look almost identical to the display units; both have keyboards and CRTs. However, the graphic unit adds the capability of color, bar graphs, and other features that a display unit does not have. Since it costs ten to fifteen times more than a display unit, a graphic unit is an optional device for ISM.

The fourth subsystem is the voting unit. A voting unit is electronic equipment that can accept vote "values" (weighted voting) and display the results. These units allow users to vote secretly (instead of by a hand vote) and provide many statistics related to the vote. This device allows the rapid and accurate tally of votes when large numbers of people are involved. Due to its high cost, the voting unit is also optional.

Module Descriptions

In this section we describe alternative modules for each of the four subsystems. To facilitate the purchase and maintenance of equipment, this selection is limited to

off-the-shelf units. There are many equipment manufacturers; the selection of the following equipment is based upon manufacturer's reputation and warranties, as well as our experiences using the equipment.

Computer Systems

One computer system that could be used for interpretive structural modeling is that of a university, local corporation, or timesharing computer company. This would involve contracting with the computer installation an access agreement that allows the use of the computer resources. There are some major economic advantages to this arrangement. First, the purchase and maintenance costs of a computer system are completely eliminated. Second, since the computer system will be accessed via the community telephone system, the resources of this large computer system could be available to every phone subscriber. Some disadvantages are that the use of this system could be restricted to certain times of the day and certain days of the week. Also, computer malfunctions can cause disruptions of computer sessions. The cost of the using the system can also be limiting. Ideally, timesharing costs should be based on a cost per access. However, some computer services also include a minimum monthly charge that could exceed the amount of time used. Since cost of a timesharing system such as this varies, no price estimate can be given.

The second computer system is the TRS-80, which falls in the \$1,000-\$4,000 price category, as of mid-1979. It is a computer system manufactured by Tandy Electronics and marketed by Radio Shack. The TRS-80 is easily maintained, since there are over 6,000 Radio Shack stores throughout the United States and in nine foreign countries. Expanding and upgrading the system are easy to do with the peripherals manufactured by Tandy and numerous small companies. A disadvantage of this system is that the expansion is limited to a single user. The execution of very large programs is limited to the amount of memory in the system.

The Cromenco System Three Computer has been selected as the third computer system. In mid-1979, its price is approximately \$10,000 to \$20,000. The System Three Computer is a table-top system which will provide fast execution of RELS tasks, as well as other tasks such as word processing, inventory, and payroll. The system is easily expandable, utilizing the industry standard S-100 bus architecture, which makes available low cost memory and peripheral interfaces. Up to eight simultaneous users are able to use the system, and no special environmentally controlled rooms are required. The one major disadvantage with the system is that it will require a maintenance contract from an outside company.

The last computer system to be discussed is the Hewlett-Packard 3000 Series II computer system. This is a full scale computer system costing over \$140,000. It is capable of handling the processing requirements of local governments and school systems. The computer can be easily expanded (although not at a low cost) to provide support for up to 53 simultaneous users. Many upgrade and expansion options are available; they permit the system to be tailored to specific needs. The large initial investment and maintenance contract prices, however, might be prohibitive. A system of this size requires a staff of personnel to operate it and to provide programming help. It should be noted that a computer system of this size is likely to be more than is needed for collective inquiry in a RELS.

Display Units

The first display unit is the Digi-Log Model Telecomputer II CRT terminal. This terminal features a built-in acoustic telephone coupler, a slave port for a printer, 40 and 80 characters per line, and a video output jack for connecting compatible video equipment, such as extra monitors and video recorders. This device costs about \$1,500, and requires a closed circuit television (CCTV) monitor (about \$150) for operation.

The second display unit is the Teleray 1000 series intelligent CRT terminal. It features user programmable functions; software selectable character size, selectable

fields including dim, inverse video, blinking, and underline; and a slave port for a printer.

requires an external acoustic coupler (about \$200), if a telephone interface is to be used. Currently, the Teleray 1000 costs about \$900.

Both types of display units are capable of driving additional video monitors. The video monitors are of the type used in a CCTV system with a composite video input. It should be noted that a large screen television can also be used as a video monitor. Any of the above model (or equivalent) large screen displays are compatible with the above units. A large screen display costs over \$4,000. However, many metropolitan areas, there are audio/video specialty shops that will rent a large screen display for about \$150 a day.

The third display unit, the Texas Instruments Model 745 Silent 700, which is a printer type device. This unit weighs 15 pounds and has its own built-in acoustic telephone coupler. It is a complete unit in itself. An optional \$50 interface allows this unit to double as an RS-232-C printer for the above two display units. With a price of \$1,500, this dual-purpose unit is an excellent choice.

The fourth display unit, the Texas Instruments Model 800 terminal is also a printer type device. This terminal features a printer speed of 120 characters per second and an upper and lower case character set. It can also be used as a hardcopy printer for an RS-232-C compatible device. Its price is approximately \$2,300.

Graphics Units

The recommended graphics units are the Tektronix 4027 and the Ramtek Micrographics color graphics terminals. These terminals feature color displays, slave printer and plotter ports, user definable fonts, and the capability of driving additional CCTV monitors. Devices such as the ones named above cost over \$10,000.

Voting Devices

At this writing, only one commercially manufactured voting device, called Consensor, is available. This unit will support up to 16 voters. It visually displays statistics such as the distribution, average weighting, and the mean of the vote

distribution. Its display device displays the information either numerically or in a histogram format. The price of this unit is about \$8,200 -- plus installation.

Alternative Configurations

In this section we suggest some alternative configurations of the above mentioned equipment. Although not all configurations will be optimum for every situation, they should serve to provide a direction for the eventual equipment selection. Not included in these configurations are the cables and the extension cords required. These items are of a variable cost and can be supplied by the equipment manufacturers.

Configuration 1 is designed around the use of a timesharing computer system. The equipment recommended is the Digi-Log CRT terminal, Texas Instruments Model 745 display unit, and six CCTV monitors. This equipment configuration allows computer sessions with up to twelve participants. By adding a large screen display, approximately fifteen to thirty participants can be accommodated. The approximate total cost is \$4,500, excluding the large screen display.

Configuration 2 is also designed around the use of a timesharing computer. It consists of equipment in configuration 1, with the addition of the Tektronix 4027 color graphics terminal. This addition costs another \$10,000 and gives color display capability.

Configuration 3 is the Radio Shack TRS-80 microcomputer system, which is recommended for single use situations. This system provides sufficient computing power for the ISM software.

Configuration 4 consists of the Cromemco System Three computer, and any combination of display and graphics units. This configuration will provide support for the ISM software as well as for word processing, inventory, and payroll with up to eight simultaneous users. The cost of the basic system is about \$10,000 for one user; additional user capability costs approximately \$2,000 each.

Configuration 5 is a full-sized computer system that can support up to 54 users and provide such services as on-line record keeping, payroll, and others. It consists of the

Hewlett-Packard 3000 series minicomputer and any combination of display and graphics units. Such a system would cost over \$160,000.

Manufacturers

Cromemco, Inc.
280 Bernado Avenue
Mountain View, CA 94040

Teleray
Division of Research, Inc.
Box 24064
Minneapolis, MN

Hewlett-Packard Co.
Computer Systems Groups
100 Wolfe Road
Cupertino, CA 95014

Radio Shack
Division of Tandy Corporation
2617 W. 7th Street
Fort Worth, TX 76107

DIGI-LOG Systems, Inc.
Babylon Road
Horsham, PA 19044

Texas Instruments
Digital Systems Division
P.O. Box 1444
Houston, TX 55424

Consensor
Applied Futures, Inc.
22 Greenwich Plaza
Greenwich, CT 06830

Strategies for Obtaining Equipment

This section presents strategies for obtaining the computer equipment needed for interpretive structural modeling. We discuss some strategies for acquiring computer equipment and recruiting personnel capable of interconnecting and operating the equipment. The reader can apply these and other appropriate strategies to configure and support a computer system suitable for interpretive structural modeling.

Utilizing Community Resources

One community resource to tap is the computer professional -- someone whose occupation is to program, design, and procure computer systems. Computer professionals may be vital to the successful implementation of ISM. Ideally, you will form a team of computer professionals with varying expertise. Systems analysts, computer engineers, and programmers all have skills to contribute. A systems analyst is concerned with the

computer as a total system to serve a useful function, such as banking, payroll, and record systems. Since this total system incorporates both the hardware (equipment) and the software (programs), a systems analyst has broad knowledge that should be valuable. The computer engineer has a college degree in either electrical or computer engineering. The engineer's job is to design the computer equipment or to configure computer equipment to perform tasks related to production, energy management, etc. Although the computer engineer is basically hardware oriented, his or her expertise is a valuable asset, especially when utilizing borrowed and donated computer equipment. The programmer's job is to program computers to perform a variety of tasks. The programmer can help maintain the software and install it on your computer system. When assembling a team of computer professionals, try to recruit at least one of each type. These people work for such institutions as banks, schools, universities, and businesses.

The second community resource to investigate is the wealth of computer equipment. Possibly, computer equipment can be begged, borrowed, or bought from universities, local companies, and local, state, and Federal agencies. A moderately large local university probably has a timesharing computer system and a meeting room equipped with the required electronic devices. Local companies might donate their old computer equipment. The donation of computer equipment by state and Federal agencies is another way to obtain equipment. A RELS is likely to qualify for a discount on the purchase of new and used computer equipment from manufacturers and local computer stores.

Probably the best way to locate sources for the use, donation, and purchase of computer equipment is to rely on the computer experts. They will know of local computer installations, equipment in their company (or another) that is slated for replacement, and local computer vendors.

Development of a Support Plan

Because of the large investment of time and money, some support plan for the computer system is necessary. Computer systems require support for both the system

hardware and the software. The support or maintenance of the hardware is taken care of during the warranty period; after that period, support is provided through a maintenance service contract or by a computer engineer or technician. The system software requires maintenance whenever errors in the programs are found. No matter how many times a computer program has been used, some latent programming errors may still exist -- especially within a large, complex program such as the interpretive structural modeling software package. The computer programmer, or another computer professional on your team, can correct these errors.

In summary, the details of the support plan depend on the specific equipment configurations you are using. In general, we recommend you develop a support plan early, allowing for as many contingencies as possible, so you can maximize the use of your computer equipment and software packages.

Programmer's Overview of the Interpretive Structural Modeling Software Package

The purpose of this programmer's overview is to document the Interpretive Structural Modeling Software (ISMS) package developed by the University of Dayton. ISMS-UD version 2 is a software package consisting of three interactive FORTRAN IV programs that aid in conducting ISM sessions. We document these programs in such a way that a computer professional can easily install them on other existing computer systems. Earlier versions of ISMS-UD have been successfully installed on various large and small computer systems, from small PDP-11's to CDC Cyber's. Included at the end of this section is a partial list of facilities that have installed ISMS-UD.

This section proceeds as follows. First, we offer a general description of ISMS-UD to acquaint the reader with the system scenario. Second, we discuss the installation of ISMS-UD from a user's point of view. Then we cover the installation details, including file formats and core storage requirements. Four attachments supplement this

documentation. Attachment 1 is the detailed documentation of the machine algorithms that comprise ISMS-UD. Attachment 2 is the interpretive structural modeling software. It is a user's manual for ISMS-UD, as it has been installed at the University of Dayton. DELTA charts for ISMS-UD are included as Attachment 3. Attachment 4, which is bound separately from this part of the appendix, consists of the FORTRAN programs for ISMS-UD.

ISMS -- The System

Interpretive structural modeling is a method that helps people think and communicate more effectively about complex issues and problems. ISMS is a series of algorithms that have been coded to allow machine execution of ISM. Warfield (1976) has presented previous discussions of these algorithms. Warfield's book will serve as an excellent reference for the documentation of the ISM algorithms.

ISMS-UD consists of three interactive FORTRAN IV programs -- ISMS-UD, CYCLE, and MAKEIT. The ISMS-UD program includes the rudiments of ISM; that is, embedding (Transitive Bordering Method) and some sophisticated digraph amending algorithms. CYCLE permits the resolution of cycles that may be included in an ISM digraph output of the ISMS-UD program. Resolution thresholds, geodetic cycles, and their associated paths are the primary outputs of CYCLE. The third program in the package, MAKEIT, allows both ISMS-UD and CYCLE to present queries to the user in an English text format. MAKEIT accepts a sequentially oriented file, usually created by the host computer text editor, and reformats it into another file that is randomly accessed by ISMS-UD and CYCLE. Figure 2 is a block diagram of the system.

ISMS-UD was designed and developed on the University of Dayton Univac Series 70/7 timesharing computer system under the VS/9 version 3.5 operating system and FORTRAN IV compiler BGFOR. The Univac system is an updated RCA Spectra, which is a hardware duplicate of an IBM 360 but with virtual storage facilities. Although virtual storage is not required for the installation of ISMS-UD, it is desirable.

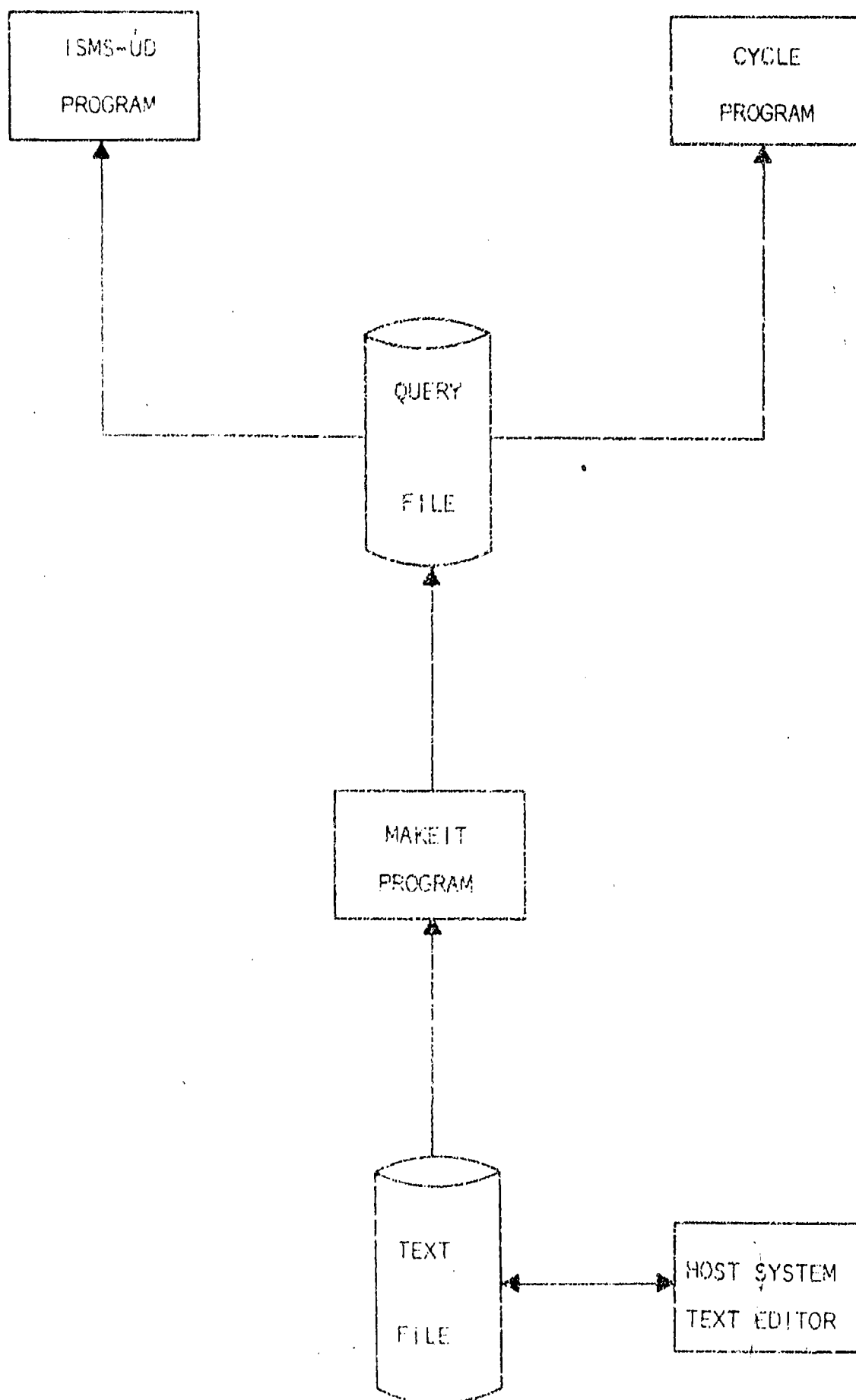


Figure 2: ISMS-UD Basic Block Diagram

ISMS -- A User's View

This section provides the installer with information needed by the user of ISMS-UD. Since the operating systems/executives on various computers differ, we present a generalized discussion of installation alternatives. The University of Dayton ISMS-UD user's guide (Attachment 2) supplements this section. The user's guide serves both as documentation of the software and as a model for user's guides for other installations.

The potential users of ISMS range from teachers in the classroom to researchers in laboratories to public administrators in the meeting room. Due to this wide diversity of users, it is important to install ISMS in such a way that little or no knowledge of the host computer job control language (JCL) is required. There are two approaches to this. One way (and probably the best) is to incorporate the operating system interface as part of the ISMS package. The software would query the user for names of disk files and check for validity of files, etc. This technique usually requires the writing of machine language routines to gain access to file handling routines via FORTRAN language. The second way is to incorporate JCL into "submit" or "catalogued procedures" (hereafter referred to as PROC files), if the operating system used supports this feature.

Following are the three PROC files that were written for the execution of the ISMS-UD software on the University of Dayton Univac Series 70 timesharing computer. These files provide disk file linkage and finally execute the respective program. The "&" (ampersand) allows parameters specified by the user to become part of the JCL.

```
/PROCEDURE N, (&SEQTEXT)
/FILE &SEQTEXT, LINK=DSET10, FCBTYPE=SAM
/FILE RND.&SEQTEXT, LINK=DSET08, FCBTYPE=ISAM, SPACE=(3,3)
/EXECUTE MAKEIT.E
/RELEASE DSET10
/RELEASE DSET08
/ENDPROC RETURN=(PRIMARY)
```

The above PROC is named "MAKEIT" and is invoked by:

```
/DO MAKEIT, (ENV.TEXT.1)
```

where ENV.TEXT.1 is the name of a user-created text file.

The PROC file first links ENV.TEXT.1 to FORTRAN unit reference 10. The second operation is to create a file named "RND.ENV.TEXT.1" as an output file to hold the randomized text. Next, the interactive program MAKEIT.E is invoked and the user then interacts with it. Upon termination, the FORTRAN logical units are released and the PROC is ended.

The most important thing to note about this PROC is the creation of an "invisible" (invisible to the user) file to hold the randomized text. The user only has to remember the name of his or her original text file to access the second "invisible" file. The next two PROC files use the "invisible" file.

```
/PROCEDURE N,(&MODEL, &TEXT)
/FILE &MODEL, LINK=DSET10, FCBTYPE=ISAM
/FILE RND.&TEXT, LINK=DSET08, FCBTYPE=ISAM
/EXECUTE ISMS-UD.E
/RELEASE DSET10
/RELEASE DSET08
/ENDPROC RETURN=(PRIMARY)
```

The above PROC is named "ISMS-UD" and invoked by:

```
/DO ISMS-UD, (ENV1.MOD, ENV.TEXT.1)
```

where "ENV1.MOD" is the file where the model information is to be stored and "ENV.TEXT.1" is the same file from above. It should be noted that the file "RND.ENV.TEXT.1" is actually linked.

```
/PROCEDURE N, (&CYCLE, &TEXT)
/FILE &CYCLE, LINK=DSET20, FCBTYPE=ISAM
/FILE RND.&TEXT, LINK=DSET08, FCBTYPE=ISAM
/EXECUTE CYCLE.E
/RELEASE DSET20
/RELEASE DSET08
/ENPROC RETURN=(PRIMARY)
```

The above PROC is named "CYCLE" and invoked by:

```
/DO CYCLE, (ENV1.CYC, ENV.TEXT.1)
```

where "ENV1.CYC" is the file where the cycle information is to be stored and ENV.TEXT.1 is the same file from above.

The greatest flexibility of ISMS-UD is its ability to access data files. The user of ISMS is primarily concerned with the models he or she has stored on disk files. For this reason, ISMS-UD should be installed on a computer system that is file oriented. This point cannot be stressed strongly enough.

Here we might note the design philosophy behind ISMS-UD. There were many earlier versions of ISMS-UD -- each of which had different designs and characteristics. ISMS-UD Version 2 incorporates what we felt to be the best operating characteristics for software of this type. ISMS-UD will not tutor the user in its use. It is imperative that the user have some knowledge of ISM and its basic principles. The software does have error traps so that the user can never input any erroneous data. The user will gain confidence in ISMS after a small amount of practice on the terminal.

File Formats

ISMS-UD's system block diagram is shown in Figure 3. ISMS-UD is an interactive system, which means that the user must respond to queries for input from the programs while they are executing. Examples of such user-provided input might be the answering of queries that specify a binary matrix representing a hierarchy, operations to be performed on this matrix, results to be displayed, and the like. For the user's convenience, the results of each matrix operation are automatically stored on the user-specified data file, allowing the programs to be terminated (normally or due to computer hardware errors) at any stage and restarted at a later time without loss of the matrix realization. There are three types of files that ISMS-UD will accept. They are: 1) sequential text files, 2) random access text files, and 3) model files. Each of the three types will be discussed below.

The sequential text file is used as input to the program MAKEIT and contains control and text records. Figure 4 shows an example of a sequential text file. The control records precede the text desired for the four types of text records -- R1 clause, R2 clause, R3 clause, and element text. The control records are identified by a "/" (slash)

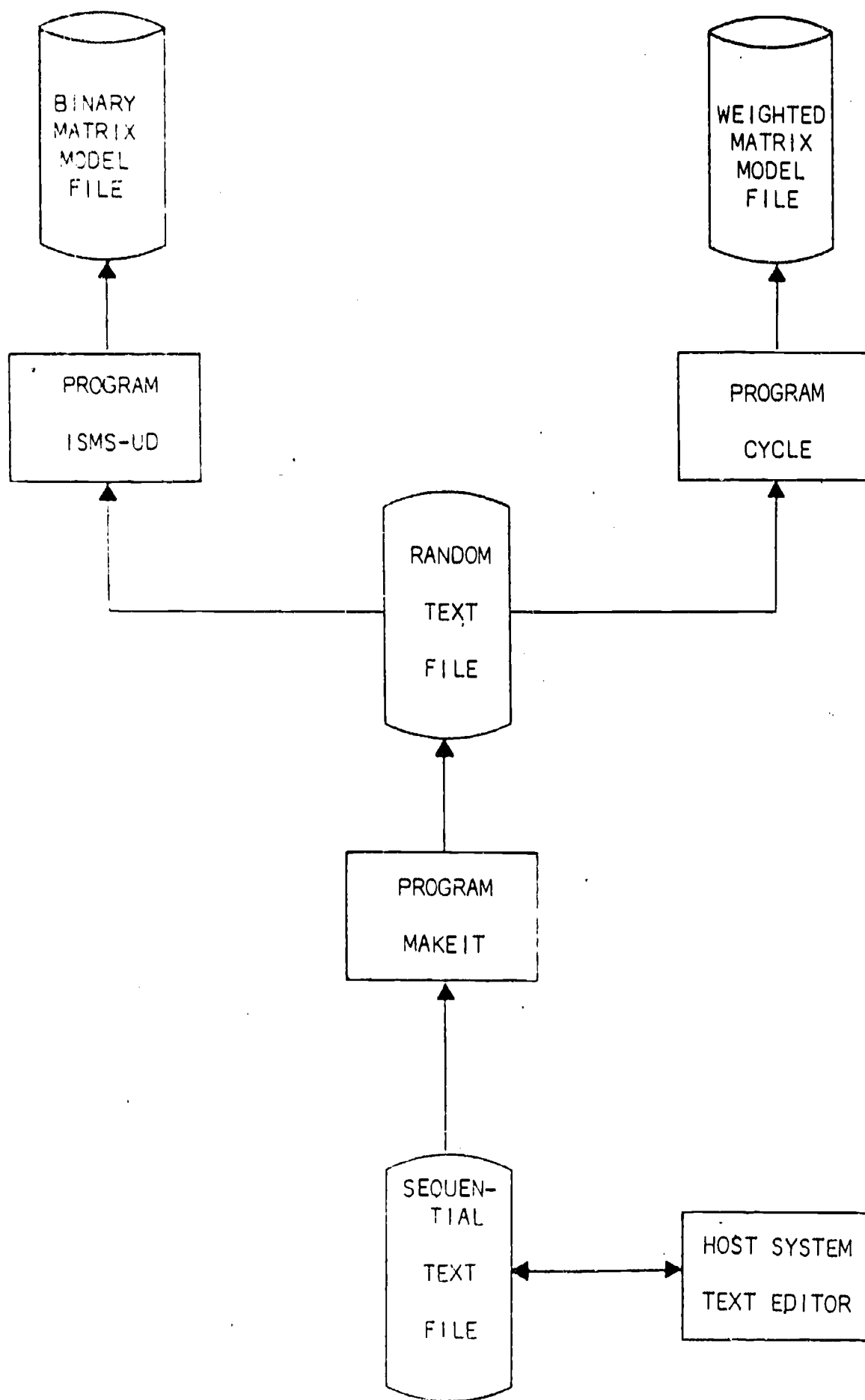


Figure 3: ISMS-UD System Block Diagram

in the first column; therefore, no text record should begin with a "/". All of the four types of text records may have up to 10 lines with a maximum of 60 characters on each line (600 characters). The sequential text file, as its name implies, is accessed sequentially by MAKEIT. The sequential text file may be created and maintained by the host computer file editor, or it could simply be a deck of computer cards.

```
/R1
DOES
/R2
AGGRAVATE OR INTENSIFY
/R3
IN MOST CASES?
/EL
CITIZEN INSECURITY IN THE ABC NEIGHBORHOOD
/EL
GOOD MARKET FOR STOLEN GOODS
/EL
INADEQUATE CRIMINAL REHABILITATION
/EL
LACK OF SUMMER JOBS FOR THE YOUTH IN THE ABC AREA
/EL
LACK OF SEVERITY OF PUNISHMENT FOR CRIMES
/EL
LOW VISIBILITY OF UNIFORMED POLICE
//END OF ELEMENT LIST
```

Figure 4: Sample Text File

The random text file is required by programs ISMS-UD and CYCLE, if the user opts to have these programs present English text queries instead of the standard numeric type. The ISM method requires that elements be accessible in random order; that is, although the elements are ordered by the user in some desired sequential order when creating the sequential text file, the introduction of user responses to relational queries determines a new nonsequential presentation of the future queries. For this reason, a random (direct) access file is utilized. The random text file, created by MAKEIT, contains all of the textual information of the sequential file, but has the property that the text of any specific element may be retrieved independently of the positioning of the file due to a

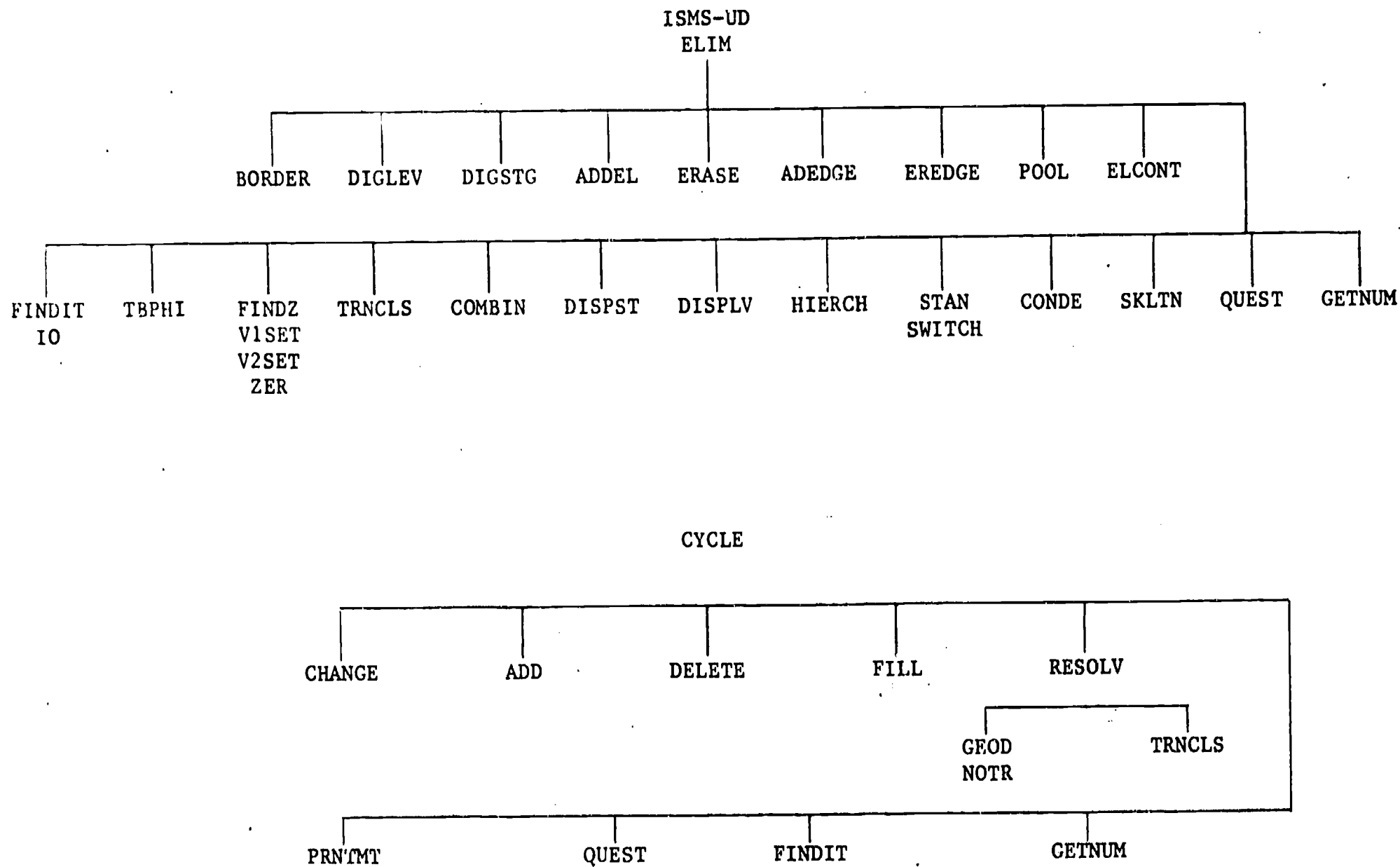
previous access. The format of the randomized text file is shown in Figure 5. Figure 6 shows the format of an individual record of the random text file.

There are two types of model files -- binary and weighted. The binary model file is used by the ISMS-UD program and contains the binary reachability matrix associated with the ISM method. The weighted model file is used by the CYCLE program and contains a weighted adjacency matrix. These two files are accessed sequentially and are opened and closed a number of times during program execution. The formats of both the binary and weighted matrix files are shown in Figure 7.

The ISMS-UD software has been designed in such a way that the binary and weighted operation matrices are core resident in order to speed up execution time. Therefore, the size of these matrix data structures directly affects the amount of core storage required due to the static allocation of memory characteristics of FORTRAN IV. On a computer with a limited amount of interactive core storage, ISMS-UD might be implemented with element and cycle resolution limits of something less than the current limit of 128 and 50. The ISMS-UD program requires 150K bytes, CYCLE requires 120K bytes, and MAKEIT requires 48K bytes on the Univac Series 70/7 computer. These core requirements should be somewhat lower on other computer systems -- the Univac system used at the University of Dayton requires high overhead for FORTRAN-run time routines. The overlay structures used for program ISMS-UD and CYCLE are shown in Figure 8.

The binary matrices used in the ISMS-UD program are of the type "LOGICAL *1". This defines the FORTRAN logical mode and allocates only one byte per matrix cell. Note that this data structure requires one fourth the amount of core storage as a "LOGICAL" matrix. The programmer should try to use the data structure in his or her computer that uses the smallest amount of directly addressable storage. The use of "IMPLICIT INTEGER *2" statements utilizes halfword precision for all integer variables, thereby further reducing the core requirements. All variables requiring full word storage have been declared as "INTEGER".

Figure 8: Suggested ISMS-UD Overlay Scheme



Organizations that have Obtained ISMS-UD

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Attachment No. 1

DOCUMENTATION OF THE ISMS ALGORITHMS

Prepared by:

David R. Yingling, Jr.

This attachment provides a description of some of the algorithms used in ISMS-UD version 2. Although these algorithms have been comprehensively tested, the descriptions should be useful for debugging and future applications of ISM. The algorithms in this section are stated in the form used by Knuth (1973).

ISMS SUBROUTINE/FUNCTION NAME - BORDER

FUNCTION - Embeds elements into an existing reachability matrix via interactive questioning.

USAGE - CALL BORDER(N,RM,INDEX,TTYIN,TTYOUT, QTYPE,SUBREL,XTWDS,SYS)

PARAMETERS

N - Input/Output integer scalar indicating the current number of elements in the argument reachability matrix.

RM - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS". This is the argument reachability matrix.

INDEX - Input/Output integer vector of length "SYS" containing the index set for "RM".

TTYIN - Input integer scalar used as unit number in FORTRAN READ statements directed to interactive terminal.

TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.

QTYPE - Input logical variable indicating type of query to be presented.
QTYPE = .FALSE. causes full text queries to be used.
QTYPE = .TRUE. causes symbolic queries to be used.

SUBREL - Input logical variable indicating type of relationship to be modeled.
SUBREL = .FALSE. implies a non-subordinate relationship is being modeled.
SUBREL = .TRUE. implies a subordinate relationship is being modeled and questioning can therefore be optimized.

TXTWDS - Input integer scalar equal to $(60/nchar) * 10 + 10$, where: $nchar$ = number of display code characters able to be contained in one machine word.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS

FTEXT - Named integer common block of length $5 * \text{TXTWDS}$.

COMMON PARAMETERS

R1 - Integer vector of dimension "TWTWDS".
R1(1-10) contain the number of machine words for each of the ten possible lines of the R1 clause
R1(11-end) contain the lines of text for the R1 clause. the text is packed with as many characters as possible in one machine word and each line starts on a word boundary.

L2 - Integer vector of length "TWTWDS". this vector is used as a scratch area.

R2 - Same attributes as "R1". the R2 clause is contained in this vector.

L2 - Same attributes as "L1".

R3 - Same attributes as "R1". the R3 clause is contained in this vector.

The common block "FTEXT" is used for full text queries. It is the calling routine's responsibility to fill vectors "R1", "R2", and "R3" if full text queries are desired.

REQUIRED ISMS ROUTINES - TBPFI, FINDZ, GETNUM, QUEST, IO, FINDIT

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - Transitive bordering on a reachability matrix. A description follows:

Algorithm TB (Transitive Bordering Algorithm). Given the number of elements currently in a logical matrix, m , the logical matrix, R , and the associated index set vectors, I , embed a new element, b , capitalizing on inference and transitivity.

- TB1. [Initialize.] Set $\text{dimphi} \leftarrow n^2$. Set vector FLAG, $\text{flag}[i] \leftarrow \text{.false.}$, vector Z, $\text{z}[i] \leftarrow \text{.false.}$, matrix PHI, $\text{phi}[i,j] \leftarrow \text{.false.}$, $i=1,2,\dots,\text{dimphi}$, $j=1,2,\dots,\text{dimphi}$.
- TB2. [Make sure physical bounds of arrays and vectors are not exceeded.] If $(n + 1) \leq \text{FORTRAN dimension sizes}$ then go to step TB3. Otherwise, type error message to interactive user and return to calling routine.
- TB3. [Get value for b from interactive user.] Prompt user, call subroutine GETNUM, put value read in b, set $\text{iln} + 1 \leftarrow b$. If $b=0$ then return to calling program.
- TB4. [Make sure that new value, b, is not duplicated.] Call subroutine FINDIT, if b has been already introduced into model matrix, issue error message and then go to step TB3.
- TB5. [Make sure that textual information for b resides on text file.] If full text queries are not being used, skip immediately to step TB6. Read record from file. If no error condition is encountered, go to step TB6. Otherwise, issue error message and then go to step TB3.
- TB6. [Form the transitive bordering inference opportunity matrix.] Form Phi (call subroutine TBPHI).
- TB7. [Calculate the row/column of Phi with the maximum inference potential, zpoint.] Call subroutine FINDZ.
- TB8. [Determine if a question filling the "Y" portion of the reachability matrix is to be asked by examining the value of zpoint.] $\text{relate} \leftarrow \text{zpoint}$. If $\text{zpoint} > m$, then go to step TB10.
- TB9. [Ask for the "X" question and check for valid response.] Call subroutine QUEST(ilrelate , $\text{ilm} + 1$). Read response from user, $\text{zzz} \leftarrow \text{.True.}$. If user typed "AB", go to step TB31, if user didn't type either a "Y" or "N", issue error message and then go to step TB9. If user typed "Y", $\text{zzz} \leftarrow \text{.false.}$ and then go to step TB11. If user typed "N" then go to step TB12.
- TB10. [Ask the "Y" question and check for valid response.] Call subroutine QUEST($\text{ilm} + 1$, $\text{ilzpoint} - m$). Read response from user, $\text{zzz} \leftarrow \text{.false.}$. If user typed "AB", go to step TB31, if user didn't type either a "Y" or "N", issue error message and then go to step TB10. If user typed "Y", $\text{zzz} \leftarrow \text{.true.}$ and then go to step TB11. If user typed "N" then go to step TB12.

TB11. [Process a yes answer to a query.] If zzz = .false., then go to step TB17, else go to step TB13.

TB12. [Process a no answer to a query.] If zzz = .false., then go to step TB17.

TB13. [Initialize.] Set J \leftarrow 1.

TB14. [Search row of Phi for inference to be entered onto R.] If phi[zpoint, J] = .false., then go to step TB16.

TB15. [Fill up R with inferred ones and zeros.] ictr2 \leftarrow J. If ictr2 > n, then ictr2 \leftarrow ictr2 - n. If J > n, then r[n + 1, ictr2] \leftarrow .true.. Otherwise, r[ictr2, m + 1] \leftarrow .false.. Set flag[J] \leftarrow .true., z[J] \leftarrow .true..

TB16. [Loop on J.] J \leftarrow J + 1. If J \leq dimphi, then go to step TB14, else go to step TB20.

TB17. [Initialize.] J \leftarrow 1

TB18. [Search column of Phi for inference to be entered onto R.] if phi[J, zpoint] = .false., then go to step TB20.

TB19. [Fill up R with inferred ones and zeros.] ictr2 \leftarrow J, if ictr2 > n, then ictr2 \leftarrow ictr2 - n. If J > n, then r[m + 1, ictr2] \leftarrow .false., otherwise r[ictr2, n + 1] \leftarrow .true.. Set flag[J] \leftarrow .true., z(J) \leftarrow .true..

TB20. [Loop on J.] J \leftarrow J + 1, if J \leq dimphi, then go to step TB18.

TB21. [Initialize.] J \leftarrow 1

TB22. [Zero out rows and columns of Phi used in steps TB19 and TB15.] If z[J] = .false., then go to step TB24.

TB23. [Zero row and column J.] phi[J, k] \leftarrow .false., phi[k, J] \leftarrow .false., k=1,2,...,dimphi.

TB24. [Loop on J.] J \leftarrow J + 1. If J \leq dimphi, then go to step TB22.

TB25. [Initialize.] J \leftarrow 1

TB26. [Check to see if Phi is all zeros via flag vector.] If flag[J] = .false., then go to step TB30.

- TB27. [Loop on J.] $J \leftarrow J + 1$. If $J \leq \text{dimphi}$, then go to step TB26.
- TB28. [Set main diagonal of R.] $r[m + 1, m + 1] \leftarrow \text{.true.}$, $n \leftarrow n + 1$.
- TB29. [Update permanent file with new R matrix.] Call subroutine IO, then go to step TB1.
- TB30. [Zero out z.] $z[J] \leftarrow \text{.false.}$, $J=1, 2, \dots, \text{dimphi}$. Go to step TB7.
- TB31. [Write out message indicating abort of bordering session.] Write message and then go to step TB1.

IBPHI

ISMS SUBROUTINE/FUNCTION NAME - TBPHI

FUNCTION - Forms the Inference Opportunity Matrix used in the Transitive Bordering Algorithm.

USAGE - CALL TBPHI(N,A,DIMPHI,SUBREL,SYS2,SYS)

PARAMETERS

- N - Input integer scalar indicating the current number of elements in matrix "A".
- A - Input logical two dimensional matrix of dimension "SYS" X "SYS". This is the current reachability matrix model.
- PHI - Output logical two dimensional matrix of dimension "SYS2" X "SYS2". This is the output inference opportunity matrix.
- DIMPHI - Input integer scalar indicating the number of rows/columns in "PHI" which contain information. This variable should be set equal to 2*"N".
- SUBREL - Input logical variable indicating type of relationship being modeled.
SUBREL = .FALSE. implies a non-subordinate relationship is being modeled.
SUBREL = .TRUE. implies a non-subordinate relationship is being modeled and the "PHI" matrix should be calculated differently.
- SYS2 - Input integer scalar used in FORTRAN DIMENSION statements. "SYS2" is equal to 2*"SYS".
- SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUT

None

ALGORITHM EMPLOYED

The algorithm employed is described below:

Algorithm FP (Form the Inference Opportunity Matrix, PHI, for Transitive bordering.) Given the number of elements in a reachability matrix, n , and the logical reachability matrix, A , form an inference opportunity matrix, PHI.

FP1. [Form the "n1" matrix.] Set PHI, $\text{phi}[i,j] \leftarrow a[i,j]$, $\text{phi}[i + n,j + n] \leftarrow a[i,j]$, $\text{phi}[i + n,j] \leftarrow \text{NOT}.a[j,i]$, $j=1,2,\dots,n$, $i=1,2,\dots,n$.

FP2. [Process request for subordinate relationships.] Set PHI, $\text{phi}[i + n,i] \leftarrow \text{true.}$, $i=1,2,\dots,n$ if subrel is true.

FP3. [Multiply.] Set $\text{phi}[i,j + n] \leftarrow \text{false.}$ Then set $\text{phi}[i,j + n] \leftarrow \text{phi}[i,j + n]$.OR.ed with the quantity $\text{phi}[i + n,k]$.AND.ed with $\text{phi}[k + n,j + n]$, $k=1,2,\dots,n$, $j=1,2,\dots,n$, $i=1,2,\dots,n$.

FP4. [Multiply.] Set $\text{phi}[i + n,j] \leftarrow \text{false.}$, then set $\text{phi}[i + n,j] \leftarrow \text{phi}[i + n,j]$.OR.ed with the quantity $\text{phi}[i,k]$.AND.ed with $\text{phi}[k,j + n]$, $k=1,2,\dots,n$, $j=1,2,\dots,n$, $i=1,2,\dots,n$.

FP5. [Set upper right half of matrix to zero.] Set $\text{phi}[i,j] \leftarrow \text{false.}$, $j=n+1,n+2,\dots,2*n$, $i=1,2,\dots,n$.

EINDZ

ISMS SUBROUTINE/FUNCTION NAME - FINDZ

FUNCTION - Finds the row/column of an inference opportunity matrix that has maximum inference potential.

USAGE - CALL FINDZ(DIMPHI, PHI, ZPOINT, FLAG, SYS2)

PARAMETERS

- DIMPHI - Input integer scalar indicating the number of rows and columns filled in the matrix "PHI".
- PHI - Input logical two dimensional matrix of dimensions "SYS2" X "SYS2". This is the inference opportunity matrix.
- ZPOINT - Output integer scalar which contains the row/column number in "PHI" with maximum inference potential. This value is returned to the calling routine.
- FLAG - Input logical vector indicating rows/columns of "PHI" which have been zeroed out as a result of the transitive bordering process. A .true. indicates that a row and column has been zeroed out.
- SYS2 - Input integer scalar used in FORTRAN DIMENSION statements. "SYS2" is equal to 2*"SYS".

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - V1SET, V2SET

REQUIRED FORTRAN ROUTINES - MINO, MAXO

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm FZ (Find a value for Z that has the maximum inference potential.) Given the number of rows and columns in a logical matrix, dimphi, the logical matrix, PHI, and a logical vector indicating rows and columns of PHI that have been zeroed out, FLAG, find the

row/column of PHI that has the maximum inference potential and set variable zpoint equal to this subscript.

- FZ1. [Form the set V.] Set vector V1, $v1[i] \leftarrow$ number of ones in column i of PHI (apply algorithm V1). Set vector V2, $v2[i] \leftarrow$ number of ones in row i of PHI (apply algorithm V2). If $flag[i]$ is .false., $i=1,2,\dots,dimphi$.
- FZ2. [Form the set V'.] Set vector MIN, $min[i] \leftarrow$ the smaller of the two values, $v1[i]$ or $v2[i]$ if $flag[i]$ is .false., $i=1,2,\dots,dimphi$.
- FZ3. [Find the maximum of the minimums of V'.] (Initialize.) Set $bigger \leftarrow 0$, $zpoint \leftarrow 0$.
- FZ4. [Fill up vector Z.] Set $bis \leftarrow$ the largest of the two values, $bigger$ or $min[i]$. If bis is greater than $bigger$, then zero out vector Z, $z[k] \leftarrow .false.$, $k=1,2,\dots,i$. If bis is less than or equal to $min[i]$, then set $z[i] \leftarrow .true.$. If bis is greater than or equal to $bigger$, then set $bigger \leftarrow bis$, $i=1,2,\dots,dimphi$. Do the above step only if $flag[i] = .false.$.
- FZ5. [Form set V" and select a "Z".] Set $bigger \leftarrow 0$. Set $bis \leftarrow$ the greater of the two values, $v1[i] + v2[i]$ or $bigger$. If $bigger$ is greater than bis , then set $bigger \leftarrow bis$ and $zpoint \leftarrow i$, $i=1,2,\dots,dimphi$. Do the above step if $z[i] = .true.$.

V1SEI

ISMS SUBROUTINE/FUNCTION NAME - V1SET

FUNCTION - Counts the number of ones in a selected column, (J), of the inference opportunity matrix.

USAGE - $X = V1SET(J, DIM, SYS2)$

PARAMETERS

| | |
|-----|----------------------------------------------------------------------------------------------------------------|
| MAT | - Input logical two dimensional matrix of dimension "SYS2" X "SYS2". This is the inference opportunity matrix. |
| J | - Input integer scalar indicating the column of "MAT" that the counting operation should be applied to. |
| DIM | - Input integer scalar indicating the last row/column of "MAT" that contains information. |
| SYS | - Input integer scalar used in FORTRAN DIMENSION statements. |

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm V1SET (Count the number of ones in the column of a matrix).
Given the number of elements currently in a matrix, d, the matrix, M, and the number of a column in M, J, count the number of ones in column J of M and assign this value to s.

V1SET1. [Initialize.] Set $s \leftarrow 0$, $i \leftarrow 1$.

V1SET2. [Count.] If $m[i, J] = 1$, then $s \leftarrow s + 1$.

V1SET3. [Loop on i.] $i \leftarrow i + 1$. If $i \leq d$, then go to step V1SET2,
else algorithm is complete.

V2SET

ISMS SUBROUTINE/FUNCTION NAME - V2SET

FUNCTION - Counts the number of ones in a selected row (i) of the inference opportunity matrix.

USAGE - $X = V2SET(MAT, I, DIM, SYS2)$

PARAMETERS

| | |
|------|----------------------------------------------------------------------------------------------------------------|
| MAT | - Input logical two dimensional matrix of dimension "SYS2" X "SYS2". This is the inference opportunity matrix. |
| I | - Input integer scalar indicating the row of "MAT" that the counting operation should be applied to. |
| DIM | - Input integer scalar indicating the last row/column of "MAT" that contains information. |
| SYS2 | - Input integer scalar used in FORTRAN DIMENSION statements. "SYS2" is equal to 2*"SYS". |

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm V2SET (Count the number of ones in the row of a matrix).

Given a logical matrix, M, the number of filled rows and columns in M, d, and the number of the row to count, i, count the total number of ones in row i of M and assign s to this value.

V2SET1. [Initialize.] $s \leftarrow 0$, $j \leftarrow 1$.

V2SET2. [Count.] If $m[i, j] = 1$, then set $s \leftarrow s + 1$.

V2SET3. [Loop on J.] $J \leftarrow J + 1$. If $J \leq d$, then go to step V2SET2.
Otherwise, algorithm is complete.

DIGLEV

ISMS SUBROUTINE/FUNCTION NAME - DIGLEV

FUNCTION - Displays a level formatted digraph of an argument reachability matrix.

USAGE - CALL DIGLEV(N,MATRIX,INDEX,TTYOUT,SYS,)

PARAMETERS

N - Input integer scalar denoting the current number of elements in the argument reachability matrix.

MATRIX - Input logical two dimensional matrix of dimension "SYS" X "SYS". This is the argument reachability matrix.

INDEX - Input integer vector of length "SYS" containing the index set for "MATRIX".

TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - HIERCH,STAN,SWITH,CONDE,ELIM,SKLTN,DISPLV

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm DIGLEV (Display the levels formatted digraph of an argument reachability matrix). Given a reachability matrix, M, associated index set vector, I, and the current number of elements in M, n, display a levels formatted digraph on the interactive terminal.

DIGLEV1. [Rearrange M and I into hierarchial form. Leave result in H and O.] Apply Algorithm HEIRCH.

- DIGLEV2. [Put H and O into standard form.] Apply Algorithm STAN.
- DIGLEV3. [Calculate condensation matrix for H and O.] Apply Algorithm CONDE.
- DIGLEV4. [Calculate skeleton matrix for H] Apply Algorithm SKLTN.
- DIGLEV5. [Print out levels formatted digraph.] Apply Algorithm DISPLV.

DISPLV

ISMS SUBROUTINE/FUNCTION NAME - DISPLV

FUNCTION - Prints a levels formatted digraph from a lower triangular skeleton matrix.

USAGE - CALL DISPLV(N,SKLTN,INDEX,LEVELS,TTYOUT, SYS)

PARAMETERS

N - Input integer scalar indicating the number of elements in the input skeleton matrix.

SKLTN - Input logical two dimensional matrix of dimension "SYS" X "SYS". This is the input lower triangularized skeleton matrix.

INDEX - Input integer vector of length "SYS" containing the index set for "SKLTN".

LEVELS - Input integer vector of length "SYS" containing the number of elements on each level.
LEVELS(I) = Number of elements on each level #i.

TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm DISPLV (Display a levels formatted digraph on the interactive terminal.) Given a logical lower triangularized skeleton matrix, S, its associated index set vector, D, the number of elements in S, n, and a vector filled with the number of elements on each level, L, display a levels formatted digraph.

- DISPLV1. [Initialize.] $i \leftarrow 1$, level $\leftarrow 0$, row $\leftarrow 1$.
- DISPLV2. [Process next level.] level \leftarrow level + 1, row \leftarrow row + 1[level]. Write level number out on interactive terminal, variable level.
- DISPLV3. [Check to see if we are processing level one.] If level = 1, then go to step DISPLV11.
- DISPLV4. [Start processing element #i.] count $\leftarrow 0$, iminus $\leftarrow i - 1$, $j \leftarrow 1$.
- DISPLV5. [Find all elements that element #i reaches to.] If $s[i,j] = 0$, then go to step DISPLV7.
- DISPLV6. [Fill scratch print vector.] count \leftarrow count + 1, list[count] $\leftarrow d[j]$.
- DISPLV7. [Loop on j.] $j \leftarrow j + 1$. If $j \leq$ iminus, then go to step DISPLV5.
- DISPLV8. [All done processing elements #i, print it and its connectives.] Write elements #i and its reachability set on interactive terminal, variables $d[i]$, list[k], $k=1, \dots, \text{count}$.
- DISPLV9. [Are we finished with the algorithm?] $i \leftarrow i + 1$. If $i \leq n$, then algorithm is complete.
- DISPLV10. [Not done with algorithm, are we done with this level on digraph?] If $i = \text{row}$, then go to step DISPLV2. Otherwise, go to step DISPLV3.
- DISPLV11. [Write out level #1 element on terminal.] Write out variable $d[i]$, then go to step DISPLV9.

HIERCH

ISMS SUBROUTINE/FUNCTION NAME - HIERCH

FUNCTION

- Rearranges an input reachability matrix into a levels partitioned hierarchical reachability matrix.

USAGE

- CALL HIERCH(N, INREA, INDXIN, REAH, INDXH, NL, LEVEL, SYS)

PARAMETERS

- N - Input integer scalar denoting number of elements in the input reachability matrix.
- INREA - Input logical two dimensional matrix of dimensions "SYS" X "SYS". This is the input reachability matrix.
- INDXIN - Input integer vector of length "SYS". This is the index set for "INREA".
- REAH - Output logical two dimension matrix of dimensions "SYS" X "SYS". This is the output hierarchical levels partitioned matrix.
- INDXH - Output integer vector of length "SYS". This is the index set for "REAH".
- NL - Output integer scalar indicating the number of hierarchical levels in "REAH".
- LEVEL - Output integer vector of length "SYS" containing the number of elements on each level.
LEVELS(I) = Number of elements on level #i.
- SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm HIERCH (Rearrange an input reachability matrix into a levels partitioned hierarchical reachability matrix.) Given a reachability matrix, R, its associated index set vector, D, the number of elements in R, n, rearrange R according to a levels partition algorithm. Leave this result in an output reachability matrix, H, and associated index set, E.

HIERCH1. [Initialize.] Copy R into H. Set local vector F to zero.
 $h[i,j] \leftarrow r[i,j]$, $f[i] \leftarrow 0$, $j=1,2,\dots,n$, $i=1,2,\dots,n$.

HIERCH2. [Initialize levels partition algorithm.] $nl \leftarrow 0$, $neap \leftarrow 0$.

HIERCH3. [Begin levels partition algorithm.] $nl \leftarrow nl + 1$, $nel \leftarrow 0$.

HIERCH4. [Find an element to process.] $i \leftarrow 1$.

HIERCH5. [Skip past element row/column already processed.] If $f[i] = 1$, then go to step HIERCH10.

HIERCH6. [Initialize subset test.] $i \leftarrow 1$.

HIERCH7. [Test to see if reachable set is not a subset of antecedent set for element #1.] If $r[i,j] = 1$.AND. $r[j,i] = 0$, then go to step HIERCH10.

HIERCH8. [Loop on j.] $j \leftarrow j + 1$. If $j \leq n$, then go to step HIERCH7.

HIERCH9. [Fill up E.] $nel \leftarrow nel + 1$, $e[neap + nel] \leftarrow d[i]$, local vector T, $t[nel] \leftarrow 1$.

HIERCH10. [Loop on i.] $i \leftarrow i + 1$. If $i \leq n$, then go to step HIERCH5.

HIERCH11. [Found all elements on this level.] $neap \leftarrow neap + nel$, vector levels, $level[nl] \leftarrow nel$.

HIERCH12. [Initialize i] $i \leftarrow 1$.
 HIERCH13. [Identify a row/column to be zeroed.] $x \leftarrow t[i]$, $f[x] \leftarrow 1$.
 HIERCH14. [Initialize zero out routine.] $j \leftarrow 1$.
 HIERCH15. [Zero out R.] $h[x,j] \leftarrow 0$, $h[j,x] \leftarrow 0$.
 HIERCH16. [Loop on j.] $j \leftarrow j + 1$. If $j \leq n$, then go to step HIERCH15.
 HIERCH17. [Loop on i.] $i \leftarrow i + 1$. If $i \leq n$, then go to step HIERCH13.
 HIERCH18. [Have all elements been processed?] If $near \leq n$, then go to step HIERCH12.
 HIERCH19. [Initialize i.] $i \leftarrow 1$.
 HIERCH20. [Initialize j.] $j \leftarrow 1$.
 HIERCH21. [Locate proper j subscript.] If $e[i] = d[j]$, then go to step HIERCH23.
 HIERCH22. [Loop on j.] $j \leftarrow j + 1$. If $j \leq n$, then go to step HIERCH21.
 HIERCH23. [Initialize k.] $k \leftarrow 1$.
 HIERCH24. [Rearrange rows of R and store result in H.] $h[i,k] \leftarrow r[j,k]$.
 HIERCH25. [Loop on k.] $k \leftarrow k + 1$. If $k \leq n$, then go to step HIERCH24.
 HIERCH26. [Loop on i.] $i \leftarrow i + 1$. If $i \leq n$, then go to step HIERCH20.
 HIERCH27. [Copy H into R.] $r[i,j] \leftarrow h[i,j]$, $j=1,2,\dots,n$,
 $i=1,2,\dots,n$.
 HIERCH28. [Initialize i.] $i \leftarrow 1$.
 HIERCH29. [Initialize j.] $j \leftarrow 1$.
 HIERCH30. [Locate proper j subscript.] If $e[i] = d[j]$, then go to step HIERCH32.

HIERCH31. [Loop on J.] $J \leftarrow J + 1$. If $J \leq n$, then go to step HIERCH30.

HIERCH32. [Initialize k.] $k \leftarrow 1$.

HIERCH33. [Rearrange columns of R and store in in H.] $h[k,i] \leftarrow r[k,J]$

HIERCH34. [Loop on k.] $k \leftarrow k + 1$. If $k \leq n$, then go to step HIERCH33.

HIERCH35. [Loop on i.] $i \leftarrow i + 1$. If $i \geq n$, then go to step HIERCH29.

HIERCH36. [Copy E into D, copy H into R.] $d[i] \leftarrow e[i]$, $r[i,J] \leftarrow h[i,J]$, $J=1,2,\dots,n$, $i=1,2,\dots,n$.

SIAN

ISMS SUBROUTINE/FUNCTION NAME - STAN

FUNCTION - Converts an input hierarchical reachability matrix into its standard form.

USAGE - CALL STAN(N,MATRIX,INDEX,NLEVEL,LEVELS,SYS)

PARAMETERS N - Input integer scalar denoting the number of elements in "MATRIX".

MATRIX - Input hierarchical reachability matrix, output standard form matrix. This is a logical two dimensional matrix with dimensions "SYS" X "SYS".

INDEX - Input/Output vector of length "SYS" containing the index set for matrix.

NLEVEL - Input integer scalar denoting the number of hierarchical levels in "MATRIX".

LEVELS - Input integer vector of length "SYS" containing the number of elements on each level.
LEVELS(I) = Number of elements on level #I.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - SWITCH

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm STAN (Compute standard form matrix). Given a hierarchical reachability matrix, H, its associated index set vector, V, the

number of elements in H, n, the number of hierarchical levels on H, nlevel, and a vector containing the number of elements on each level, L, compute the standard form of H.

STAN1. [Initialize.] end $\leftarrow 0$, i $\leftarrow 1$.

STAN2. [Continue to check to make sure all non-cycle elements are on top.] start \leftarrow end + 1, end \leftarrow end + 1[i]

STAN3. [Check to see if number of elements is two or less.] If 1[i] .LE. 2, then go to step STAN13.

STAN4. [Find and move all non-cycle elements together on each level.]
ii \leftarrow start

STAN5. [Initialize row check.] row \leftarrow start, last $\leftarrow 0$.

STAN6. [Initialize column check.] col \leftarrow start, nones $\leftarrow 0$.

STAN7. [Count number of ones.] If h[row,col] = 1, then nones \leftarrow nones + 1.

STAN8. [Loop on col.] col \leftarrow col + 1. If col .LE. end, then go to step STAN7.

STAN9. [Check to see if a switch should be performed.] If nones .LT. last, then apply algorithm SWITCH.

STAN10. [Reset variable last.] If nones .GE. last, then set last \leftarrow nones.

STAN11. [Loop on row.] row \leftarrow row + 1. If row .LE. end, then go to step STAN6.

STAN12. [Loop on ii.] ii \leftarrow ii + 1. If ii .LE. end, then go to step STAN5.

STAN13. [Loop on i.] i \leftarrow i + 1. If i .LE. nlevel, then go to step STAN2.

STAN14. [Calculate variable nminus.] nminus \leftarrow n - 1.

STAN15. [Set switch flag to 0.] swit $\leftarrow 0$.

STAN16. [Initialize i.] i $\leftarrow 1$.

STAN17. [Initialize j.] iPlus \leftarrow i + 1, j \leftarrow iPlus.

STAN18. [Check for ones above main diagonal.] If $h[i,j] = 0$, then go to step STAN21.

STAN19. [Check to see if a switch is needed.] $j_{\text{minus}} \leftarrow j - 1$. If $i = j_{\text{minus}}$, then go to step STAN22.

STAN20. [Apply algorithm SWITCH to switch row and col j with row and col $j-1$.] Apply algorithm SWITCH. Set $\text{swit} \leftarrow 1$, $j \leftarrow j - 1$, then go to step STAN19.

STAN21. [Loop on j .] $j \leftarrow j + 1$. If $j \leq n$, then go to step STAN18.

STAN22. [Loop on i .] $i \leftarrow i + 1$. If $i \leq n_{\text{minus}}$, then go to step STAN17.

STAN23. [Check to see if any switching was required.] If $\text{swit} = 1$, then go to step STAN15.

ISMS SUBROUTINE/FUNCTION NAME - SKLTN

FUNCTION - Calculates the nonredundant adjacency matrix (skeleton matrix) given a condensation matrix.

USAGE - CALL SKLTN(N,MATRIX,SYS)

PARAMETERS N - Input integer scalar denoting the number of elements in "MATRIX".

MATRIX - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS". This is the input condensation matrix/output skeleton matrix.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is similar to the one described by R.K. Shyamasundar in IEEE Transactions on Systems, Man, and Cybernetics, Vol. SMC-8, No. 2, February, 1978. It is described below.

Algorithm SKLTN (Calculate the nonredundant adjacency matrix). Given a condensation matrix, C, and the number of elements contained in C, n, calculate the skeleton matrix.

SKLTN1. [Initialize.] nminus \leftarrow n - 1, i \leftarrow 2.

SKLTN2. [Initialize.] iminus \leftarrow i - 1, j \leftarrow 1.

SKLTN3. [Check reachability of node j to node i.] If $c[i,j] = 0$, then go to step SKLTN7.

SKLTN4. [Initialize.] $iplus \leftarrow i + 1$, $k \leftarrow iplus$.

SKLTN5. [Add all nodes to row i that can be reached from node i .] If $c[k,j] = 1$. .AND. $c[k,i] = 1$, then $c[k,j] = 0$.

SKLTN6. [Loop on k .] $k \leftarrow k + 1$. If k .LE. n , then go to step SKLTN5.

SKLTN7. [Loop on j .] $j \leftarrow j + 1$. If j .LE. $iminus$, then go to step SKLTN3.

SKLTN8. [Loop on i .] $i \leftarrow i + 1$. If i .LE. $nminus$, then go to step SKLTN2.

SWITCH

ISMS SUBROUTINE/FUNCTION NAME - SWITCH

FUNCTION - Exchanges two rows and columns on a binary matrix.

USAGE - CALL SWITCH(N,MATRIX,INDEX,ROW,SYS)

PARAMETERS

- N - Input integer scalar denoting the number of elements in "MATRIX".
- MATRIX - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS".
- INDEX - Input/Output vector of length "SYS" containing the index set of "MATRIX".
- SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm SWITCH (Exchange two rows and columns of a binary matrix).
Given a binary matrix, M, its associated index set, R, the number of elements in M, n, and the number of a row and column to be switched; s, switch row and column s with row and column s-1.

SWITCH1. [Initialize.] other \leftarrow s - 1, i \leftarrow 1.

SWITCH2. [Switch rows.] temp \leftarrow m[s,i], m[s,i] \leftarrow m[other,i],
m[other,i] \leftarrow temp.

SWITCH3. [Loop on i.] i \leftarrow i + 1. If i \leq n, then go to step SWITCH2.

SWITCH4. [Initialize.] $i \leftarrow 1$.

SWITCH5. [Switch the columns.] $\text{temp} \leftarrow m[i,s]$, $m[i,s] \leftarrow m[i,\text{other}]$,
 $m[i,\text{other}] \leftarrow \text{temp}$.

SWITCH6. [Loop on i .] $i \leftarrow i + 1$. If $i \leq n$, then go to step
SWITCH5.

SWITCH7. [Switch index set.] $\text{temp} \leftarrow r[s]$, $r[s] \leftarrow r[\text{other}]$, $r[\text{other}] \leftarrow \text{temp}$.

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ISMS FUNCTION/SUBROUTINE NAME - ZER

FUNCTION - Zeros out a logical vector.

USAGE - CALL ZER(VECTOR,I,SYS)

PARAMETERS

| | |
|--------|--------------------------------------------------------------------------------------|
| VECTOR | - Input/Output logical vector of length "SYS". |
| I | - Input integer scalar denoting starting index of values to be retained in "VECTOR". |
| SYS | - Input integer scalar used in FORTRAN DIMENSION statements. |

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm ZER (Zero out selected positions on a logical vector).

Given a logical vector, V, and the value of the first index of V to retain, i, zero out the vector.

ZER1. [Initialize.] ii \leftarrow i - 1, J \leftarrow 1.

ZER2. [Zero out.] v[J] \leftarrow 0.

ZER3. [Loop on J.] J \leftarrow J + 1. If J .LE. ii, then go to step ZER2.

IO

ISMS SUBROUTINE/FUNCTION NAME - IO

FUNCTION - Reads and writes ISMS information on disk file.

USAGE - CALL IO(N,MAT,INDEX,UNITNO,READ,SYS)

PARAMETERS N - Input/Output integer scalar contains the number of elements in "MAT".

MAT - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS". This is the ISM binary matrix.

INDEX - Input/Output integer vector of length "SYS". This is the index set vector for "MAT".

UNITNO - Input integer scalar used as unit number in FORTRAN READ and WRITE statements for disk file.

READ - Input logical scalar used for determining if a read or write operation is desired.

READ = .TRUE. means a read operation is requested.

READ = .FALSE. means a write operation is requested.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm 10 (Read and write ISM information on disk file). Given a logical switch variable, r , read or write the ISM binary matrix, index set, and number of elements.

I01. [Determine if a read or write operation is desired.] If $r = 1$, then go to step I03.

I02. [Write operation.] Rewind file. Write n , MAT, and INDEX. Algorithm is complete.

I03. [Read operation.] Rewind file. Read n , MAT, and INDEX. Algorithm is complete.

DISESI

ISMS SUBROUTINE/FUNCTION NAME - DISPST

FUNCTION - Prints a stages formatted digraph given an upper triangularized skeleton matrix.

USAGE - CALL DISPST(N,MATRIX,INDEX,STAGES,NSTAGE,TTYOUT,SYS)

PARAMETERS

N - Input integer scalar equal to the number of elements in "MATRIX".

MATRIX - Input logical two dimensional matrix of dimensions "SYS" X "SYS". This is the upper triangularized skeleton matrix.

INDEX - Input integer vector of length "SYS" containing the index set of "MATRIX".

STAGES - Input integer vector of length "SYS" containing the number of elements on each stage.
STAGES(I) = Number of elements on stage #I.

NSTAGE - Input integer scalar equal to the total number of stages on "MATRIX".

TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

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ALGORITHM EMPLOYED

- The algorithm employed is described below:

Algorithm DISPST (Print a stages formatted digraph). Given the number of elements currently in an upper triangularized skeleton matrix, n , the skeleton matrix, S , the index set for S , T , a vector containing the number of elements on each stage, R , and the total number of stages, v , print the stages formatted digraph.

DISPST1. [Initialize.] row $\leftarrow 1$, stage $\leftarrow 0$, $i \leftarrow 1$.

DISPST2. [Initialize this stage.] stage \leftarrow stage + 1, row \leftarrow row + $r[\text{stage}]$. Write out stage number on interactive terminal, variable stage.

DISPST3. [Check to see if last stage is to be processed.] If stage = v , then go to step DISPST13.

DISPST4. [Initialize element search.] count $\leftarrow 0$, istart $\leftarrow i + 1$, $j \leftarrow$ istart.

DISPST5. [Find all elements that element # i reaches.] If $s[i,j] = 0$, then go to step DISPST7.

DISPST6. [Found one, keep a record of it.] count \leftarrow count + 1, fill local vector LIST, $l[\text{count}] \leftarrow t[j]$.

DISPST7. [Loop on j .] $j \leftarrow j + 1$. If $j \leq n$, then go to step DISPST5.

DISPST8. [Special printing routine for elements with no reachability set.] If count = 0, then go to step DISPST13.

DISPST9. [All done processing element # i , print on terminal.] Write out on interactive terminal, $t[i]$, $l[k]$, $k=1,2,\dots,\text{count}$.

DISPST10. [Increment i .] $i \leftarrow i + 1$.

DISPST11. [Finished with algorithm?] If $i > n$, then algorithm is complete.

DISPST12. [Finished with this stage?] If $i = \text{row}$, then go to step DISPST2, else go to step DISPST3.

DISPST13. [Process last stage elements.] Write out on interactive terminal $t[i]$, then go to step DISPST10.

EINDII

ISMS SUBROUTINE/FUNCTION NAME - FINDIT

FUNCTION

- Determines if two elements are contained in the index set of a matrix.

USAGE

- CALL FINDIT(N,N1,N2,S1,S2,INDEX,FOUND1,FOUND2,SYS)

PARAMETERS

- N - Input integer scalar denoting number of elements contained in "INDEX".
- N1 - Input integer scalar equal to the first element number to be found.
- N2 - Input integer scalar equal to the second element number to be found.
- S1 - Output integer scalar which is set equal to the subscript of "INDEX" that "N1" is found.
- S2 - Output integer scalar which is set equal to the subscript of "INDEX" that "N2" is found.
- INDEX - Input integer vector of length "SYS". This is an index set vector for a matrix.
- FOUND1 - Output logical scalar set to .true. if "N1" was found, set to .false. otherwise.
- FOUND2 - Output logical scalar set to .true. if "N2" was found, set to .false. otherwise.
- SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS

- None

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REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm FINDIT (Find the subscript values for two elements in any index set vector). Given two element numbers to be found, n1 and n2, the number of elements in a given index set, n, the index set vector, I, find the vector index for each, s1 and s2, and set a logical variable for each, f1 and f2, if found.

FINDIT1. [Initialize.] f1 \leftarrow 0, f2 \leftarrow 0, s1 \leftarrow 0, s2 \leftarrow 0, J \leftarrow 1.

FINDIT2. [Search.] If n1 = i[J], then set s1 \leftarrow J. If n2 \leftarrow i[J], then set s2 \leftarrow J.

FINDIT3. [Loop on J.] J \leftarrow J + 1. If J .LE. n, then go to step FINDIT2.

FINDIT4. [Set logical variables.] If s1.gt. 0, then set f1 \leftarrow 1. If s2 .GT. 0, then set f2 \leftarrow 1.

DIGSIG

ISMS SUBROUTINE/FUNCTION NAME - DIGSTG

FUNCTION - Displays a stages formatted digraph from a given reachability matrix.

USAGE - CALL DIGSTG(N,MATRIX,INDEX,TTYOUT,SYS)

PARAMETERS N - Input integer scalar denoting the number of elements in "MATRIX".

MATRIX - Input logical two dimensional matrix of dimensions "SYS" X "SYS". This is the input reachability matrix.

INDEX - Input integer vector of length "SYS". This is the index set for "MATRIX".

TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - HIERCH,STAN,SWITCH,CONDE,ELIM,SKLTN,DISPST

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm DIGSTG (Display the digraph of a reachability matrix in a stages format). Given a reachability matrix, R, its associated index set, X, and the number of elements in R, n, display the digraph in a stages format.

DIGSTG1. [Check for error condition.] If n .GT. 0, then go to step DIGSTG3.

DIGSTG2. [Write out error message.] Write error to interactive terminal and then algorithm is complete.

DIGSTG3. [Initialize.] $i \leftarrow 1$.

DIGSTG4. [Initialize.] $j \leftarrow 1$.

DIGSTG5. [Transpose input reachability matrix.] Set local logical matrix, L . $t[i,j] \leftarrow r[j,i]$.

DIGSTG6. [Loop on j .] $j \leftarrow j + 1$. If $j \leq n$, then go to step DIGSTG5.

DIGSTG7. [Loop on i .] $i \leftarrow i + 1$. If $i \leq n$ then go to step DIGSTG4.

DIGSTG8. [Copy T into R .] $r[i,j] \leftarrow t[i,j]$, $j=1,2,\dots,n$, $i=1,2,\dots,n$.

DIGSTG9. [Levels partition T .] Apply Algorithm HIERCH.

DIGSTG10. [Put T into standard form.] Apply Algorithm STAN.

DIGSTG11. [Compute condensation matrix of T .] Apply Algorithm CONDE.

DIGSTG12. [Calculate skeleton matrix of T .] Apply Algorithm SKLTN.

DIGSTG13. [Initialize.] $i \leftarrow 1$.

DIGSTG14. [Initialize.] $j \leftarrow 1$.

DIGSTG15. [Transpose T , leave result in R .] $r[i,j] \leftarrow t[j,i]$.

DIGSTG16. [Loop on j .] $j \leftarrow j + 1$. If $j \leq n$, then go to step DIGSTG15.

DIGSTG17. [Loop on i .] $i \leftarrow i + 1$. If $i \leq n$, then go to step DIGSTG14.

DIGSTG18. [Print out stages digraph.] Apply Algorithm DISPST.

ELCONT

ISMS SUBROUTINE/FUNCTION NAME - ELCONT

FUNCTION - Performs the elementary contraction process.

USAGE - CALL ELCONT(N, REA, INDEX, TTYIN, TTYOUT, SYS)

PARAMETERS

- N - Input/Output integer scalar equal to the number of elements in "REA".
- REA - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS".
- INDEX - Input/Output integer vector of length "SYS" containing the index set of "REA".
- TTYIN - Input integer scalar used as unit number in FORTRAN READ statements directed to interactive terminal.
- TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.
- SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - GETNUM, FINDIT, HIERCH, STAN, SWITCH, CONDE, ELIM, SKLTN, COMBIN, TRNCLS, IO

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm ELCONT (Contract two adjacent elements on different levels or stages). Given a reachability matrix, R, its associated index

set vectors, T, and the number of elements in R, n, ask the user for two elements and perform the elementary contraction process.

- ELCONT1. [Accept two number from interactive terminal.] Apply Algorithm GETNUM.
- ELCONT2. [Check to see if numbers are zero.] If $u = 0$ or $v = 0$, then algorithm is complete.
- ELCONT3. [Check to see if numbers are in index set.] Apply Algorithm FINDIT. If u and/or v are/is not found, then write appropriate error message on interactive terminal and then go to step ELCONT1.
- ELCONT4. [Set nn .] $nn \leftarrow n$.
- ELCONT5. [Calculate non-redundant adjacency matrix for "REA".] Apply Algorithm HIERCH, Algorithm STAN, Algorithm CONDE, Algorithm SKLTN. Fill up local vector Z with index set of non-redundant adjacency matrix and local matrix A with matrix.
- ELCONT6. [Check to see if u and v are adjacent.] Apply Algorithm FINDIT in order to obtain the row and column subscripts for u and v , iu , iv . If u and v are found on A and $a[iu, iv] = 1$, then go to step ELCONT8.
- ELCONT7. [Element u is not adjacent to element v .] Write appropriate error message on interactive terminal and then go to step ELCONT1.
- ELCONT8. [Accept new index value from interactive terminal for contracted elements.] Apply Algorithm GETNUM. If $newnam = 0$, then the algorithm is complete.
- ELCONT9. [Check to see if $newnam$ is already in index set.] Apply Algorithm FINDIT. If $newnam$ is already used, then write appropriate error message on interactive terminal and then go to step ELCONT8.
- ELCONT10. [Combine elements u and v and use $newnam$ as index value.] Apply Algorithm COMBIN.
- ELCONT11. [Calculate reachability of new element.] Apply Algorithm TRNCLS.
- ELCONT12. [Update permanent file with new matrix.] Apply Algorithm IO and then go to step ELCONT1.

CONDE

ISMS SUBROUTINE/FUNCTION NAME - CONDE

FUNCTION - Computes the condensation matrix of a given standard form matrix.

USAGE - CALL CONDE(N,MATRIX,INDEX,LEVELS,TTYOUT,TYPE,SYS)

PARAMETERS

N - Input/Output integer scalar denoting the number of elements in the input standard form matrix/number of elements in the output condensation matrix.

MATRIX - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS". This is the input standard form matrix/output condensation matrix.

INDEX - Input/Output integer vector of length "SYS". This is the index set of the input standard form matrix/output condensation matrix.

LEVELS - Input/Output integer vector of length "SYS". This vector contains the number of elements on each level of the input standard form matrix/output condensation matrix.

TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.

TYPE - Input logical scalar used to determine if the printing of cycles at the interactive terminal is to be performed.
TYPE = .TRUE. means Print cycles.
TYPE = .FALSE. means not to print cycles.

SYS - Input integer scalar used in FORTRAN
DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - ELIM

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described
below:

Algorithm CONDE (Compute the condensation matrix of a given standard form matrix.) Given a standard form matrix, S, its associated index set vector, R, the number of elements in S, n, a vector with the number of elements on each level on S, L, and a variable that tells whether or not to print the cycles, b, reduce all cycle sets to a single proxy element by eliminating elements to finally yield the condensation matrix.

CONDE1. [Initialize i.] $i \leftarrow 1$.

CONDE2. [Initialize.] $\text{count} \leftarrow 1$, $j \leftarrow i + 1$, local scratch vector T, $t[i] \leftarrow r[i]$.

CONDE3. [Check for a one above main diagonal (cycle).] If $s[i,j] = 0$, then go to step CONDE11.

CONDE4. [Put found cycle element in print out list and eliminate from matrix.] $\text{count} \leftarrow \text{count} + 1$, $t[\text{count}] \leftarrow r[j]$, apply algorithm ELIM.

CONDE5. [Initialize.] $\text{posin} \leftarrow 0$, $ii \leftarrow 1$.

CONDE6. [Find proper element to change in L.] $\text{posin} \leftarrow \text{posin} + 1[i]$. If $\text{posin} \geq j$, then go to step CONDE8.

CONDE7. [Loop on ii.] $ii \leftarrow ii + 1$. If $ii \leq n$, then go to step CONDE6.

CONDE8. [Reduce number of elements on level where an element was eliminated.] $l[ii] \leftarrow l[ii] - 1$.

CONDE9. [Any more elements in this cycle set?] If $s[i,j] = 1$.AND. $j \leq n$, then go to step CONDE4,

CONDE10. [Write cycle out to interactive terminal.] If $b = 1$, then write out $t[iii]$, $iii=1,2,\dots,\text{count}$.

CONDE11. [Loop on i.] $i \leftarrow i + 1$. If i .LT. n , then go to step
CONDE2.

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ISMS SUBROUTINE/FUNCTION NAME - GETNUM

FUNCTION - Reads n unsigned integers from an interactive terminal in a free format.

USAGE - CALL GETNUM(ARRAY,N,TTYIN,TTYOUT)

PARAMETERS

ARRAY - Output integer vector of length "N". This vector contains the number as read from the terminal
ARRAY(1) = First number, etc.

N - Input integer equal to the number of integers to be read.

TTYIN - Input integer scalar used as unit number in FORTRAN READ statements directed to interactive terminal.

TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINE - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm GETNUM (Read "n" unsigned integers in a free format). Given the vector to store the integers, A, the number of integers to read, n, read "n" unsigned integers from the interactive terminal.

GETNUM1. [Read in string from terminal.] Read into local vector B.

GETNUM2. [Initialize.] 1 \leftarrow n, a[1] \leftarrow 0, power \leftarrow 0, i \leftarrow 1.

GETNUM3. [Search backwards through string looking for delimiters.] $k \leftarrow 81 - 1$. If $b[k]$ = blank or a comma, then go to step GETNUM9.

GETNUM4. [Initialize.] $j \leftarrow 1$, $ii \leftarrow j - 1$.

GETNUM5. [Check against holerith constant vector (nums).] If $b[k]$.NE. $nums[j]$, then go to step GETNUM7.

GETNUM6. [Construct integer.] $a[i] \leftarrow a[i] + (ii * (10^{**power}))$, $power \leftarrow power + 1$, then go to step GETNUM12.

GETNUM7. [Loop on j.] $j \leftarrow j + 1$. If j .LE. 10, then go to step GETNUM5.

GETNUM8. [Character found was not numeric.] Write error message and then go to step GETNUM1.

GETNUM9. [Check for end of number.] If $a[i] = 0$, then go to step GETNUM12.

GETNUM10. [Decrement l and check for completion.] $l \leftarrow l - 1$. If $l = 0$, then go to step GETNUM13.

GETNUM11. [Initialize.] $a[i] \leftarrow 0$, $power \leftarrow 0$.

GETNUM12. [Loop on i.] $i \leftarrow i + 1$. If i .LE. 80, then go to step GETNUM3.

GETNUM13. [Initialize.] $i \leftarrow 1$.

GETNUM14. [Check magnitude of each number.] If $a[i]$.GT. 99999, then go to step GETNUM16.

GETNUM15. [Loop on.] $i \leftarrow i + 1$. If i .LE. n , then go to step GETNUM14. Otherwise, algorithm is complete.

GETNUM16. [Numbers too large.] Write out error message and then go to step GETNUM1.

COMBIN

ISMS SUBROUTINE/FUNCTION NAME - COMBIN

FUNCTION - Combines two rows and columns in a reachability matrix.

USAGE - CALL COMBIN(N, REA, INDEX, IU, IV, NEWNAM, SYS)

PARAMETERS

| | |
|--------|----------------------------------------------------------------------------------------------------------------------|
| N | - Input/Output integer scalar equal to the number of elements in "REA". |
| REA | - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS". This is the argument reachability matrix. |
| INDEX | - Input/Output integer vector of length "SYS" containing the index set of "REA". |
| IU | - Input integer scalar equal to the row/column subscript of the first elements to be combined. |
| IV | - Input integer scalar equal to the row/column subscript of the second element to be combined. |
| NEWNAM | - Input integer scalar equal to the integer to be used in the index set for the combined elements. |
| SYS | - Input integer scalar used in FORTRAN DIMENSION statements. |

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - ELIM

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED

- The algorithm employed is described below:

Algorithm COMBIN (Replace row and column iv with the boolean sum of row and column iu and iv). Given the number of elements in a reachability matrix, n , the reachability matrix, R , its associated index set vector, T , the two row/column indices to combine, iu and iv , and the value to be replaced in the index set for the combined element, $newnam$, combine iu and iv via boolean sum method and replaced index name with $newnam$.

COMBIN1. [Initialize.] $i \leftarrow 1$.

COMBIN2. [Replace row iv with boolean sum of rows iu and iv .] $r[iv,i] \leftarrow r[iu,i] \text{ .OR. } r[iv,i]$.

COMBIN3. [Replace column iv with boolean sum of columns iu and iv .] $r[i,iv] \leftarrow r[i,iu] \text{ .OR. } r[i,iv]$.

COMBIN4. [Loop on i .] $i \leftarrow i + 1$. If $i \text{ .LE. } n$, then go to step COMBIN2.

COMBIN5. [Replace iv 's index with $newnam$.] $t[iv] \leftarrow newnam$.

COMBIN6. [Erase row, column, and index for iu .] Apply Algorithm ELIM.

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ISMS SUBROUTINE/FUNCTION NAME - QUEST

FUNCTION

- Displays the queries used during an embedding session with ISM.

USAGE

- CALL QUEST(EL1,EL2,TTYOUT,QTYPE,TXTWDS)

PARAMETERS

- EL1 - Input integer scalar which is equal to the index value of the first element to be displayed.
- EL2 - Input integer scalar which is equal to the index value of the second element to be displayed.
- TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.
- QTYPE - Input integer scalar used to determine if full text or symbolic queries
QTYPE = .TRUE. means that symbolic queries will be used.
QTYPE = .FALSE. means that full text queries will be used.
- TXTWDS - Input integer scalar which is equal to the number of machine words required to hold 600 display code characters plus 10.

COMMON BLOCKS

- FTEXT - Named integer common block

See the description under subroutine BORDER for "FTEXT" parameters

REQUIRED ISMS ROUTINES

- None

REQUIRED FORTRAN ROUTINES

- None

ALGORITHM EMPLOYED

- The algorithm employed is described below:

Algorithm QUEST (Display the ISM queries in either full text or symbolic formats). Given the index values of each element to be displayed, e1 and e2, a switch variable which determines what type of queries to use, q, and a common block filled with the introductory, relational, and qualifying phrases, display the query.

QUEST1. [Determine query type to be presented.] If $q = 1$, then go to step QUEST12.

QUEST2. [Calculate direct access file offset.] $i1 \leftarrow e1 + 4$, $i2 \leftarrow e2 + 4$.

QUEST3. [Read elements' text file and put in common block.] Read record i1 and place in "L1". Read record i2 and place in "L2".

QUEST4. [Initialize.] $offset \leftarrow 0$, $i \leftarrow 1$.

QUEST5. [Initialize.] $i1 \leftarrow 0$, $i2 \leftarrow 0$, $j \leftarrow 1$.

QUEST6. [Check length indicator.] (NOTE: a local vector, B, is equivalent to the common block.) If $b[j + offset] = 0$, then go to step QUEST9.

QUEST7. [Compute length and location of print line.] $length \leftarrow b[j + offset]$, $i1 \leftarrow i2 + 1$, $i2 \leftarrow i1 + length - 1$.

QUEST8. [Print line of interactive terminal.] Write out $b[c + offset + 10]$, $c = i1, \dots, i2$.

QUEST9. [Loop on j.] $j \leftarrow j + 1$. If $j \leq 10$, then go to step QUEST6.

QUEST10. [Update offset into common block.] $offset \leftarrow offset + txtwds$.

QUEST11. [Loop on i.] $i \leftarrow i + 1$. If $i \leq 5$, then go to step QUEST5. Otherwise algorithm is complete.

QUEST12. [Present symbolic (numeric) queries.] Write out e1, e2.

ISMS SUBROUTINE/FUNCTION NAME - ELIM

FUNCTION - Eliminates an element from a given binary matrix.

USAGE - CALL ELIM(N,MATRIX,INDEX,DELETE,SYS,)

PARAMETERS

N - Input/Output integer scalar denoting the number of elements in the argument matrix.

MATRIX - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS". This is the matrix to be operated on.

INDEX - Input/Output integer vector of length "SYS". This is the index set of "MATRIX".

DELETE - Input integer scalar denoting the subscript of the row and column of the element to be eliminated.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm ELIM (Eliminate an element from a matrix). Given an argument matrix, M, its associated index set vector, R, the number of elements in M, n, and the row and column to eliminate, d, remove row and column d on M and shift the matrix up and to the left to get rid of blank row and column.

ELIM1. [Initialize.] $n_{\text{minus}} \leftarrow n - 1$.
 ELIM2. [Check to see if row and column to be deleted is last logical position on M.] If d .EQ. n , then go to step ELIM14.
 ELIM3. [Initialize row1.] $\text{row1} \leftarrow d$.
 ELIM4. [Initialize col.] $\text{row2} \leftarrow \text{row1} + 1$, $\text{col} \leftarrow 1$.
 ELIM5. [Move all columns below "d" over by 1.] $m[\text{row1}, \text{col}] \leftarrow m[\text{row2}, \text{col}]$.
 ELIM6. [Loop on col.] $\text{col} \leftarrow \text{col} + 1$. If col .LE. n , then go to step ELIM5.
 ELIM7. [Initialize row.] $\text{row} \leftarrow 1$.
 ELIM8. [Move all rows below "d" up by 1.] $m[\text{row}, \text{row1}] \leftarrow m[\text{row}, \text{row2}]$.
 ELIM9. [Loop on row.] $\text{row} \leftarrow \text{row} + 1$. If row .LE. n , then go to step ELIM8.
 ELIM10. [Loop on row.] $\text{row1} \leftarrow \text{row1} + 1$, then go to step ELIM4.
 ELIM11. [Initialize row1.] $\text{row1} \leftarrow d$.
 ELIM12. [Fix up index set.] $\text{row2} \leftarrow \text{row1} + 1$, $r[\text{row1}] \leftarrow r[\text{row2}]$.
 ELIM13. [Loop on row1.] $\text{row1} \leftarrow \text{row1} + 1$. If row1 .LE. n_{minus} , then go to step ELIM12.
 ELIM14. [Set n to reflect deleted element.] $n \leftarrow n - 1$.

ADDEL

ISMS SUBROUTINE/FUNCTION NAME - ADDEL

FUNCTION - Adds elements to a reachability matrix.

USAGE - CALL ADDEL(N,MATRIX,INDEX,TTYIN,TTYOUT,
SYS)

PARAMETERS

N - Input/Output integer scalar indicating number of elements currently on "MATRIX".

MATRIX - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS". This is the binary reachability matrix that elements will be added to.

INDEX - Input/Output integer vector of length "SYS" containing the index set of "MATRIX".

TTYIN - Input integer scalar used as unit number in FORTRAN READ statements directed to interactive terminal.

TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - GETNUM,FINDIT

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm ADDEL (Adds elements to a reachability matrix). Given the number of elements currently in a reachability matrix, n, the

reachability matrix, R, and its associated index set vector, T, accept and add element numbers from interactive terminal.

ADDEL1. [Make sure memory limits are not exceeded.] If $n \geq \text{SYS}$, then write error message and then algorithm is complete.

ADDEL2. [Accept element number to be added.] Apply Algorithm GETNUM to get $n1$. If $n1 = 0$, then algorithm is complete.

ADDEL3. [Make sure that $n1$ does not already exist in index set.] Apply Algorithm FINDIT. If $n1$ already exists, then write out error message and then go to step ADDEL2.

ADDEL4. [Add element.] $n \leftarrow n + 1$, $t[n] \leftarrow n1$, $r[j,n] = 0$, $r[n,j] = 0$, $j=1,2,\dots,n$. Set $r[n,n] = 1$ and then go to step ADDEL1.

ISMS SUBROUTINE/FUNCTION NAME - POOL

FUNCTION - Combines two elements on the same level or stage that are not connected.

USAGE - CALL POOL(N, REA, INDEX, TTYIN, TTYOUT, SYS)

PARAMETERS

N - Input/Output integer scalar equal to the number of elements in "REA".

REA - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS".

INDEX - Input/Output integer vector of length "SYS" that contains the index set of "REA".

TTYIN - Input integer scalar used as unit number in FORTRAN READ statements directed to interactive terminal.

TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - GETNUM, FINDIT, HIERCH, COMBIN, TRNCLS

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm POOL (Combines two elements on the same level or stage).
 Given a reachability matrix, R, its associated index set vector, T,
 and the number of elements in R, n, ask the user for two elements
 and perform the pooling process.

POOL1. [Accept two elements from interactive terminal.] Apply Algorithm GETNUM.

POOL2. [Check if numbers are zero.] If $u = 0$ or $v = 0$, then algorithm is complete.

POOL3. [Check if numbers are in index set.] Apply Algorithm FINDIT. If u and/or v are/is not found, then write appropriate error message and then go to step POOL1.

POOL4. [Initialize.] $stages \leftarrow 0$.

POOL5. [Calculate the number of elements on each level of "REA".] Apply Algorithm HIERCH to fill up local vector L and H .

POOL6. [Initialize.] $f1 \leftarrow 0$, $f2 \leftarrow 0$, $start \leftarrow 1$, $i \leftarrow 1$.

POOL7. [Initialize.] $end \leftarrow start + l[i] - 1$, $J \leftarrow start$.

POOL8. [Determine if u and v are on same level or stage.] If $h[J] = u$, then $f1 \leftarrow 1$. If $h[J] = v$, then $f2 \leftarrow 1$.

POOL9. [Loop on J .] $J \leftarrow J + 1$. If $J \leq end$, then go to step POOL8.

POOL10. [Check if u and v are on same level or stage.] If $f1 = 1$ and $f2 = 1$, then go to step POOL19.

POOL11. [Check if either u or v was found on this level or stage.] If $f1 = 1$ or $f2 = 1$, then go to step POOL13.

POOL12. [Check next level.] $start \leftarrow end + 1$, $i \leftarrow i + 1$. If $i \leq \text{number of levels or stages}$ then go to step POOL7.

POOL13. [Not on same level, see if on same stages.] If $stages = 1$, then go to step POOL18.

POOL14. [Transpose R to obtain number of elements on each stage.] Set local matrix Z , $z[i,j] \leftarrow r[j,i]$, $i=1,2,\dots,n$, $j=1,2,\dots,n$.

POOL15. [Copy transposed matrix into R .] $r[i,j] \leftarrow z[i,j]$, $i=1,2,\dots,n$, $j=1,2,\dots,n$.

POOL16. [Calculate number of elements on each stage of "REA".] Apply Algorithm HIERCH to fill up local vectors L and H .

- POOL17. [Transpose again in order to get back original matrix.] $r[i, j] \leftarrow z[i, i]$, $i=1,2,\dots,n$, $j=1,2,\dots,n$. Set stages, stages $\leftarrow 0$, then go to step POOL6.
- POOL18. [Pooling error message, not on same level or stage.] Write appropriate error message to interactive terminal and then go to step POOL1.
- POOL19. [Determine row and column subscripts for u and v on "REA".] Apply Algorithm FINDIT to set variables i_u and i_v .
- POOL20. [Accept index set name for pooled elements from interactive terminal.] Apply Algorithm GETNUM to get newnam.
- POOL21. [Check to see if newnam is zero.] If newnam = 0, then algorithm is complete.
- POOL22. [Check to see if newnam is already in index set.] Apply Algorithm FINDIT. If newnam is already in index set, then write appropriate error message on interactive terminal and then go to step POOL20.
- POOL23. [Combine elements u and v and use newnam as index in index set.] Apply Algorithm COMBIN.
- POOL24. [Calculate reachability of new element.] Apply Algorithm TRNCLS and then go to step POOL1.

EREDGE

ISMS SUBROUTINE/FUNCTION NAME - EREDGE

FUNCTION - Erases an edge from the minimum edge digraph.

USAGE - CALL EREDGE(N, REA, INDEX, TTYIN, TTYOUT, SYS)

PARAMETERS

- N - Input integer scalar indicating the number of elements currently in "REA".
- REA - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS". This is the binary reachability matrix containing the minimum edge digraph.
- INDEX - Input integer vector of length "SYS" containing the index set for "REA".
- TTYIN - Input integer scalar used as unit number in FORTRAN READ statements directed to interactive terminal.
- TTYOUT - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.
- SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - GETNUM, FINDIT, HIERCH, STAN, SWITCH, CONDE, ELIM, SKLTN, TRNCLS, IO

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm EREDGE (Erase an edge on the minimum edge digraph). Given the number of elements currently in a reachability matrix, n, the

reachability matrix containing the minimum edge disgraph, R, and its associated index set vector, T, accept edges to be eliminated from interactive user.

- EREDGE1. [Accept edge to be eliminated.] Apply Algorithm GETNUM to set n1 and n2.
- EREDGE2. [See if termination is requested.] If $n1 = 0$ or $n2 = 0$, then algorithm is complete.
- EREDGE3. [Check to see if n1 and n2 are in system index set.] Apply Algorithm FINDIT. If n1 and/or n2 are/is not in system index set, then issue appropriate error message and then go to step EREDGE1.
- EREDGE4. [Check to see if n1 and n2 are members of a cycle.] If $r[ii, jj] = 1$ and $r[jj, ii] = 1$, issue error message and then go to step EREDGE1.
- EREDGE5. [Calculate non/redundant adjacency matrix for R.] Apply Algorithms HIERCH, STAN, CONDE, SKLTN.
- EREDGE6. [Check to see if n1 and n2 are on minimum edge disgraph.] Apply Algorithm FINDIT. If n1 and/or n2 are/is not on minimum edge disgraph, then issue appropriate error message and then go to step EREDGE1.
- EREDGE7. [Check to see if edge from n1 to n2 exists on minimum edge disgraph.] If $r[ii, jj] = 0$, then issue error message and then go to step EREDGE1.
- EREDGE8. [Obtain index positions of n1 and n2 on R.] Apply Algorithm FINDIT.
- EREDGE9. [Initialize.] $l \leftarrow 1$.
- EREDGE10. [Check to see if l is not a member of the antecedent set of n1.] If $r[l, ii] = 0$, then go to step EREDGE14.
- EREDGE11. [Initialize.] $k \leftarrow 1$.
- EREDGE12. [Check to see if k is not a member of the reachability set of n2.] If $r[jj, k] = 0$, then go to step EREDGE13, else set $r[l, k] \leftarrow 0$.
- EREDGE13. [Loop on k.] $k \leftarrow k + 1$. If $k \leq n$, then go to step EREDGE12.

EREDGE14. [Loop on l.] $l \leftarrow l + 1$. If $l \leq n$, then go to step EREDGE10.

EREDGE15. [Calculate reachability.] Apply Algorithm TRNCLS.

EREDGE16. [Update permanent file.] Apply Algorithm IO and then go to step EREDGE1.

ISMS SUBROUTINE/FUNCTION NAME - AEDGE

FUNCTION - Adds edges on the minimum edge disgraph.

USAGE - CALL AEDGE(N, REA, INDEX, TTYIN, TTYOUT,
SYS)

PARAMETERS

| | |
|--------|--------------------------------------------------------------------------------------------------------------------------------------------------|
| N | - Input/Output integer scalar indicating the number of elements currently on "REA". |
| REA | - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS". This is the reachability matrix containing the minimum edge disgraph. |
| INDEX | - Input integer vector of length "SYS" containing the index set of "REA". |
| TTYIN | - Input integer scalar used as unit number in FORTRAN READ statements directed to interactive terminal. |
| TTYOUT | - Input integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal. |
| SYS | - Input integer scalar used in FORTRAN DIMENSION statements. |

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - GETNUM, FINDIT, TRNCLS

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm AEDGE (Add an edge on the minimum edge disgraph). Given the number of elements currently in a reachability matrix, n, the reachability matrix containing the minimum edge disgraph, R, and its

associated index set, T , ask the interactive user for edges (relationships) to be added on the minimum edge digraph.

- AEDGE1. [Accept two element numbers from interactive terminal.] Apply Algorithm GETNUM to obtain $n1$ and $n2$.
- AEDGE2. [Check for termination directive.] If $n1 = 0$ or $n2 = 0$, then go to step AEDGE5.
- AEDGE3. [Check to see if $n1$ and $n2$ are members of index set.] Apply Algorithm FINDIT. If $n1$ and/or $n2$ are/is not in index set, write error message on interactive terminal and then go to step AEDGE1.
- AEDGE4. [Put edge in.] $r[i,j] \leftarrow 1$, then go to step AEDGE1.
- AEDGE5. [Transitively close matrix.] Apply Algorithm TRNCLS and then algorithm is complete.

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ERASE

ISMS SUBROUTINE/FUNCTION NAME - ERASE

FUNCTION - Erases elements from a binary matrix.

USAGE - CALL ERASE(N,MATRIX,INDEX,TTYIN,TTYOUT,
SYS)

PARAMETERS N - Input/Output integer scalar indicating
the number of elements currently in
"MATRIX".

MATRIX - Input/Output logical two dimensional
matrix of dimensions "SYS" X "SYS".
This is the argument binary matrix.

IDEX - Input/Output integer vector of length
"SYS" containing the index set of
"MATRIX".

TTYIN - Input integer scalar used as unit number
in FORTRAN READ statements directed to
interactive terminal.

TTYOUT - Input integer scalar used as unit number
in FORTRAN WRITE statements directed
to interactive terminal.

SYS - Input integer scalar used in FORTRAN
DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - ELIM,GETNUM,FINDIT

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described
below:

Algorithm ERASE (Erase elements from a binary matrix). Given the
number of elements currently in a binary matrix, n, the binary

matrix, R, and its associated index set vector, T, accept and erase elements numbers from interactive terminal and erase them.

- ERASE1. [Check to make sure n is not .LE. 0.] If n .LE. 0, then write out error message on interactive terminal and algorithm is complete.
- ERASE2. [Accept an element number from interactive terminal.] Apply Algorithm GETNUM to obtain n1. If n1 = 0, then algorithm is complete.
- ERASE3. [Check to see if n1 is a member of index set.] Apply Algorithm FINDIT. If n1 is not a member, issue appropriate error message and then go to step ERASE2.
- ERASE4. [Erase element from matrix.] Apply Algorithm ELIM and then go to step ERASE2.

ISMS SUBROUTINE/FUNCTION NAME - TRNCLS

FUNCTION - Transitively closes a binary matrix to yield a reachability matrix.

USAGE - CALL TRNCLS(N,MATRIX,SYS)

PARAMETERS N - Input integer scalar indicating number of elements in "MATRIX".

MATRIX - Input/Output logical two dimensional matrix of dimensions "SYS" X "SYS". this is the argument binary matrix.

SYS - Input integer scalar used in FORTRAN DIMENSION statements.

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm TRNCLS (Transitive closure). Given the number of elements currently in a matrix, n, and the binary matrix, M, calculate and rearrange M into a reachability matrix.

TRNCLS1. [Initialize.] i \leftarrow 1.

TRNCLS2. [Initialize element #i search.] nones \leftarrow 0, last \leftarrow 0.

TRNCLS3. [Initialize.] k \leftarrow 1.

TRNCLS4. [Find all ones in reachability set of element #i.] If $m[i,k] = 0$, then go to step TRNCLS6.

TRNCLS5. [Found a one, keep track of it.] nones \leftarrow nones + 1, fill local vector C, $c[\text{nones}] \leftarrow k$.

TRNCLS6. [Loop on k.] $k \leftarrow k + 1$. If $k \leq n$, then go to step TRNCLS4.

TRNCLS7. [Check to see if any new ones were added from last time through.] If $\text{nones} = \text{last}$, then go to step TRNCLS14.

TRNCLS8. [No, compute new elements in reachability set of element #I by transitivity.] $\text{last} \leftarrow \text{nones}$, $l \leftarrow 1$.

TRNCLS9. [Initialize.] $k \leftarrow c[l]$, $j \leftarrow 1$.

TRNCLS10. [Fill up matrix.] If $m[i, k] = 1$ and $m[k, j] = 1$, then set $m[i, j] \leftarrow 1$.

TRNCLS11. [Loop on j.] $j \leftarrow j + 1$. If $j \leq n$, then go to step TRNCLS10.

TRNCLS12. [Loop on l.] $l \leftarrow l + 1$. If $l \leq \text{nones}$, then go to step TRNCLS9.

TRNCLS13. [Continue processing.] $\text{nones} \leftarrow 0$, then go to step TRNCLS3.

TRNCLS14. [Loop on i.] $i \leftarrow i + 1$. If $i \leq n$, then go to step TRNCLS2, else algorithm is complete.

ISMS SUBROUTINE/FUNCTION NAME - PACK

FUNCTION - Packs as many display characters as possible into one machine word.

USAGE - CALL PACK(NWORDS,NCHAR,CARD)

PARAMETERS,

- NWORDS - Input integer scalar equal to the number of machine words required to hold 600 characters plus 10 additional words.
- NCHAR - Input integer scalar equal to the number of characters able to be stored in one machine word.
- CARD - Input integer vector of length 60 containing the text to be packed, stored one character per word.

COMMON BLOCKS - Blank common is utilized.

COMMON PARAMETERS - See the description for subroutine MAKEIT.

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm PK (Packing Algorithm). Given a maximum of 60 characters stored in one word, pack as many characters as possible into one word and store the packed string starting on the first available word boundary.

PK1. [Initialize.] pos \leftarrow pos + 1, i \leftarrow 1

PK2. [Compute number of characters to process.] If card(61 - i) is not a blank, then go to step PK4.

PK3. [Loop on i.] i \leftarrow i + 1. If i \leq 61, then go to step PK2, else i = 59.

PK4. [Calculate actual length.] $len \leftarrow 61 - i$

PK5. [Encode text (i.e., PACK).] $length(pos) \leftarrow (len - 1) / nchar + 1$,
 $end \leftarrow start + length(pos)$. Store reformatted text into
"packed", $packed(i)$, $i=start, end$. $start \leftarrow start + length(pos)$
and algorithm is complete.

ISMS SUBROUTINE/FUNCTION NAME - WRITE

FUNCTION - Writes a record onto the direct access query file.

USAGE - CALL WRITE(LOGPOS,NWORDS)

PARAMETERS
 LOGPOS - Input integer scalar denoting the logical position of the record in the file.
 NWORDS - Input integer scalar equal to the number of machine words required to hold 600 characters plus 10 additional words.

COMMON BLOCKS - Blank common is utilized.

COMMON PARAMETERS - See the description for subroutine MAKEIT.

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm WR (Write Algorithm). Given a record of size "nwords", write the record onto the query file in the proper position.

WR1. [Determine if an element text is to be written.] If logpos < 5, then go to step WR3.

WR2. [An element text is to be written.] $n \leftarrow n + 1$

WR3. [Write record.] Write record onto file in position "logpos".

WR4. [Reset record associated variables.] start $\leftarrow 1$, pos $\leftarrow 0$, length(1) $\leftarrow 0$, $i=1,2,\dots,10$ and algorithm is complete.

MAKEIT

ISMS SUBROUTINE/FUNCTION NAME - MAKEIT

FUNCTION - Converts a sequential access file containing queries for an ISM session into a direct access file for random retrieval.

USAGE - CALL MAKEIT

PARAMETERS - None

COMMON BLOCKS - Blank common is utilized.

COMMON PARAMETERS

TTYIN - Integer scalar used as unit number in FORTRAN READ statements directed to interactive terminal.

TTYOUT - Integer scalar used as unit number in FORTRAN WRITE statements directed to interactive terminal.

N - Integer scalar equal to the number of element text records processed.

TOTAL - Integer scalar equal to the number of machine words required for the current element text record.

POS - Integer scalar which denotes the line number (1 thru 10) of the current element text record being processed.

START - Integer scalar which is equal to the first available subscript of "PACKED" for packing in the current element text record.

LENGTH - Integer vector of length 10 words.
LENGTH(I) = number of machine words required to store line#I of the element text record currently being processed.

PACKED - Integer vector of length 200 words containing the text of the element text record with each line stored on a word boundary.

REQUIRED ISMS ROUTINES - DOIT, SHOW, PACK, WRITE, GETNUM

REQUIRED FORTRAN ROUTINES - ENCODE (if available)

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm MK (Makeit). Given a line oriented sequential text file with textual and control information, reformat and write the information to a direct access file.

MK1. [Initialize.] nchar \leftarrow number of bits per word/number of bits required to represent one character, nwords \leftarrow (60/nchar)*10

MK2. [Construct file.] Apply Algorithm DOIT.

MK3. [Write number of records on first record.] Write, variable n.

MK4. [Ask user if he wants to show elements.] Write message and read response. If user does not want to see his elements then go to step MK6.

MK5. [Show elements.] Apply Algorithm SHOW.

MK6. [Terminate.] Algorithm is complete.

DOIT

ISMS SUBROUTINE/FUNCTION NAME - DOIT

FUNCTION - Creates the random text file.

USAGE - CALL DOIT(NCHAR,NWORDS,CARD)

PARAMETERS

| | |
|--------|-----------------------------------------------------------------------------------------------------------------------|
| NCHAR | - Input integer scalar equal to the number of characters able to be stored in one machine word. |
| NWORDS | - Input integer scalar equal to the number of machine words required to hold 600 characters plus 10 additional words. |
| CARD | - Input integer vector of length 60 words. |

COMMON BLOCKS - Blank common is utilized.

COMMON PARAMETERS - See the description for subroutine MAKEIT.

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm DO (Doit Algorithm). Given a sequential text file with text and control cards, read each record and form a random access text file for full text queries.

DO1. [Initialize.] start \leftarrow 1, flag \leftarrow 0, n \leftarrow 0, length(i) \leftarrow 0, i=1,2,...,10.

DO2. [Read a record from text file.] Read in sixty characters into "card" with each character on a word boundary. On end of file, algorithm is complete.

DO3. [Check for control card.] If CARD(1) = "/" then go to step DO5.

DO4. [Pack text.] The record read was a text record so apply Algorithm PK and then go to step DO2.

- D05. [Check to see if current record should be written.] If flag = 0 then go to step D07.
- D06. [Write current record to random text file.] Apply Algorithm WRITE.
- D07. [Parse control card.] $logpos \leftarrow 0$. If CARD(2) not equal to an "R" then go to step D010.
- D08. [Determine logical position for next record.] If CARD(3) = i, then $logpos \leftarrow i + 1$, $i=1,2,3$.
- D09. [Check for error on control card.] If $logpos < 2$, then go to step D013, else flag = 1 and go to step D02.
- D010. [Continue parsing control card.] If CARD(2) is not equal to an "E", then go to step D012.
- D011. [Continue parsing control card.] If CARD(3) = "L", then $logpos \leftarrow n + 5$. If $logpos$ is less than 2, then go to step D013, else set flag = 1 and go to step D02.
- D012. [Check for end of file.] If CARD(2) = "/", then algorithm is complete.
- D013. [Write error on interactive terminal.] Write "control card error" and then algorithm is complete.

SHOW

ISMS SUBROUTINE/FUNCTION NAME - SHOW

FUNCTION - Displays all or user specified pairs of elements as they might appear during an ISM session.

USAGE - CALL SHOW

COMMON BLOCKS SHOWB - Named integer common block of length 5*"NWORDS."
- blank common is also utilized. See the description for subroutine MAKEIT.

COMMON PARAMETERS - See the description for subroutine BORDER for the list of parameters.
- SHOWB has been made a common block to guarantee the contingency of storage.

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm SH (Show Algorithm). Given a direct access query file, display user specified sets of queries.

SH1. [Initialize.] $a11s \leftarrow 0$, $e11 \leftarrow 0$, $e12 \leftarrow 0$

SH2. [Read in relational clauses 1,2, and 3.] Read R1, R2, R3.

SH3. [See if user wants to display all elements.] Ask user, if yes, then go to step SH13.

SH4. [Ask user which two elements to present.] Ask user, if one or two zeros are typed, then algorithm is complete. Otherwise, make sure that elements text exist. If not, then issue an error message and then go to step SH4.

SH5. [Read in elements text.] $i1 \leftarrow e11 + 4$, $i2 \leftarrow e12 + 4$, read in records $i1$ and $i2$.

SH6. [Initialize for typing out.] offset \leftarrow 0, i \leftarrow 1

SH7. [Initialize for ten phrases.] J \leftarrow 1, i1 \leftarrow 0, i2 \leftarrow 0

SH8. [Check length indicator.] If SHOWB(J + offset) = 0, then go to step SH10.

SH9. [Print out this phrase.] length \leftarrow SHOWB(J + OFFSET), i1 \leftarrow i2 + 1, i2 \leftarrow i1 + LENGTH + 1, WRITE OUT SHOWB(c + offset + 10), c=i1, ..., i2.

SH10. [Loop on J.] J \leftarrow J + 1. If J < 11, then go to step SH8.

SH11. [Loop on i.] offset \leftarrow offset + nwords, i \leftarrow i + 1. If i < 6, then go to step SH7.

SH12. [Check for desired printing of all element pairs.] If alls = 0, then go to step SH4.

SH13. [Come here when printing all element pairs.] e11 \leftarrow e12 + 1, e12 \leftarrow e12 + 1. If e11 > n, then algorithm is complete, if e12 > n, then set e12 \leftarrow 1, alls = 1, then go to step SH5.

ADD

ISMS SUBROUTINE/FUNCTION NAME - ADD

FUNCTION - Adds elements to the weighted matrix.

USAGE - CALL ADD(N,MAT,INDEX)

PARAMETERS

| | |
|-------|---------------------------------------------------------------------------------------------------|
| N | - Input/Output integer scalar indicating the current number of elements in the weighted matrix. |
| MAT | - Input/Output integer two dimensional matrix of dimensions 50 X 50. This is the weighted matrix. |
| INDEX | - Input/Output integer vector of length 50 words containing the index set for "MAT". |

COMMON BLOCKS INFO - Named integer common block 3 words long.

COMMON PARAMETERS - See the description for subroutine FRNTMT.

REQUIRED ISMS ROUTINES - FINDIT,GETNUM,QUEST

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm ADD (Add an element to a weighted matrix). Given the number of elements currently in a weighted matrix, n , the weighted matrix, W , and its associated index set, T , add an element to the weighted matrix.

ADD1. [Ask user for new element to be added.] Prompt user and then apply algorithm GETNUM.

ADD2. [Make sure that number is less than or equal to 9999 for fortran formatting.] If number read in is > 9999 , issue error message and then go to step ADD1.

ADD3. [Check for a zero input.] If number typed was a zero, then algorithm is complete.

ADD4. [Check for duplicate element.] Apply Algorithm FINDIT. If element is already in index set, the issue error message and then go to step ADD1.

ADD5. [Put new element into matrix.] $n \leftarrow n + 1$, $t[n] \leftarrow$ new number.

ADD6. [See if user wants to fill up relationships for new element.] Prompt user. If user types an "N", then go to step ADD1.

ADD7. [Initialize column fill.] $i \leftarrow 1$

ADD8. [Present question.] Apply Algorithm QUEST($t[i]$, $t[n]$).

ADD9. [Get weight value.] Apply Algorithm GETNUM.

ADD10. [Make sure weight value is less than 10.] If weight value is greater than 9, issue error message and then go to step ADD8.

ADD11. [Set matrix.] $w[i,n] \leftarrow$ weight value

ADD12. [Loop on i.] $i \leftarrow i + 1$. If i is less than or equal to $n - 1$, then go to step ADD8.

ADD13. [Initialize row fill.] $i \leftarrow 1$

ADD14. [Present question.] Apply Algorithm QUEST($t[n]$, $t[i]$)

ADD15. [Get weight value.] Apply Algorithm GETNUM.

ADD16. [Make sure weight value is less than 10.] If weight value is greater than 9, issue error message and then go to step ADD14.

ADD17. [Set matrix.] $w[n,i] \leftarrow$ weight value

ADD18. [Loop on i.] $i \leftarrow i + 1$. If i is less than or equal to $n - 1$, then go to step ADD14, else go to step ADD1.

ESNIMI

ISMS SUBROUTINE/FUNCTION NAME - PRNTMT

FUNCTION

- Prints out all weighted relationships included in a weighted matrix that are greater than or equal to an input weight value.

USAGE

- CALL PRNTMT(N,MAT,INDEX,THRESH,SELPNT)

PARAMETERS

- N - Input integer scalar equal to the number of elements on "MAT".
- MAT - Input integer two dimensional matrix of dimensions 50 X 50. This is the weighted matrix.
- INDEX - Input integer vector of length 50 containing the index set for "MAT".
- THRESH - Input integer scalar used as cut-off threshold. All relationships greater than or equal to "THRESH" will be printed if "SELPNT" = .TRUE..
- SELPNT - Input logical variable used to determine if only relationships greater than or equal to "THRESH" should be printed (SELPNT = .TRUE.) or all relationships (SELPNT = .FALSE.).

COMMON BLOCKS,

- INFO - Named integer common block 3 words long.

COMMON PARAMETERS

- QTYPE - Input logical variable indicating the type of query to be presented.
QTYPE = .FALSE. causes full text queries to be printed.
QTYPE = .TRUE. causes symbolic queries to be printed.
- TTYIN - Input integer scalar used as unit number in FORTRAN READ statements directed to interactive terminal.
- TTYOUT - Input integer scalar used as unit number in in FORTRAN WRITE statements directed to interactive terminal.

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm PT (Print weighted matrix). Given the number of elements in a weighted matrix, n , the weighted matrix, W , and its associated index set vector, M , print out all or just relationships greater than or equal to a threshold, t .

PT1. [Initialize.] If $selent = 1$, then print special title. Set $i, i \leftarrow 1$.

PT2. [Begin row search.] Set $ctr \leftarrow 0, j \leftarrow 1$.

PT3. [Check for printing greater than or equal to threshold.] If $selent = 0$, then go to step PT5.

PT4. [Check for relationships greater than or equal to the threshold.] If $w[i,j] \geq t$, then go to step PT7.

PT5. [Don't print diagonal.] If $i = j$, then go to step PT7.

PT6. [Store values in local array for printing.] $ctr \leftarrow ctr + 1$, $list[ctr,1] \leftarrow m[j]$, $list[ctr,2] \leftarrow w[i,j]$.

PT7. [Loop on j .] $j \leftarrow j + 1$. If $j \leq n$, then go to step PT3.

PT8. [Check for $ctr=0$.] If $ctr = 0$, then write out $m[i]$ and then go to step PT10.

PT9. [Write out row relationships.] Write out $m[i]$, $list[k,1]$, $list[k,2]$, $k=1,2,\dots,ctr$.

PT10. [Loop on i .] $i \leftarrow i + 1$. If $i \leq n$, then go to step PT2, else algorithm is complete.

CHANGE

ISMS SUBROUTINE/FUNCTION NAME - CHANGE

FUNCTION - Changes weights in the weighted matrix via user interaction.

USAGE - CALL CHANGE(N,MAT,INDEX)

PARAMETERS

N - Input integer scalar indicating the current number of elements in "MAT".
MAT - Input/Output integer two dimensional matrix of dimensions 50 X 50 containing the weighted matrix.
INDEX - Input integer vector of length 50 words containing the index set for "MAT".

COMMON BLOCKS

INFO - Named integer common block of length 3 words.

COMMON PARAMETERS

- See the description for subroutine PRNTMT.

REQUIRED ISMS ROUTINES

- FINDIT,GETNUM

REQUIRED FORTRAN ROUTINES

- None

ALGORITHM EMPLOYED

- The algorithm employed is described below:

Algorithm CH (Change Algorithm). Given the number of elements currently in a weighted matrix, n, the weighted matrix, W, and its associated index set, T, allow changes in the matrix via user interaction.

CH1. [Prompt user.] Write message to interactive terminal.

CH2. [Read change.] Apply Algorithm GETNUM. Read the element numbers and weight change (3 numbers).

CH3. [Check for termination.] If any of the element numbers are zero, algorithm is complete.

- CH4. [Check to see if element numbers are valid.] Apply Algorithm FINDIT. If one or both numbers are invalid, write an error message and then go to step CH1.
- CH5. [Make sure weight typed is 9 or less.] If weight value is greater than 9, write error message and then go to step CH1.
- CH6. [Change weight.] $w[x,y] \leftarrow \text{weight}$, where x and y are outputs of algorithm FINDIT. Go to step CH1.

EILL

ISMS SUBROUTINE/FUNCTION NAME - FILL

FUNCTION - Helps the user fill the weighted matrix by presenting the required queries.

USAGE - CALL FILL(N,MAT,INDEX,R1,R2)

PARAMETERS

| | |
|-------|----------------------------------------------------------------------------------------------------|
| N | - Input integer scalar indicating the current number of elements in "MAT". |
| MAT | - Input/Output integer two dimensional matrix of dimensions 50 X 50p. This is the weighted matrix. |
| INDEX | - Input integer vector of length 50 words containing the index set for "MAT". |
| R1 | - Input/Output integer scalar used as the row restart indice. |
| R2 | - Input/Output integer scalar used as the column restart indice. |

COMMON BLOCKS

| | |
|------|-------------------------------------------------|
| INFO | - Named integer common block of length 3 words. |
|------|-------------------------------------------------|

COMMON PARAMETERS - See the description for subroutine PRNTMT.

REQUIRED ISMS ROUTINES - QUEST,GETNUM

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm FL (Fill Algorithm). Given the number of elements currently in a weighted matrix, n, the weighted matrix, W, its associated index set, T, and two restart parameters r1 and r2, allow the filling of all positions (with the exception to the diagonal) on the weighted matrix.

FL1. [Check to see if this is a restart.] If r1 or r2 is greater than zero, go to step FL12.

FL2. [Initialize.] row \leftarrow 1, col \leftarrow 1, r1 \leftarrow 0, r2 \leftarrow 0

FL3. [Initialize row loop.] i \leftarrow row

FL4. [Initialize column loop.] j \leftarrow col

FL5. [Don't process diagonal.] If i = j, then go to step FL10.

FL6. [Present query.] Apply Algorithm QUEST(t[i],t[j]).

FL7. [Obtain weight.] Apply Algorithm GETNUM.

FL8. [Check for valid weight input.] If weight is equal to 10, then go to step FL19. If weight is greater than 9, then issue error message and go to step FL6.

FL9. [Set matrix.] w[i,j] \leftarrow weight value

FL10. [Loop on j.] j \leftarrow j + 1. If j is greater than or equal to n, then go to step FL5.

FL11. [Loop on i.] i \leftarrow i + 1. If i is less than or equal to n, then go to step FL4, else algorithm is complete.

FL12. [Restart questioning.] i \leftarrow r1, j \leftarrow r2

FL13. [Present query.] Apply Algorithm QUEST(t[i],t[j]).

FL14. [Obtain weight.] Apply Algorithm GETNUM.

FL15. [Check for valid weight input.] If weight is equal to 10, then go to step FL19. If weight is greater than 9, then issue error message and go to step FL13.

FL16. [Set matrix.] w[i,j] \leftarrow weight value

FL17. [Loop on j.] j \leftarrow j + 1. If j is less than or equal to n, then go to step FL13.

FL18. [Reset restart parameters.] row \leftarrow 1, col \leftarrow r1 + 1, r1 \leftarrow 0, r2 \leftarrow 0, then go to step FL3.

FL19. [Generate restart parameters.] r1 \leftarrow i, r2 \leftarrow j, then algorithm is complete.

RESOLV

ISMS SUBROUTINE/FUNCTION NAME - RESOLV

FUNCTION - Resolves the threshold of a weighted matrix.

USAGE - CALL RESOLV(N,MAT,INDEX,THRESH)

PARAMETERS

| | |
|--------|---------------------------------------------------------------------------------------------|
| N | - Input integer scalar indicating the current number of elements in "MAT". |
| MAT | - Input integer two dimensional matrix of dimension's 50 X 50. This is the weighted matrix. |
| INDEX | - Input integer vector of length 50 words containing the index set for "MAT". |
| THRESH | - Input integer scalar equal to the maximum threshold to be used (usually equal to 9). |

COMMON BLOCKS

| | |
|------|-------------------------------------------------|
| INFO | - Named integer common block of length 3 words. |
|------|-------------------------------------------------|

COMMON PARAMETERS - See the description for subroutine PRNTMT.

REQUIRED ISMS ROUTINES - TRNCLS, PRNTMT, GEOD, NOTR

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm R5 (Resolve Algorithm). Given the number of elements currently in a weighted matrix, n , the weighted matrix, W , its associated index set, T , and the maximum threshold, r , resolve the maximum threshold of the matrix.

R51. [Initialize.] $z \leftarrow n$

R52. [Initialize.] $i \leftarrow 1$

R53. [Begin constructing binary adjacency matrix.] $j \leftarrow 1$

- RS4. [Don't process main diagonal.] Set $a[i,j] \leftarrow 0$. If $i = j$, then go to step RS6.
- RS5. [Check threshold.] If $w[i,j]$ is less than z , then go to step RS7.
- RS6. [Set binary matrix position.] $a[i,j] \leftarrow 1$
- RS7. [Loop on j.] $j \leftarrow j + 1$. If j is less than or equal to n , then go to step RS4.
- RS8. [Loop on i.] $i \leftarrow i + 1$. If i is less than or equal to n , then go to step RS3.
- RS9. [Transitively close matrix.] Apply Algorithm TRNCLS to set reachability matrix b .
- RS10. [Is binary matrix all ones?] If $b[i,j] = 0$, then go to step RS14, $j=1,2,\dots,n$, $i=1,2,\dots,n$.
- RS11. [Tell user cycle is resolved.] Write z .
- RS12. [Print universal matrix.] Apply Algorithm PRNTMT.
- RS13. [Print geodetic paths.] Apply Algorithm GEOD and then algorithm is complete.
- RS14. [Select next lowest threshold.] $z \leftarrow z - 1$, then go to step RS2.

GEOD

ISMS SUBROUTINE/FUNCTION NAME - GEOD

FUNCTION - Outputs the geodetic cycle paths contained within a weighted matrix.

USAGE - CALL GEOD(ADJ,N,INDEX)

PARAMETERS

| | |
|-------|------------------------------------------------------------------------------------------------------------------------------------------|
| ADJ | - Input logical two dimensional matrix of dimensions 50 X 50. This is the binary threshold adjacency matrix formed in subroutine RESOLV. |
| N | - Input integer scalar indicating the number of elements in "ADJ". |
| INDEX | - Input integer vector of length 50 containing the index set for "ADJ". |

COMMON BLOCKS

| | |
|------|-------------------------------------------------|
| INFO | - Named integer common block of length 3 words. |
|------|-------------------------------------------------|

COMMON PARAMETERS - See the description for subroutine PRNTMT.

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm GE (Geodetic Algorithm). Given the number of elements currently in a binary threshold matrix, n , the binary threshold matrix, A , and its associated index set vector, T , determine and print all geodetic cycle sets.

GE1. [Initialize.] $i \leftarrow 1$, $nbn \leftarrow 1$

GE2. [Initialize G1 formation.] $j \leftarrow 1$

GE3. [Check for a zero.] If $a[i,j] = 0$, then go to step GE6.

GE4. [Set G and B.] $a[i,j] \leftarrow 1$, $b[i,j] \leftarrow 1$

GE5. [Subtract identity.] If $i = j$, then $a[i,j] = 0$.

GE6. [Loop on j.] $j \leftarrow j + 1$. If j is less than or equal to n , then go to step GE3.

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GE7. [Loop on i.] $i \leftarrow i + 1$. If i is less than or equal to n , then go to step GE2.

GE8. [Initialize.] $i \leftarrow 1$

GE9. [Initialize.] $j \leftarrow 1$

GE10. [Zero C.] $c(j,i) \leftarrow 0$

GE11. [Initialize.] $k \leftarrow 1$

GE12. [Set switch.] $ma \leftarrow 0$. If $a[i,j] = 0$, then $ma \leftarrow 1$.

GE13. [Compute $A^{**}nbn-1$.] $c[j,i] \leftarrow ma * b[k,i] + c[j,i]$

GE14. [Loop on k.] $k \leftarrow k + 1$. If k is less than or equal to n , then go to step GE12.

GE15. [Check for invalid entry.] If $c[j,i]$ is not equal to 0, then set $c[j,i] \leftarrow 1$.

GE16. [Loop on j.] $j \leftarrow j + 1$. If j is less than or equal to n , then go to step GE10.

GE17. [Loop on i.] $i \leftarrow i + 1$. If i is less than or equal to n , then go to step GE9.

GE18. [Initialize for $G^{**}nbn$.] $nbn \leftarrow nbn + 1$, $i \leftarrow 1$

GE19. [Initialize.] $j \leftarrow 1$

GE20. [Calculate $G^{**}nbn$.] $g[i,j] \leftarrow g[i,j] + nbn * (c[i,j] - b[i,j])$, $b[i,j] \leftarrow c[i,j]$.

GE21. [Loop on j.] $j \leftarrow j + 1$. If j is less than or equal to n , then go to step GE20.

GE22. [Loop on i.] $i \leftarrow i + 1$. If i is less than or equal to n , then go to step GE19.

GE23. [Is G formed yet ?] If nbn is less than $n - 1$, then go to step GE8.

GE24. [Initialize path computation.] $i \leftarrow 2$, $ncpt \leftarrow 0$, write heading on terminal.

GE25. [Set limit for j.] $lim \leftarrow i - 1$, $j \leftarrow 1$

GE26. [Find cycle distance.] nb \leftarrow s[i,j], na \leftarrow s[j,i]

GE27. [If distance is zero, don't process.] If na or nb = 0, then go to step GE47.

GE28. [Keep account of paths.] nbnd \leftarrow 1, l[nbnd] \leftarrow i

GE29. [Is path longer than 1 ?] If nb is greater than 1, then go to step GE31.

GE30. [Store end link.] nbnd \leftarrow nbnd + 1, l[nbnd] \leftarrow j, then go to step GE36.

GE31. [Initialize search for last link.] lmm \leftarrow nb - 1, last \leftarrow 0, in \leftarrow 1

GE32. [Find link back to i.] nbnd \leftarrow nbnd + 1, l[nbnd] \leftarrow next link (Apply Algorithm NOTR), last \leftarrow l[nbnd].

GE33. [Unsuccessful ?] If ix = 1, then write out error message and go to step GE36.

GE34. [Loop on in.] in \leftarrow in + 1. If in is less than or equal to lmm, then go to step GE32.

GE35. [Store end link.] nbnd \leftarrow nbnd + 1, l[nbnd] \leftarrow j

GE36. [Is path complete ?] If na is greater than 1, then go to step GE38.

GE37. [Store return link.] l[nbnd + 1] \leftarrow i, then go to step GE43.

GE38. [Initialize search for last link.] lmn \leftarrow na - 1, last \leftarrow 0, in \leftarrow 1

GE39. [Find link back to j.] nbnd \leftarrow nbnd + 1, l[nbnd] \leftarrow next link (Apply Algorithm NOTR), last \leftarrow l[nbnd].

GE40. [Unsuccessful ?] If ix = 1, then write out error message and go to step GE43.

GE41. [Loop on in.] in \leftarrow in + 1. If in is less than or equal to lmn, then go to step GE39.

GE42. [Store end link.] l[nbnd + 1] \leftarrow i

GE43. [Initialize for printout.] ncpt \leftarrow ncpt + 1, lmc \leftarrow nbnd + 1, nx \leftarrow 1

- GE44. [Fix up print vector for index set used.] $\text{num} \leftarrow 1[\text{nx}]$, $\text{p}[\text{nx}] \leftarrow \text{t}[\text{num}]$
- GE45. [Loop on nx.] $\text{nx} \leftarrow \text{nx} + 1$. If nx is less than or equal to lmc , then go to step GE44.
- GE46. [Print out cycle path on terminal.] $\text{nxx} \leftarrow \text{na} + \text{nb}$, write ncpt , nxx , $\text{t}[\text{i}]$, $\text{t}[\text{j}]$, $\text{p}[\text{k}]$, $\text{k}=1,2,\dots,\text{lmc}$.
- GE47. [Loop on j.] $\text{j} \leftarrow \text{j} + 1$. If j is less than or equal to lim , then go to step GE26.
- GE48. [Loop on i.] $\text{i} \leftarrow \text{i} + 1$. If i is less than or equal to n, then go to step GE25, else algorithm is complete.

NOIR

ISMS SUBROUTINE/FUNCTION NAME - NOTR

FUNCTION - Returns the next element in the geodetic path.

USAGE - X = NOTR(G, ICL, NBLG, ILG, NBCL, IX, NBN, LAST)

| | | |
|------------|------|-------------------------------------------------------------------------------------------------------------------------------|
| PARAMETERS | G | - Input integer two dimensional matrix of dimensions 50 X 50. This is the distance matrix. |
| | ICL | - Input integer scalar equal to the starting search index. |
| | NBLG | - Input integer scalar equal to the distance of the link. |
| | ILG | - Input integer scalar equal to the ending search index. |
| | NBCL | - Input integer scalar equal to the distance of the desired link from the starting index. |
| | IX | - Output integer scalar used as an error switch for incomplete paths. IX = 0 means complete path. IX = 1 means no link found. |
| | LAST | - Input integer scalar equal to the last link found. This is required to keep the algorithm searching down the same path. |

COMMON BLOCKS - None

REQUIRED ISMS ROUTINES - None

REQUIRED FORTRAN ROUTINES - None

ALGORITHM EMPLOYED - The algorithm employed is described below:

Algorithm NOTR. Given the number of paths in a distance matrix, n, the distance matrix, G, and other information, determine the next element in the geodetic cycle.

NOTR1. [Initialize.] $ix \leftarrow 0$, $ibi \leftarrow 0$, $i \leftarrow 1$
 NOTR2. [Find all elements with distance $nbis$ from ils .] If $ils[i, ils]$ is not equal to $nbis$, then go to step NOTR4.
 NOTR3. [Keep track of elements with same distance.] $ibi \leftarrow ibi + 1$, $l[ibi] \leftarrow i$
 NOTR4. [Loop on i .] $i \leftarrow i + 1$. If i is less than or equal to n , then go to step NOTR2.
 NOTR5. [If none found, then algorithm is complete.] If $ibi = 0$, then algorithm is complete.
 NOTR6. [Initialize.] $i \leftarrow 1$
 NOTR7. [Search for an element of proper distance.] If $g[ic1, i]$ is not equal to $nbch$, then go to step NOTR13.
 NOTR8. [See if this is the first link.] If $last = 0$, then go to step NOTR10.
 NOTR9. [Keep going on same path.] If $g[last, i]$ is not equal to 1, then go to step NOTR13.
 NOTR10. [Initialize.] $k \leftarrow 1$
 NOTR11. [See if this i is of proper distance.] If $i = l[k]$, then go to step NOTR15.
 NOTR12. [Loop on k .] $k \leftarrow k + 1$. If k is less than or equal to ibi , then go to step NOTR11.
 NOTR13. [Loop on i .] $i \leftarrow i + 1$. If i is less than or equal to n , then go to step NOTR7.
 NOTR14. [No link found.] $ix \leftarrow 1$, then algorithm is complete.
 NOTR15. [Return link.] $notr \leftarrow l[k]$, then algorithm is complete.

Attachment No. 2

INTERPRETIVE STRUCTURAL MODELING SOFTWARE

Prepared by:

David R. Yingling, Jr.

This User's Manual was developed for the ISM software on the University of Dayton computer. It illustrates the type of instructions that should be given to users wishing to utilize the ISM software.

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PREFACE

This manual* is intended to serve as a user's guide for the Interpretive Structural Modeling Software Package developed by the Engineering and Public Policy Group at the University of Dayton, and as implemented on UD's Univac Series 70 computer operating under VS/9. This version of the manual was written under the assumption that its user is familiar with the ISM** methodology and supportive concepts but has had little or no previous experience on a computer, specifically the UD computer. Therefore, all basic procedures like LOGGING ON, etc. have been included so as to make this manual a stand alone guide.

It is suggested that if the user plans to make extensive use of text files that he or she obtain Manual #4.1 (EDT) from the University of Dayton Office for Computing Activities (OCA). This manual is available free of charge and will make using the Univac File Editor (EDT) and text file editing much easier.

If you find any discrepancies in this manual or ISMS-UD, please write to the address on the front page or call Monday thru Friday 0800 to 1500 Eastern time. If the telephone is busy or no answer, "leave word" at and we will return your call.

*The table of contents may be found at the end.

**Interpretive Structural Modeling (ISM) is developed in John N. Warfield's Societal Systems: Planning, Policy, and Complexity, Wiley-Interscience: New York, 1976.

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INTRODUCTION

The University of Dayton ISMS Version 2.0, as implemented on UD's time sharing computer system, consists of three FORTRAN IV programs which perform embedding and amending, element text processing, and cycle resolution. The following is a brief description of each program.

ISMS-UD - is the program that provides embedding and amending facilities for Interpretive Structural Models. ISMS-UD maintains a reachability matrix permanent file. A query file supplies the information necessary for full text queries.

CYCLE - is the program that resolves the cycles contained in an Interpretive Structural Model. CYCLE creates and maintains a matrix containing numeric weights denoting the strength of relationships between elements of a cycle. A query file supplies the information necessary for full text queries.

MAKEII - is a program that restructures an EDT-created text file into a random access format required by the ISM method. The random access format is written onto the query file.

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UD_TIME_SHARING_COMPUTER_SYSTEM

SYSTEM_HARDWARE

The UD Computing Facility owns a Univac Series 70/7 Time-Sharing Computer System. This third generation virtual memory multiprogramming system has 750K bytes of main core storage with a backing store of 6.4 million bytes. Time-sharing input/output to the processor is handled by a communications controller attached to one of the multiplexor channels. The communications controller presently allows for 18 remote hardwire terminals and 33 remote dial-up terminals. Ten of the 33 dial-up ports are 300 baud with the remaining ports 110 baud.

SYSTEM_SOFTWARE

The Univac Series 70/7 computer presently runs under the operating system VS/9 Version 3.5. VS/9 is a group of programs and subprograms which control input, compilation, assembly, loading, execution, and output of all programs submitted to the computer as well as the allocation of system resources.

UNIVAC_FILE_EDITOR (EDT)

If the user desires to have programs ISMS-UD and CYCLE present queries in an English text format, a two step procedure must be followed. Step 1: a sequential element text file must be created (using EDT). Step 2: the sequential element text file must then be converted into a query file (using MAKEIT).

EDT is a program that permits the creation and modification of sequential element text files. EDT is invoked under VS/9 and responds to simple commands prefixed with an @ ("at sign") which initiate, maintain, correct, and complete file construction. Although EDT is a comprehensive editing package capable of performing varied tasks, it presents minimal concern to the ISMS-UD user. A sample run using EDT can be seen in the Appendix of this document. For a more comprehensive discussion of EDT, the user is referred to OCA Manual #4.1 (EDT) available free of charge at the UD data center, Miriam Hall.

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PROCEDURES AND IMPORTANT SYSTEM CONCEPTSSCHEDULING THE USE OF THE UD COMPUTER

The Univac Series 70/7 is a multiprogramming computer system; that is, it is servicing many users all at the same time. Consequently, the more people using the system, the slower the turnaround time (response time). It is suggested that the user not schedule the use of the computer during prime time (i.e., 1 P.M. to 5 P.M. weekdays) and near the end of the UD academic terms (usually first two weeks in both December and April). It has also been experienced near the end of the terms that dial-up ports are extremely hard to get, so these words to the wise --PLAN AHEAD-- START EARLY!!!

AUTOMATIC LOGOFF FEATURE

In order to keep dial-up ports from being tied up or occupied by an inactive terminal user, the UD computer has attached to it a device which will AUTOMATICALLY LOGOFF A TERMINAL INACTIVE FOR NINE MINUTES. This is important to remember because sometimes the queries presented by ISMS-UD require a long period of time for thought and discussion. During this discussion time, the computer under control of ISMS-UD is waiting for an answer to the query; that is, the terminal is inactive. If an input/output is not done within nine minutes after the start of the read, the terminal is automatically logged off and the ISM is partially lost. It is suggested that the terminal manager or user keep an eye on the time and, if the nine minute limit draws near while discussion continues, type an invalid input to ISMS-UD. ISMS-UD will inform the user of the invalid input and re-prompt with the same query. The nine minute timer is then reset and the ISM saved.

SPECIAL CONTROL KEYSEnd_of_Transmission_Key

Anyone who has used an interactive computer terminal knows that some special key on the terminal is used to signal the computer to take action on the typed input. On most computers, this is the RETURN key. The UD computer does not use the RETURN key, but instead, the depressing

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of two keys simultaneously. These two keys are the CTRL and C keys, denoted as CTRL C. That is, when the user wants to type a command to the computer, he or she first types the command and then sends the input to the computer by holding down the CTRL key and typing a "C". It will be assumed from here on that the user understands this principle.

Error-Correcting Keys

Two input error correcting keys are supported by the Univac hardware. They are:

"CTRL X" -- cancels the input line. The Processor responds with %CNCL to indicate the line was canceled.

"SHIFT O" or "UNDERLINE" -- is essentially a backspace key. When typed, it tells the processor to ignore the last typed character before the SHIFT O or UNDERLINE. (The type of terminal used dictates which is appropriate.)

Both error correcting keys may be used while running the operating system and ISMS-UD.

PARITY AND DUPLEX SETTINGS FOR DIAL-UP OPERATION

The parity setting on the terminal must be EVEN.

The duplex setting on the acoustic coupler (and terminal if applicable) must be HALF.

PERTINENT TELEPHONE NUMBERS

The telephone numbers below should be kept handy.

- 110 BAUD
- 110 BAUD
- 300 BAUD
- SYSTEM STATUS/SCHEDULE (recording)
- HELP!!!!

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HOW_TO_LOGON

There are two types of terminals that can be used on the UD computer: the hardwired terminal which is directly connected to the computer, and the acoustically coupled terminal which is connected to the computer via the telephone system. Each type of terminal requires a different LOGON procedure which will subsequently be discussed.

PROCEDURE FOR LOGGING ON A HARDWIRED TERMINAL

1. Turn on the power switch on the terminal.
2. Set the DUPLEX switch to HALF.
3. Set the LOCAL/LINE switch to LINE.

Note: On a Teletype Model 33, turn the control knob (below and to the right of the keyboard) counterclockwise instead of the above three steps.

4. While depressing the CTRL key, type a "C". The terminal should type:

```
%E222 PLEASE LOGON
/
```

5. The user should then type

```
/LOGON userid#,,C'password'
```

followed by a "CTRL C". The computer will then type various information about the task and return with a slash (/).

PROCEDURE FOR LOGGING ON AN ACOUSTICALLY COUPLED TERMINAL

1. Turn on the acoustic coupler's on/off switch if using a terminal with an external coupler (this switch is usually unmarked).

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2. Turn on the power switch on the terminal.
3. Set the PARITY switch on the terminal to EVEN.
4. Set the DUPLEX switch on the coupler and/or terminal to HALF.
5. Set the LOCAL/LINE switch on the coupler and/or terminal to LINE.

Note: On a Teletype Model 33, turn the control knob (below and to the right of the keyboard) counterclockwise instead of steps 2 and 5 above.

6. Set the BAUD switch to the desired position (either 110 or 300 baud -- sometimes denoted as 10 characters/sec. or 30 characters/sec.) on the terminal if so equipped.
7. Dial the computer's telephone number for the baud rate you wish to use (see "Pertinent Phone numbers"...remember to dial 9 first if using a phone on campus). The phone should ring twice and the computer will answer with a high pitched tone.
8. Upon hearing the tone, immediately place the phone into the cradle. There should be some directive on the coupler indicating at which end the phone cord should be.
9. The CARRIER or SIGNAL light on the coupler will be lit and the terminal will print:

```
%E222 PLEASE LOGON
```

```
/
```

10. The user should then type

```
/LOGON userid#,,C'password'
```

followed by a "CTRL C". The computer will then type various information about the task and return with a slash (/).

Provided that the user types a syntactically correct LOGON command containing a valid userid# and password, the system should be ready and waiting.

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LOGON PROBLEMS

If the computer doesn't answer, it is probably not operational, and you may obtain a recorded status/schedule report (see "pertinent telephone numbers").

If a busx signal is received, this indicates all available lines are in use. Try again in a few minutes.

If the computer answers, prints nothing, and quickly hangs up, the parity setting is usually incorrect or the phone was incorrectly placed onto the coupler.

If the computer responds normally but the terminal does not print what you type, the DUPLEX setting is incorrect.

The above are some of the more common problems with the LOGON procedure. There are many variations of the above problems which have not been discussed. If your LOGON problem is temporary, that is, you successfully LOGged ON yesterday but have not been able today, feel free to call the UD Office for Computing Activities HELP line (see "pertinent telephone numbers") for help. You could be informing them of a system error. If you have never been able to successfully LOGON, please call us, the Engineering and Public Policy Group at _____ or _____, and we will be happy to discuss your problem.

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ISMS-UD_FULL_TEXT_QUERY_FACILITY

If the user desires to have programs ISMS-UD and CYCLE present queries in an English text format, two things are necessary. First, a sequential element text file must be created using the Univac file editor EDT. The sequential element text file consists of control information, and the English text for the relational expression and for each element to be considered.

Second, the sequential element text file must be converted into a random access format required by the ISM method. Program MAKEIT performs this function on the sequential element text file and writes the random access format onto a query file.

FORMAT OF THE INPUT SEQUENTIAL ELEMENT TEXT FILE

The format of an example input sequential text file can be seen in its entirety in the Appendix. Certain slash (/) keywords are used to identify the records to be used for the five outputs typed on the terminal for each question when full text queries are used. The five outputs are:

- 1) introductory clause
- 2) element a
- 3) relational clause
- 4) element b
- 5) qualifying clause

An example of the above is:

```
DOES  
DEVELOPING SOCIAL INCENTIVES TO LIMIT  
HUMAN BIRTHS  
HELP  
TO ASSURE EACH FAMILY SUFFICIENT LAND FOR  
THEIR FOOD NEEDS  
IN THE SAHEL REGION OF AFRICA ?
```

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The input sequential text for the above output is:

```

/R1
DOES
/R2
HELP
/R3
IN THE SAHEL REGION OF AFRICA ?
/EL
DEVELOPING SOCIAL INCENTIVES TO LIMIT
HUMAN BIRTHS
/EL
TO ASSURE EACH FAMILY SUFFICIENT LAND FOR
THEIR FOOD NEEDS
// END

```

The slant (/) control keywords define the type of clause for the line(s)* immediately following them. These definitions are listed below.

```

/R1 = introductory clause
/R2 = relational clause
/R3 = qualifying clause
/EL = element text
// = physical end of sequential element text file

```

It is important for you to understand that MAKEIT orders the element texts sequentially, starting with one and incrementing by one. That is, element number ten in the list is defined by the text immediately following the tenth /EL card, and so on.

*The maximum number of lines is 10. The minimum number of lines is 1.

*The maximum number of characters on one line is 60. The minimum number of characters on one line is 1.

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FULLTEXT.QUERIES.CHECKLIST

To use full text queries, follow this checklist:

- 1) Using EDT, type in the element list to form a sequential element text file.
- 2) Run program MAKEIT with the sequential text file created in EDT (during step 1).
- 3) When typing the "DO" command executing ISMS-UD or CYCLE, be sure to type the name of the sequential text file as the second operand.
- 4) Answer "Y" to the "FULL TEST QUERIES DESIRED" question that is presented in ISMS-UD or CYCLE.

TEXT.DISPLAYING.MINIS

The following are some of the optional capabilities of the ISMS-UD full text query facility.

CLEARING THE SCREEN OF A CRT

If a CRT (cathode ray tube) terminal is being used for an ISM session, it is possible to have ISMS-UD clear the screen of the display before each query. To do this, all that is necessary is to make the screen clearing character the first character of the introductory clause. The clearing character is usually a CTRL L on most CRTs.

To clear the screen before each query, type a CTRL L as the first character of the introductory clause.

SPACING BETWEEN FRAMING CLAUSES

The term "framing clauses" denotes the combination of the introductory clause, relational clause, and qualifying clause. The output using full text queries is always in the form shown on page 10. It is

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sometimes desirable to
framing clauses, e.g.

the query with spacing between the

DOES

DEVELOPING SOCIAL INCENTIVES TO LIMIT
HUMAN BIRTHS

HELP

TO ASSURE EACH FAMILY SUFFICIENT LAND FOR
THEIR FOOD NEEDS

IN THE SAHEL REGION OF AFRICA ?

in order to increase readability. This can be achieved by inserting a
line with one blank character 1) after the introductory clause, 2)
before and after the relational clause, and 3) before the qualifying
clause.

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INTERPRETIVE STRUCTURAL MODELING SOFTWARE PACKAGE

The Interpretive Structural Modeling Software Package consists of three interactive FORTRAN IV programs which perform operations required to construct and modify Interpretive Structural Models. The following is a description of each program. References to John Warfield's book, Societal Systems: Planning, Policy, and Complexity, will be made to help the user locate the theory used.

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PROGRAM_MAKEIT

PURPOSE_OF_MAKEIT

Program MAKEIT accepts a sequential element text file, reconstructs it into a random (direct) access format for full text queries, and writes this information onto a query file. The program is conversational in nature to allow the user to display (show) on his or her terminal all elements, or specific pairs of elements, as they might appear during an ISM session.

If you do not know what a sequential element text file is, see "ISMS-UD FULL TEXT QUERY FACILITY".

HOW_TO_USE_MAKEIT

1. To invoke program MAKEIT, use the operating system command /DO MAKEIT. This command has one required operand - the name of the input sequential element text file (i.e., the name used in the EDT @WRITE command). The output query file is automatically named and initialized by prefixing the sequential element text file name with "RND.". The user does not have to worry about this file naming and creation.

Example.... /DO MAKEIT,(textfilename)

IF...an invalid file name* is typed, the operating system will produce these error messages and will not let you run MAKEIT:

```
% D531 INVALID FILENAME.  COMMAND TERMINATED.
% E015 ERROR IN PRECEEDING CMD - CMDS IGNORED TILL STEP OR LOGOFF.
% E804 ENDPROC RETURNED TO PRIMARY.
```

*A valid file name is 1 to 53 characters with no initial numerics and no embedded blanks.

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Retype the "/DO" COMMAND.

IF...the syntax of the /DO MAKEIT command is incorrect, the operating system will reject the command and type this error message:

% E140 OPERAND SYNTATICAL ERROR; REENTER THE DO COMMAND

IF...the /DO MAKEIT command is typed syntactically correct with a valid file name, the operating system will type these messages indicating that MAKEIT is being loaded for use:

% P500 LOADING VER# 001 OF ISMS.
FORTRAN IV PROGRAM MAKEIT STARTED --- MM/DD/YY

2. MAKEIT now begins reading the sequential element text file named in the /DO MAKEIT command.

IF...an unsuccessful read by MAKEIT occurs, the operating system will type this error message:

MAKEIT TERMINATED: TOO MUCH DATA REQUESTED FROM RECORD : P1-CTR
= NNNNNNNN.
SELECT DEBUG OUTPUT OPTIONS (D,S,A,B,C,HELP)

Type a "CTRL C" and the computer will return to the operating system mode enabling the user to correct the problem.

Reasons for an unsuccessful read operation are:

1. the @WRITE'textfilename':1-60: command in EDT was not used.
2. the file name specified in the /DO MAKEIT command was not a sequential element text file.

IF...MAKEIT encounters invalid sequential element text file syntax, these error messages will be printed and MAKEIT will terminate prematurely:

ERROR INVALID SLANT KEYWORD ENCOUNTERED
RE-EDIT TEXT FILE TO CORRECT

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MAKEIT TERMINATED

Re-edit the sequential element text file in EDT and try again.

IF...no problems are encountered, MAKEIT will print this message:

PERMFILE HAS BEEN CREATED FOR FULL TEXT QUERIES
SHOW(Y/N) ?

*

3. Answer "Y" or "N".

IF...the user types "N", MAKEIT will terminate and return control to the operating system.

IF...the user types "Y", this message will be typed:

SHOW ALL ELEMENTS ? (Y/N)

*

4. Answer "Y" or "N".

IF...the user types "Y", all of the elements are shown, i.e., printed at the terminal as they might appear during an ISM session. MAKEIT terminates and returns control to the operating system after showing all elements.

IF...the user types "N", this message will be typed:

SHOW WHICH ELEMENTS ?

*

5. Type the numbers of the two elements that you desire to see. MAKEIT will keep accepting and showing element pairs until zeros are typed. At that time, MAKEIT will terminate and return control to the operating system.

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PROGRAM_ISMS-UD

PURPOSE_OF_ISMS-UD

Program ISMS-UD is essentially the hub of the ISMS-UD Version 2.1 software package. ISMS-UD allows the user to embed the ISM utilizing the theory of transitive bordering on a reachability matrix* as well as modify the ISM. Elements that are adjacent on the digraph map, or elements that are on the same level or stage of the hierarchy and are not connected are able to be modified**. ISMS-UD maintains and operates on a reachability matrix which allows a digraph to be obtained after each operation on the matrix. ISMS-UD also allows the user to embed the ISM by using his own queries or the computer's queries.

HOW_TO_USE_ISMS-UD

1. To invoke ISMS-UD, use the operating system command /DO ISMS-UD. This command must include the following two operands: 1) the name of the permanent file which is to contain the ISM, and 2) the name of the sequential element text file for the full text queries option.

Example: /DO ISMS-UD,(modelfilename,textfilename)

IF...full text queries are not going to be used, type "JUNK" for the second operand.

IF...no model file exists, the operating system will create one and give it the name typed for operand one.

*John N. Warfield, *Societal Systems: Planning, Policy, and Complexity* (New York: Wiley-Interscience, 1976), p. 237.

**Warfield, p. 356.

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IF...a file name is not typed for both operands, the operating system will produce these error messages and will not let you run ISMS-UD:

% D531 INVALID FILENAME. COMMAND TERMINATED.
% E015 ERROR IN PRECEEDING CMD - CMDS IGNORED TILL STEP OR LOGOFF.
% E804 ENDPROC RETURNED TO PRIMARY.

IF...an invalid file name is entered, the operating system will respond with one of the following error messages:

% E146 ILLEGAL SYMBOLIC PARAMETERS DETECTED WHILE PROCESSING
PROCEDURE FILE; PROCEDURE FILE TERMINATED.

or

% E144 SYMBOLIC PARAMETER OPERAND ERROR, REENTER THE DO COMMAND.

IF...the syntax of the /DO ISMS-UD is incorrect, the operating system will reject the command and type this error message:

% E140 OPERAND SYNTATICAL ERROR; REENTER THE DO COMMAND.

After each of these error messages, retype the "/DO" COMMAND.

IF...the /DO ISMS-UD command is correctly typed, the operating system will type these messages indicating that ISMS-UD has been loaded for use:

% P500 VER# 2.0 OF ISMS-UD LOADED AT LOCATION 000000.

FORTRAN IV PROGRAM ISMSUD STARTED --- MM/DD/YY

2. ISMS-UD is now in control and it asks:

NEW SYSTEM ? (Y/N)

*

A "new system" is an ISM that has never been initialized. Answer "Y" or "N".

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IF..."Y" is typed, ISMS-UD will assume a new ISM is being generated and proceed to the next question.

IF..."N" is typed, ISMS-UD will try to read in the ISM from the model file named in operand one of the /DO ISMS-UD command. An unsuccessful attempt to read the ISM will result in one of the following error messages to be typed:

ISMSUD TERMINATED: READ OPERATION ON A NON-EXISTENT ISAM
FILE: P1-CTR = NNNNNNNN.
SELECT DEBUG OUTPUT OPTIONS (D,S,A,B,C,HELP)

BORD TERMINATED: ISAM SEQUENTIAL ERROR ODAE: P1- CTR =
NNNNNNNN.
SELECT DEBUG OUTPUT OPTIONS (D,S,A,B,C,HELP)

Type a "CTRL C" and the computer will return to the operating system enabling you to retype the /DO ISMS-UD command.

3. Next, ISMS-UD types:

TYPE ISMS-UD CMD ? (OR "HELP")

*

Type either a two letter ISMS-UD command described below or the keyword "HELP" in order to obtain a list of valid ISMS-UD commands.

ISMS-UD_COMMANDS

ISMS-UD recognizes the following two letter keywords as command input:

BO COMMAND

The BO command allows the user to invoke the transitive bordering algorithm to embed an element into the ISM. ISMS-UD will first ask:

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FULL TEXT QUERIES DESIRED ? (Y/N)

*

Answer "Y" or "N".

IF...the user types "Y", and the filename for operand two of the /DO ISMS-UD command is: 1) not the name of a sequential element text file, or 2) the name of a sequential element text file that has not undergone the program MAKEIT sequence, the operating system will type this error message:

ISMSUD TERMINATED: ISAM ERROR 0D9A: P1-CTR = NNNNNNNN.
SELECT DEBUG OUTPUT OPTIONS (D,S,A,B,C,HELP)

Type a "CTRL C" and the computer will return to the operating system mode which will enable you to retype the /DO ISMS-UD command.

Next, ISMS-UD asks:

SUBORDINATION RELATION ? (Y/N)

*

Answer "Y" or "N".

IF...you are in doubt, answer "N". A subordination relation is a transitive relation where if Element A is related to Element B, then Element B is not related to Element A. This is an optimization for the transitive bordering algorithm which eliminates the possibility of feedback (cycles) and reduces the number of questions required to embed an element for this bordering session only. Caution: be sure that you understand the meaning of a subordination relation. In most cases a subordination relation is not desired. When in doubt, answer "N".

Note: Once these questions have been answered in an ISM session, they cannot be changed until a TE command is received and a new session started.

The next question is:

TYPE NEXT ELEMENT NUMBER OR 0 FOR BREAK ?

*

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Type the number of the first element (if a new system) or the next element (if an old system) to be included in the ISM.

IF...a new system is being modeled, ISMS-UD will form a 1 X 1 reachability matrix and respond with the same question again to get a second element to begin the questioning process.

IF...a non-numeric entry is typed, ISMS-UD will respond with:

ERRORINPUT NOT NUMERIC--RETRY
*

IF...the number typed is already in the system, ISMS-UD will give this error message:

ERROR NNNN ALREADY IN SYSTEM INDEX

and ask for another element number.

IF...a 0 (zero) is typed, ISMS-UD will ask for the next command keyword.

IF...a valid element is typed, ISMS-UD will present the queries necessary to embed the element.

When queries are presented, answer with either "Y", "N", or "AB". ISMS-UD will retype the response it received and

IF...a "Y" or "N" is received, ISMS-UD will insert a one or zero, respectively, in the proper location in the reachability matrix and present the next query.

IF..."AB" is received, ISMS-UD will abort the bordering algorithm for the element being processed and ask for the next command keyword.

IF...an invalid response is received, ISMS-UD will issue an error message and present the query again.

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After the new element is properly embedded, ISMS-UD will ask for another element number to begin the bordering algorithm again.

BOQ COMMAND

The BOQ command allows the user to invoke the transitive bordering algorithm to embed an element into the ISM by using his or her own queries. (Note: If this is the first time the transitive bordering algorithm is being started, the "full text" and "subordination relation" questions will be asked as in the BO command; however, if they have already been asked, ISMS-UD will assume the answers to be the same as in the BO command.)

Next, ISMS-UD will ask for elements as it did in the BO command. When a valid new element is accepted, ISMS-UD will ask:

WHICH ELEMENTS TO BE COMPARED ? (TYPE 0 FOR AUTOMATIC)

*

Now type the elements to be used in the query in the order to be compared (i.e. if 2,3 is typed, the query will ask if 2 is related to 3). Remember...the new element must be one of the two elements to be compared.

IF...the relation between the elements typed has already been answered, ISMS-UD will respond with:

SORRY, THIS QUESTION HAS BEEN EITHER ASKED ALREADY OR
FILLED BY INFERENCE EL1 R EL2 = YES(NO)

IF...0 (zero) is typed, ISMS-UD will send a message acknowledging the 0 and will begin to issue queries automatically. The automatic queries option only lasts for the new element. When the next element is entered, ISMS-UD will again ask for elements to be compared.

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DI COMMAND

The DI command directs ISMS-UD to extract and print the digraph of the current reachability matrix ISM in a levels format. No other user interaction is required for the execution of this command.

DIS COMMAND

The DIS command directs ISMS-UD to extract and print the digraph of the current reachability matrix ISM in a stages format. No other user interaction is required for the execution of this command.

ADD COMMAND

The ADD command allows the user to add elements to the ISM. This command should be used in conjunction with the AE command to add elements and their relationships. ISMS-UD will type:

TYPE ELEMENTS TO BE ADDED ?
*

Type the number of the element to be added. ISMS-UD will keep accepting elements to be added until a zero is typed.

ELIM COMMAND

The ELIM command permits the elimination of elements from the ISM. ISMS-UD types this message:

TYPE ELEMENT NUMBERS TO BE ERASED ?
*

Type the number of the element you wish to eliminate from the ISM. ISMS-UD will keep accepting element numbers to be eliminated until a zero is typed.

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AE COMMAND

The AE command allows the user to add a relationship (edge) on the ISM. ISMS-UD will type:

TYPE <A REACHES TO B>

*

Type the two elements to be connected. The edges to be added are accepted until a zero is typed for one or both of the elements.

EE COMMAND

The EE command allows the user to eliminate a relationship (edge) on the ISM. ISMS-UD will type:

TYPE <A REACHES TO B> TO BE ERASED

*

Type the two elements to be disconnected. The edges to be disconnected are accepted until a zero is typed for one or both of the elements.

IF...either element is not on the minimum edge digraph (i.e. the elements are not connected in any way), ISMS-UD will issue an error message and ask for next command keyword.

IF...the edge entered does not exist (i.e. elements are connected but not directly), ISMS-UD will issue the following error message:

ERROR THE EDGE FROM NNN TO MMM DOES NOT

EXIST ON MINIMUM EDGE DIGRAPH

IF...the elements typed are in a cycle, ISMS-UD will type:

ERROR NNN AND MMM ARE IN A CYCLE

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"ELIM" NNN AND RE-ENTER USING THE "BO" COMMAND

Eliminate either one of the elements from the system and re-embed using either the BO, BOQ, or ADD and AE command.

PO COMMAND

The PO (Pooling) command permits the user to combine unconnected elements on the same level or stage of a hierarchy into one single element. ISMS-UD will type:

TYPE TWO ELEMENTS TO BE POOLED ?

*

Type the two elements to be pooled.

IF...the elements to be pooled are not on the same level or stage, ISMS-UD will issue the following error message and ask for two elements to be pooled.

ERROR NNN AND MMM ARE NOT ON THE SAME LEVEL

OR STAGE

IF...the two elements pass the error checks. ISMS-UD will type

NEW INDEX NAME ?

*

Type the number that the new element is to take on.

IF...the new index number typed is already in the system, ISMS-UD will inform the user of the error and ask for another new index number.

Elements to be pooled are accepted until two zeros are typed.

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EC COMMAND

The EC (Elementary Contraction) command allows the user to combine two adjacent elements on different levels or stages of a hierarchy into one single element. ISMS-UD will type:

TYPE TWO ELEMENTS TO BE CONTRACTED ?

*

Type the two elements to be contracted.

IF...the elements are not adjacent, ISMS-UD will issue the following error message and ask for two elements to be contracted.

ERROR NNN IS NOT ADJACENT TO MMM

IF...it is possible to contract the elements, ISMS-UD will type:

NEW INDEX NAME ?

*

Type the number that the new element is to take on.

IF...the new index number typed is already in the system, ISMS-UD will inform the user of the error and ask for another new index number.

Elements to be contracted are accepted until two zeros are typed.

TE COMMAND

The TE command directs ISMS-UD to terminate and return control to the operating system.

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ISMS-UD-RESIARIS

Occasionally, problems may arise with a computer system which will render it inoperative. These occurrences are usually infrequent and short. The UD computing facility maintains a 96% up time (i.e., the computer is operational 96% of the scheduled operation time). Even with this excellent record, the restart capabilities of ISMS-UD should be noted. The ISMS-UD software has been designed in such a way that the model permanent file is updated after an operation onto the ISM is completed. The updating operation occurs immediately before this message is typed:

TYPE ISMS-UD CMD ? OR ("HELP")

*

It can be seen that a small part of the ISM can be lost if the computer crashes (breaks down) while executing any of the embeddings or amending commands. Restarting is very easy however. See the procedure below.

How to restart a ISMS-UD session after a time-sharing system failure:

1. Type the /DO ISMS-UD command with the name of the model permanent file and the sequential element text file to be restarted.
2. Answer "N" to the "new system" question. ISMS-UD now has a copy of the ISM.
3. (OPTIONAL) Type "DI" to display the digraph of the restarted ISM.
4. Type the ISMS-UD command you were using when the computer crashed.
5. Continue normally.

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PROGRAM_CYCLE

PURPOSE_OF_CYCLE

Program CYCLE resolves cycles contained in an Interpretive Structural Model. The user creates and maintains a weighted adjacency matrix that denotes the strength of relationships between elements in the cycle. CYCLE allows the user to obtain listings of the resolution threshold, threshold matrix, and the geodetic paths*.

HOW_TO_USE_CYCLE

1. To invoke CYCLE, use the operating system command /DO CYCLE. This command must include the following two operands: 1) the name of the permanent file which is to contain the weighted matrix, and 2) the name of the sequential element text file for the full text queries option.

Example... /DO CYCLE,(weightedmatrixname,textfilename)

IF...full text queries are not going to be used, type "JUNK" for the second operand.

IF...no weighted matrix file exists, the operating system will create one and give it the name typed for operand 1.

IF...an invalid filename or a filename is not typed for both operands, the operating system will produce these error messages and will not let you run CYCLE:

% D531 INVALID FILENAME. COMMAND TERMINATED.
% E015 ERROR IN PRECEEDING CMD - CMDS IGNORED TILL STEP OR LOGOFF.
% E804 ENDPROC RETURNED TO PRIMARY.

*Warfield, pp. 328-43

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Retype the "/DO" COMMAND.

IF...the syntax of the /DO CYCLE command is incorrect, the operating system will reject the command and type this error message:

% E140 OPERAND SYNTATICAL ERROR; REENTER THE DO COMMAND.

IF...the /DO CYCLE command is correctly typed, the operating system will type these messages indicating that CYCLE is being loaded for use:

% P500 LOADING VER# 001 OF ISMS.
FORTRAN IV PROGRAM CYCLE STARTED --- MM/DD/YY

2. CYCLE is now in control and it asks:

FULL TEXT QUERIES DESIRED ? (Y/N)

*

Answer "Y" or "N".

IF...the user types "Y" and the file name for operand two of the /DO CYCLE command is: 1) not the name of a sequential element text file, or 2) the name of a sequential element text file that has not undergone the program MAKEIT sequence, the operating system will type the following error message after attempting to print the first full text query:

CYCLE TERMINATED: ISAM ERROR 0D9A: P1-CTR = NNNNNNNN.
SELECT DEBUG OUTPUT OPTIONS (D,S,A,B,C,HELP)

Type a "CTRL C" and the computer will return control to the operating system enabling you to retype the /DO CYCLE command.

3. The next question is:

NEW SYSTEM ? (Y/N)

*

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A "new system" is a weighted matrix that has never been initialized. Answer "Y" or "N".

IF..."N" is typed, CYCLE will try to read in the weighted matrix from the file named in operand one of the /DO CYCLE command. An unsuccessful attempt to read in the matrix will result in the following error message to be printed:

EXLST = OPENERR

CYCLE will terminate and return control to the operating system.

IF..."Y" is typed, CYCLE will then type:

NUMBER OF ELEMENTS (50 MAX.) ?

*

Type the number of elements that are contained in the argument cycle. CYCLE will then ask:

REGULAR INDEXING OF ELEMENTS DESIRED ? (Y/N)

*

Type "Y" or "N".

IF...the user types "Y", CYCLE assigns sequential numbers to the elements in the cycle. That is, the indexes for the elements are assigned as 1, 2, ... n (where: n is the number of elements).

IF...the user types "N", CYCLE then types:

ENTER INDEXES ONE AT A TIME

*

The user is now allowed to specify his own index numbers for the elements in the cycle. Type the integers that define the index set one per line.

The advantage of being able to specify index set numbers is apparent when utilizing the full text queries option. For example, suppose elements 3, 7, 59, and 62 of a particular ISM were in a cycle. If full text

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queries were desired, the user would have to specify his index set as 3, 7, 59, and 62 in order to be able to use the same text file as he did for the embedding session. If "regular indexing" were mistakenly used, the text for elements 1, 2, 3, and 4 would appear whenever full text queries were required and would lead to a great deal of confusion.

4. Next, CYCLE types:

TYPE CYCLE COMMAND (OR "HELP")

*

Type either a two letter CYCLE command or the keyword "HELP" in order to obtain a list of valid CYCLE commands.

CYCLE_COMMANDS

CYCLE recognizes the following two letter keywords as command input:

AL COMMAND

The AL command allows the user to add elements to the weighted matrix. CYCLE types this message:

ELEMENT NUMBER TO BE ADDED ?

*

Type the number of the element you wish to add. CYCLE will then ask:

BORDER ON THIS ELEMENT ? (Y/N)

*

Answer "Y" or "N".

IF...the user types "Y", the weights pertaining to this element can be conveniently assigned by "bordering" down the side and bottom of the weighted matrix. CYCLE will present queries to the user.

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IF...the user types "N", the weights are not able to be "bordered" in and must be entered via the CH command.

Added element numbers are accepted until a zero is typed.

CH COMMAND

The CH command allows the user to change any weight in the weighted matrix. CYCLE types this message:

ENTER A REACHES TO B, WEIGHT (3 NUMBERS)

*

Type the element indexes defining the relationship and the new weight. CYCLE will keep accepting changes until three zeros are typed.

The CH command can also be used to initially insert weights into the matrix.

DL COMMAND

The DL command allows the deletion of elements from the weighted matrix. CYCLE types this message:

ELEMENT NUMBER TO BE ERASED ?

*

Type the number of the element you desire to delete from the weighted matrix. CYCLE will keep accepting element numbers to be deleted until a zero is typed.

FI COMMAND

The FI command facilitates the loading of the weighted adjacency matrix by automatically presenting all the queries necessary to fill

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the matrix. In addition, if the full text queries option was selected, the FI command will use the text (rather than just the numbers) of the elements during the automatic process. The user should respond to each query by typing an integer (less than 10) which represents the weight of the relationship between the elements.

EI_COMMAND_RESTARTING. When using the FI command, the user must answer $(N^2)-N$ questions in order to completely fill the weighted matrix (where: "N" is the number of elements in the weighted matrix). A restart/termination feature is available since the number of questions could become quite large. To use the restart feature, simply type a weight of 10 for the element pair you wish to terminate with. CYCLE can now be terminated and resumed at a later time. If the FI command is typed, questioning will resume with the element pair for which the weight of 10 was specified. The restart feature may be used as many times as required.

PR COMMAND

The PR command prints the weighted matrix in its present state on the terminal. No other user interaction is required for the execution of this command.

RE COMMAND

The RE command is essentially the main command of CYCLE. This command calculates the resolution threshold of the current weighted matrix. The resolution threshold, threshold matrix, and geodetic paths are printed at the terminal. No other user interaction is required for the execution of this command.

TE COMMAND

The TE command directs CYCLE to terminate and return control to the operating system.

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APPENDIX_SAMPLE_RUNS

For each of the following sample program runs, the following text file is used for the full text queries option. The text file is presented as the user would type it in using EDT. Comments pertaining to each sample run will immediately follow the listing of the sample run.

EDI_SAMPLE_RUN

/EXEC EDT

% P500 VER# 10B OF EDT LOADED AT LOCATION 000000.

*** EDITOR LOADED (VER 10B), BEGIN TYPE IN:

```
1.0000 /R1
2.0000 DOES
3.0000 /R2
4.0000 WEIGH MORE THAN
5.0000 /R3
6.0000 IN MOST CASES?
7.0000 /EL
8.0000 A BREATH OF AIR (1)
9.0000 /EL
10.0000 AN OLD TENNIS SHOE (2)
11.0000 /EL
12.0000 A FIRE TRUCK (3)
13.0000 /EL
14.0000 A FEATHER (4)
15.0000 /EL
16.0000 AN ELEPHANT (5)
17.0000 // END
18.0000 @WRITE TEXT.SAMPLE:1-60:
TEXT.SAMPLE IS IN THE CATALOG
OVERWRITE ? (Y,N) Y
```

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18.0000 @HALT

NOTE: If the text file has never been written to the disc, the overwrite message following the "@WRITE" command will not be presented. When the computer returns a slash ("/") after the "@HALT" command, the ISMS-UD programs can be executed.

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MAKEIT_SAMPLE_RUN

/DO MAKEIT.(TEXT.SAMPLE)
% P500 VER# 001 OF ISMS LOADED AT LOCATION 000000.
FORTRAN IV PROGRAM MAKEIT STARTED --- 04/10/79

PERMFILE HAS BEEN CREATED FOR FULL TEXT QUERIES
SHOW(Y/N) ?
*Y

SHOW ALL ELEMENTS(Y/N) ?
*Y
DOES
A BREATH OF AIR (1)
WEIGH MORE THAN
AN OLD TENNIS SHOE (2)
IN MOST CASES ?
DOES
A FIRE TRUCK (3)
WEIGH MORE THAN
A FEATHER (4)
IN MOST CASES ?
DOES
AN ELEPHANT (5)
WEIGH MORE THAN
A BREATH OF AIR (1)
IN MOST CASES ?
DOES
AN OLD TENNIS SHOE (2)
WEIGH MORE THAN
A FIRE TRUCK (3)
IN MOST CASES ?
DOES
A FEATHER (4)
WEIGH MORE THAN
AN ELEPHANT (5)
IN MOST CASES ?
** FORTRAN ** CALL EXIT

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NOTE. If an odd number of elements exists in the text file, as in the preceding example, MAKEIT will show each element twice so as to display queries until the end of file occurs after a complete query is presented.

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ISMS=UD_SAMPLE_BUN

/DO ISMS-UD.(MODEL.SAMPLE,TEXT.SAMPLE)
% P500 VER# 002 OF ISMS-UD LOADED AT LOCATION 000000.
FORTRAN IV PROGRAM ISMSUD STARTED --- 04/10/79

NEW SYSTEM ? (Y/N)
*Y

TYPE ISMS-UD CMD ? (OR "HELP")
*HELP

ISMS - U D C O M M A N D S

EMBEDDING

BO - TRANSITIVE BORDERING METHOD
BOQ - TRANSITIVE BORDERING WITH SELECTABLE QUERIES

DISPLAYING

DI - DISPLAY MINIMUM EDGE DIGRAPH
IN A LEVELS FORMAT

DIS - DISPLAY MINIMUM EDGE DIGRAPH
IN A STAGES FORMAT

SUBSTANTIVE AMENDING

ADD - ADD ELEMENTS

ELIM - ELIMINATE ELEMENTS

AE - ADD EDGES (RELATIONSHIPS)
ON THE MINIMUM EDGE DIGRAPH .

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EE - ERASE EDGES (RELATIONSHIPS)
ON THE MINIMUM EDGE DIGRAPH

FORMAT AMENDING

PO - POOL ELEMENTS

EC - ELEMENTARY CONTRACTION

TERMINATION

TE - TERMINATE ISMS-UD PROGRAM

NOTE

HELP - REPRINTS ABOVE LIST

TYPE ISMS-UD CMD ? (OR "HELP")
*BO

FULL TEXT QUERIES DESIRED ? (Y/N)
*Y
SUBORDINATION RELATION ? (Y/N)
*N

TYPE NEXT ELEMENT NUMBER OR 0 FOR BREAK ?
*1

TYPE NEXT ELEMENT NUMBER OR 0 FOR BREAK ?
*2

DOES
A BREATH OF AIR (1)
WEIGH MORE THAN
AN OLD TENNIS SHOE (2)
IN MOST CASES ?
*N N

DOES
AN OLD TENNIS SHOE (2)
WEIGH MORE THAN

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A BREATH OF AIR (1)
IN MOST CASES ?
*Y Y

TYPE NEXT ELEMENT NUMBER OR 0 FOR BREAK ?
*0

TYPE ISMS-UD CMD ? (OR "HELP")
*DI

LEVEL NO. 1

1

LEVEL NO. 2

2 => 1,

TYPE ISMS-UD CMD ? (OR "HELP")
*BOQ

TYPE NEXT ELEMENT NUMBER OR 0 FOR BREAK ?
*3
WHICH ELEMENTS TO BE COMPARED ? (TYPE 0 FOR AUTOMATIC)
*3,2

DOES

A FIRE TRUCK (3)
WEIGH MORE THAN
AN OLD TENNIS SHOE (2)
IN MOST CASES ?
*Y Y

WHICH ELEMENTS TO BE COMPARED ? (TYPE 0 FOR AUTOMATIC)
*3,1

SORRY, THIS QUESTION HAS BEEN EITHER ASKED ALREADY OR FILLED
BY INFERENCE 3 R 1 = YES
WHICH ELEMENTS TO BE COMPARED ? (TYPE 0 FOR AUTOMATIC)

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*0

0 ACKNOWLEDGED, BEGINNING AUTOMATIC QUERIES

DOES

AN OLD TENNIS SHOE (2)

WEIGH MORE THAN

A FIRE TRUCK (3)

IN MOST CASES ?

*N N

TYPE NEXT ELEMENT NUMBER OR 0 FOR BREAK ?

*0

TYPE ISMS-UD CMD ? (OR "HELP")

*DI

LEVEL NO. 1

1

LEVEL NO. 2

2 => 1,

LEVEL NO. 3

3 => 2,

TYPE ISMS-UD CMD ? (OR "HELP")

*ADD

TYPE ELEMENTS TO BE ADDED ?

*4

*5

*0

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TYPE ISMS-UD CMD ? (OR "HELP")
 *DI

LEVEL NO. 1

1
4
5

LEVEL NO. 2

2 => 1,

LEVEL NO. 3

3 => 2,

TYPE ISMS-UD CMD ? (OR "HELP")
 *AE

TYPE <A REACHES TO B>
 *4,1
 *5,3
 *3,5
 *0,0

TYPE ISMS-UD CMD ? (OR "HELP")
 *DI

CYCLE ON 3, 5,

LEVEL NO. 1

1

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LEVEL NO. 2

2 => 1,
4 => 1,

LEVEL NO. 3

3 => 2,

TYPE ISMS-UD CMD ? (OR "HELP")
*EETYPE <A REACHES TO B> TO BE ERASED
*4,1
*3,5***ERROR*** 3 AND 5 ARE IN A CYCLE.
"ELIM" 3 AND RE-ENTER USING THE "BO" COMMANDTYPE ISMS-UD CMD ? (OR "HELP")
*DI

CYCLE ON 3, 5,

LEVEL NO. 1

1
4

LEVEL NO. 2

2 => 1,

LEVEL NO. 3

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3 => 2,

TYPE ISMS-UD CMD ? (OR "HELP")
*PO

TYPE TWO ELEMENTS TO BE POOLED ?
*3,5
NEW INDEX NUMBER ?
*6

TYPE TWO ELEMENTS TO BE POOLED ?
*0,0

TYPE ISMS-UD CMD ? (OR "HELP")
*DI

LEVEL NO. 1

1
4

LEVEL NO. 2

2 => 1,

LEVEL NO. 3

6 => 2,

TYPE ISMS-UD CMD ? (OR "HELP")
*EC

TYPE TWO ELEMENTS TO BE CONTRACTED ?
*2,1
NEW INDEX NUMBER ?
*7

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TYPE TWO ELEMENTS TO BE CONTRACTED ?

*0,0

TYPE ISMS-UD CMD ? (OR "HELP")

*DI

LEVEL NO. 1

7

4

LEVEL NO. 2

6 => 7,

TYPE ISMS-UD CMD ? (OR "HELP")

*ELIM

TYPE ELEMENT NUMBERS TO BE ERASED ?

*4

*6

*0

TYPE ISMS-UD CMD ? (OR "HELP")

*DI

LEVEL NO. 1

7

TYPE ISMS-UD CMD ? (OR "HELP")

*TE

**FORTRAN ** STOP

NOTE: Whenever a 0 (zero) is to be entered to discontinue an ISMS-UD command, a "CTRL C" can be typed instead of the zero to save time.

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the computer recognizes this input as a zero. This also holds true for the MAKEIT (when showing which elements) and CYCLE programs.

The above sample run is a very simplified-example of how each ISMS-UD command affects the ISM. During an actual ISM session, a digraph would not usually be printed until the session is completed. The commands can be used in any logical order.

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CYCLE_SAMPLE_RUN

/DO CYCLE,(WEIGHT.SAMPLE,TEXT.SAMPLE)
% P500 VER# 001 OF ISMS LOADED AT LOCATION 000000.
FORTRAN IV PROGRAM CYCLE STARTED --- 04/17/79

FULL TEXT QUERIES DESIRED ? (Y/N)

*Y

NEW SYSTEM ? (Y/N)

*Y

NUMBER OF ELEMENTS (50 MAX.) ?

*4

REGULAR INDEXING OF ELEMENTS DESIRED ? (Y/N)

*Y

TYPE CYCLE COMMAND (OR "HELP")

*HELP

HELP MESSAGE

AL - ADD AN ELEMENT TO SYSTEM
DL - DELETE AN ELEMENT FROM SYSTEM
FI - FILL SYSTEM (ASSIGN WEIGHTS)
CH - CHANGE WEIGHT OF A RELATIONSHIP
RE - RESOLVE SYSTEM
PR - PRINT SYSTEM OUT
TE - TERMINATE SESSION
HELP - REPRINTS ABOVE LIST

TYPE CYCLE COMMAND (OR "HELP")

*AL

ELEMENT NUMBER TO BE ADDED ?

*5

BORDER ON THIS ELEMENT ? (Y/N)

*Y

A BREATH OF AIR (1)
HAS BEEN RELATED TO
AN ELEPHANT (5)
WEIGHT ?

*2

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AN OLD TENNIS SHOE (2)
HAS BEEN RELATED TO
AN ELEPHANT (5)
WEIGHT ?

*8

A FIRE TRUCK (3)
HAS BEEN RELATED TO
AN ELEPHANT (5)
WEIGHT ?

*7

A FEATHER (4)
HAS BEEN RELATED TO
AN ELEPHANT (5)
WEIGHT ?

*8

AN ELEPHANT (5)
HAS BEEN RELATED TO
A BREATH OF AIR (1)
WEIGHT ?

*4

AN ELEPHANT (5)
HAS BEEN RELATED TO
AN OLD TENNIS SHOE (2)
WEIGHT ?

*5

AN ELEPHANT (5)
HAS BEEN RELATED TO
A FIRE TRUCK (3)
WEIGHT ?

*6

AN ELEPHANT (5)
HAS BEEN RELATED TO
A FEATHER (4)
WEIGHT ?

*2

ELEMENT NUMBER TO BE ADDED ?

*0

TYPE CYCLE COMMAND (OR "HELP")

*PR

| | | | | |
|-----|-------|-------|-------|-------|
| 1=> | 2(0), | 3(0), | 4(0), | 5(9), |
| 2=> | 1(0), | 3(0), | 4(0), | 5(8), |
| 3=> | 1(0), | 2(0), | 4(0), | 5(7), |
| 4=> | 1(0), | 2(0), | 3(0), | 5(8), |
| 5=> | 1(4), | 2(5), | 3(6), | 4(2), |

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TYPE CYCLE COMMAND (OR "HELP")

*DL

ELEMENT NUMBER TO BE ERASED ?

*5

*0

TYPE CYCLE COMMAND (OR "HELP")

*PR

| | | | |
|-----|-------|-------|-------|
| 1=> | 2(0), | 3(0), | 4(0), |
| 2=> | 1(0), | 3(0), | 4(0), |
| 3=> | 1(0), | 2(0), | 4(0), |
| 4=> | 1(0), | 2(0), | 3(0), |

TYPE CYCLE COMMAND (OR "HELP")

*FI

A BREATH OF AIR (1)
HAS BEEN RELATED TO
AN OLD TENNIS SHOE (2)
WEIGHT ?

*7

A BREATH OF AIR (1)
HAS BEEN RELATED TO
A FIRE TRUCK (3)
WEIGHT ?

*8

A BREATH OF AIR (1)
HAS BEEN RELATED TO
A FEATHER (4)
WEIGHT ?

*9

AN OLD TENNIS SHOE (2)
HAS BEEN RELATED TO
A BREATH OF AIR (1)
WEIGHT ?

*7

AN OLD TENNIS SHOE (2)
HAS BEEN RELATED TO
A FIRE TRUCK (3)
WEIGHT ?

*6

AN OLD TENNIS SHOE (2)
HAS BEEN RELATED TO
A FEATHER (4)
WEIGHT ?

*7

A FIRE TRUCK (3)

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HAS BEEN RELATED TO
A BREATH OF AIR (1)
WEIGHT ?

*0

A FIRE TRUCK (3)
HAS BEEN RELATED TO
AN OLD TENNIS SHOE (2)
WEIGHT ?

*9

A FIRE TRUCK (3)
HAS BEEN RELATED TO
A FEATHER (4)
WEIGHT ?

*7

A FEATHER (4)
HAS BEEN RELATED TO
A BREATH OF AIR (1)
WEIGHT ?

*9

A FEATHER (4)
HAS BEEN RELATED TO
AN OLD TENNIS SHOE (2)
WEIGHT ?

*6

A FEATHER (4)
HAS BEEN RELATED TO
A FIRE TRUCK (3)
WEIGHT ?

*8

TYPE CYCLE COMMAND (OR "HELP")

*PR

| | | | |
|-----|-------|-------|-------|
| 1=> | 2(7), | 3(8), | 4(9), |
| 2=> | 1(7), | 3(6), | 4(7), |
| 3=> | 1(0), | 2(9), | 4(7), |
| 4=> | 1(9), | 2(6), | 3(8), |

TYPE CYCLE COMMAND (OR "HELP")

*CH

ENTER A REACHES TO B, WEIGHT (3 NUMBERS)

*3,1,8

*0,0,0

TYPE CYCLE COMMAND (OR "HELP")

*PR

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```

1=> 2(7), 3(8), 4(9),
2=> 1(7), 3(6), 4(7),
3=> 1(8), 2(9), 4(7),
4=> 1(9), 2(6), 3(8),

```

```

TYPE CYCLE COMMAND (OR "HELP")
*RE

```

```

***CYCLE RESOLVED***
THRESHOLD==> 7

```

```

                THRESHOLD MATRIX
1=> 2(7), 3(8), 4(9),
2=> 1(7), 4(7),
3=> 1(8), 2(9), 4(7),
4=> 1(9), 3(8),

```

| NUMBER | #LINKS | ELEMENTS | PATH |
|--------|--------|----------|---------|
| 1 | 2 | 2, 1 | 2 1 2 |
| 2 | 2 | 3, 1 | 3 1 3 |
| 3 | 3 | 3, 2 | 3 2 1 3 |
| 4 | 2 | 4, 1 | 4 1 4 |
| 5 | 3 | 4, 2 | 4 1 2 4 |
| 6 | 2 | 4, 3 | 4 3 4 |

```

TYPE CYCLE COMMAND (OR "HELP")
*TE
** FORTRAN ** CALL EXIT

```

NOTE: As in the ISMS-UD program, a "CTRL C" can be used instead of zeros when terminating a CYCLE command.

Also as with ISMS-UD, the above example is a very simplified run of the CYCLE program. The user may wish to resolve the matrix several times during the session to obtain the threshold of resolution desired. The commands may also be executed in any logical order as with the ISMS-UD commands.

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Attachment No. 3

DELTA CHARTS FOR ISMS

Prepared by:

David R. Yingling, Jr.

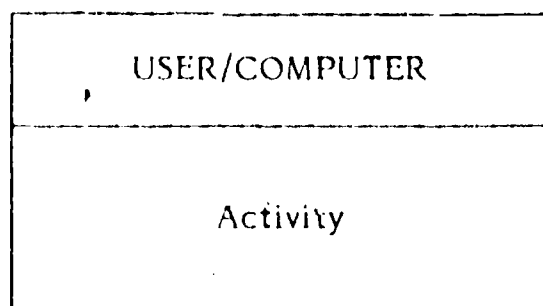
User DELTA Charts

This attachment consists of DELTA charts that describe the user and machine decision interactions that occur while operating ISMS. These DELTA charts are provided for both the programmer and the user in order to facilitate the installation and the use of ISMS. At first glance, you might notice that the DELTA charts look like programming flow charts. However, these charts do not document program flow. Instead, they document the actions of both the user and computer at a typical ISM computer session, therefore, the DELTA charts should be used only for instructional and system overview purposes.

How to Read the DELTA Charts

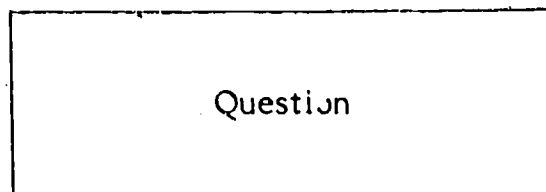
The charts depict the flow of activities at a typical ISM computer session. They convey a great deal of information in a highly structured format and a relatively small amount of space. This information includes activities, decisions, time flow, logic connections, and who is responsible for each activity or decision.

Several symbols appear on the charts; these are explained and illustrated in the next few paragraphs.* Activity boxes are the most common symbol used on the charts. An activity box is divided into two parts; the lower part shows an activity and the top part shows who is responsible for carrying out the activity. In ISM, either the user or the computer is responsible for each activity.

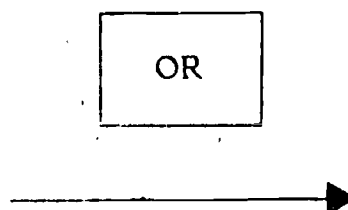


*This material is adapted from Warfield (1976, pp. 421-425).

A decision box is similar to an activity box, but it only has one part. By answering the question shown on the decision box, the user chooses among several alternatives paths leading from the decision box to subsequent boxes.

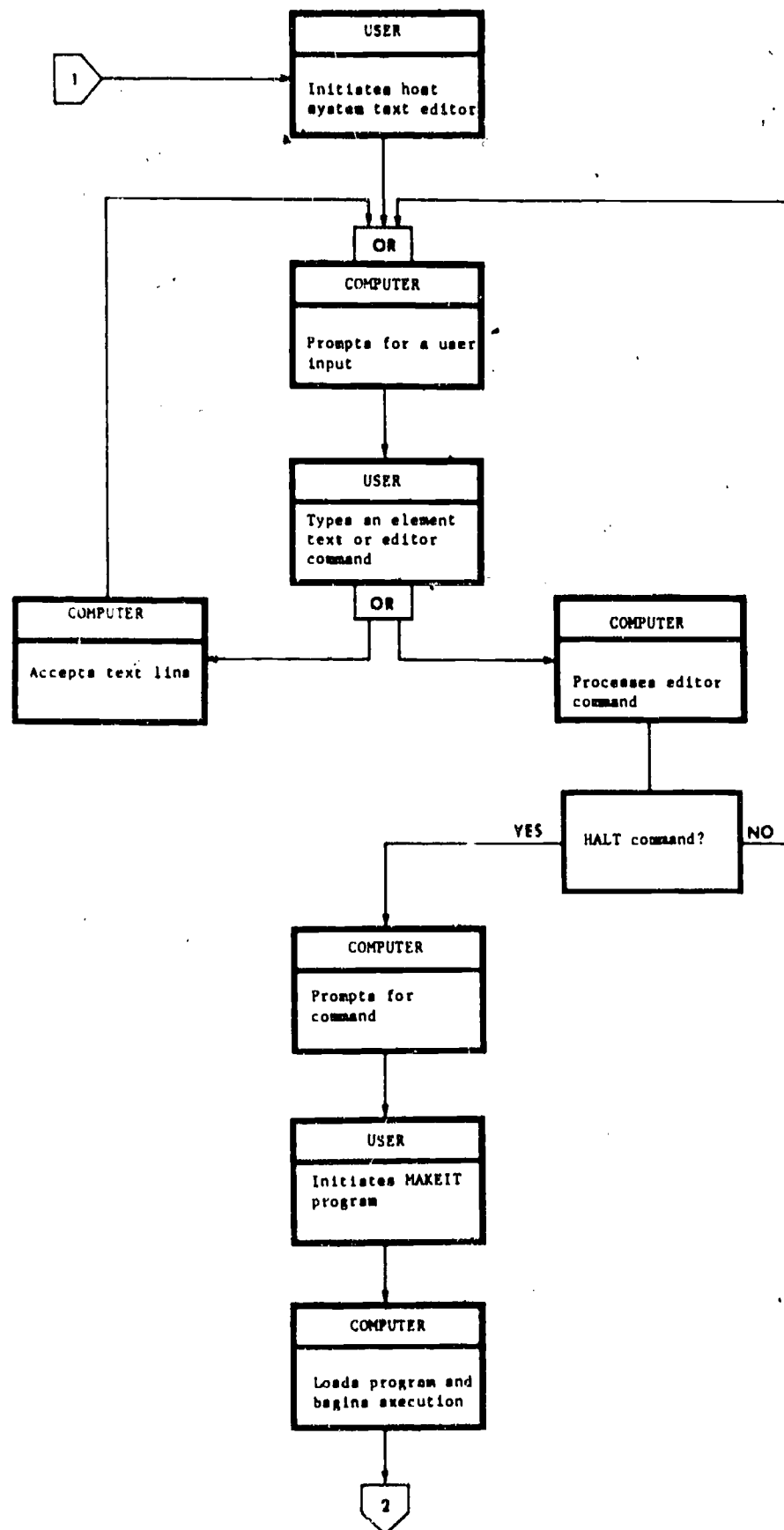


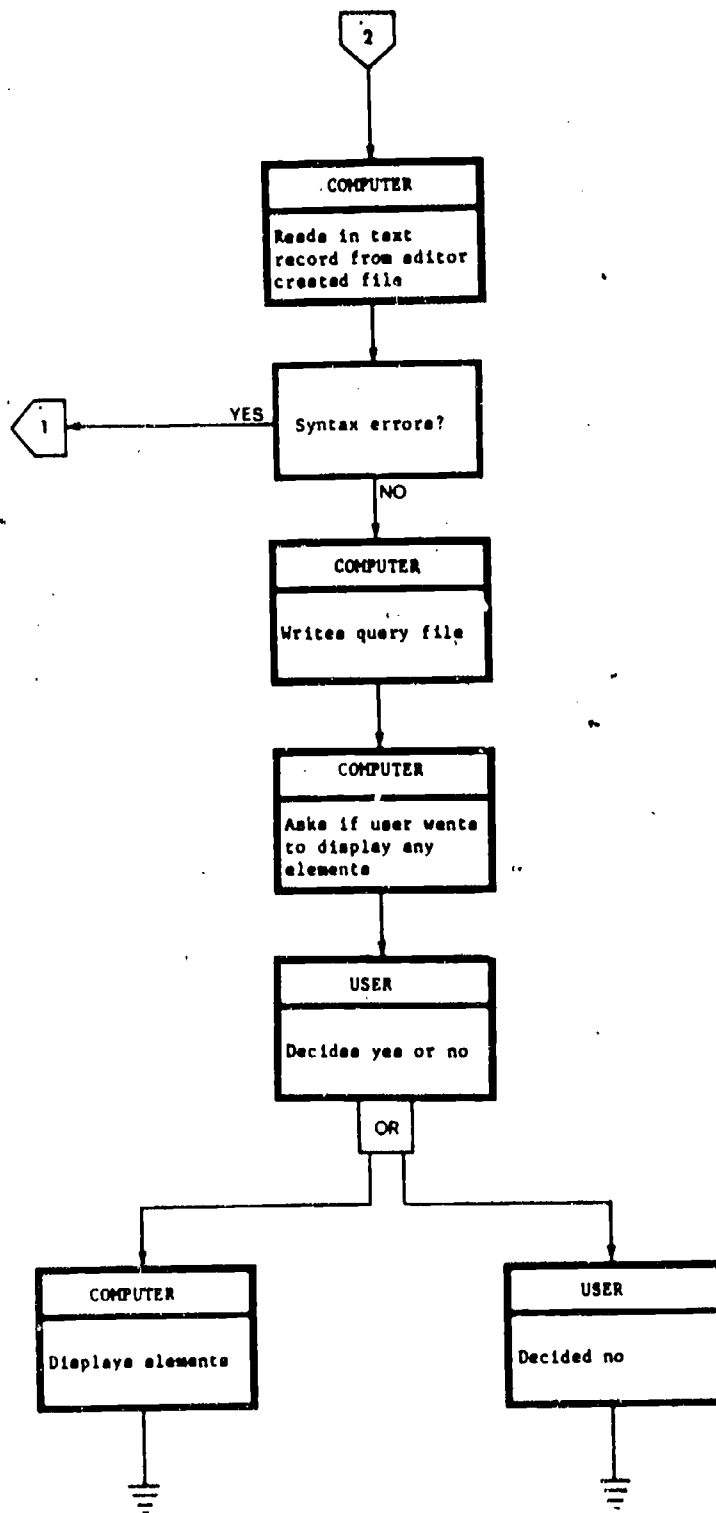
The OR box is interpreted as an "exclusive OR." One and only one of the preceding activities or decisions can occur at a given time. The lines that join the various boxes represent only the flow of time, except at the output of a decision box where lines also represent the various decisions that could be made. In that case, the lines are labeled; usually with either YES or NO.



2.00

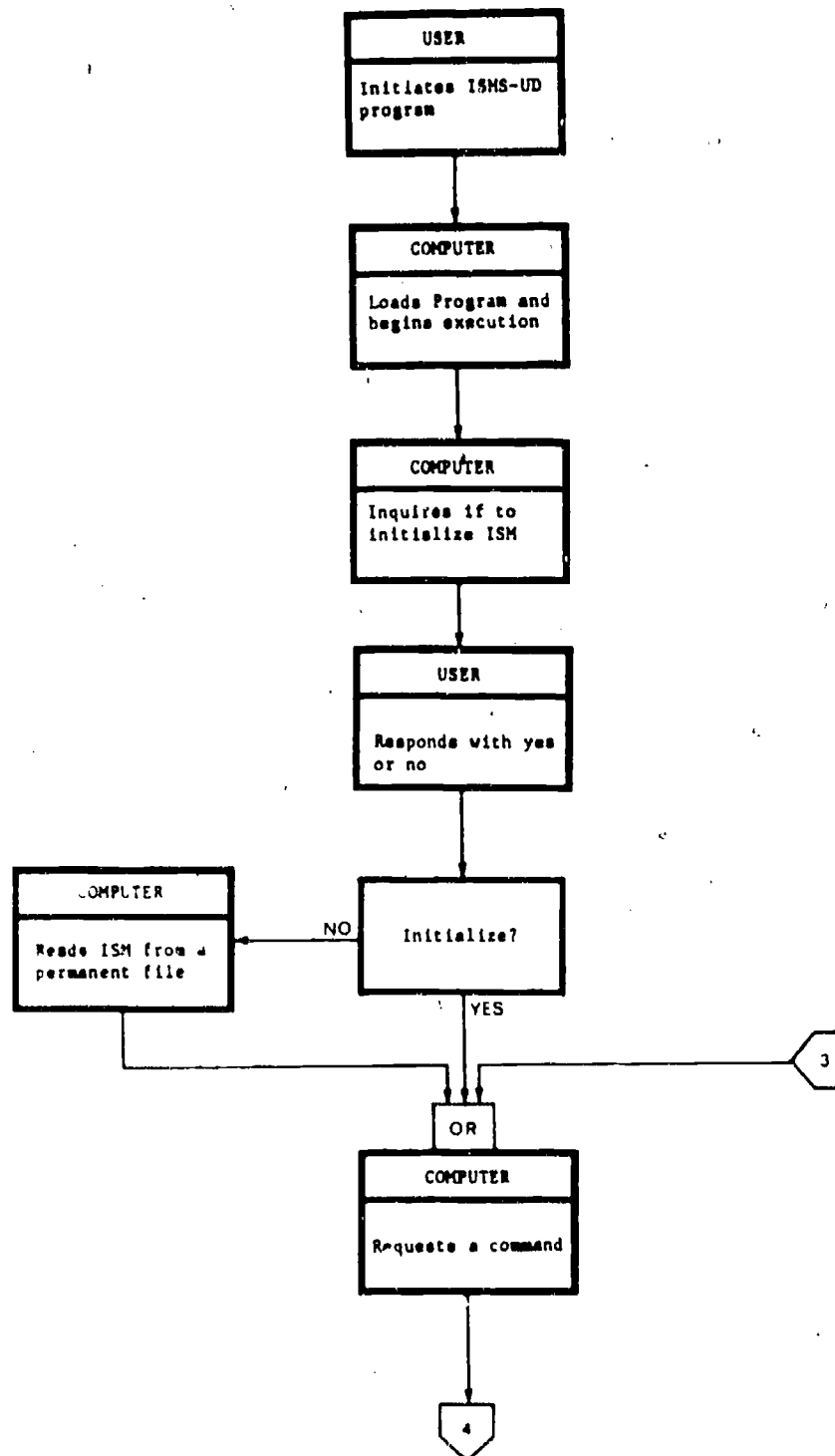
ISMS-US FULL TEXT QUERY FACILITY USE

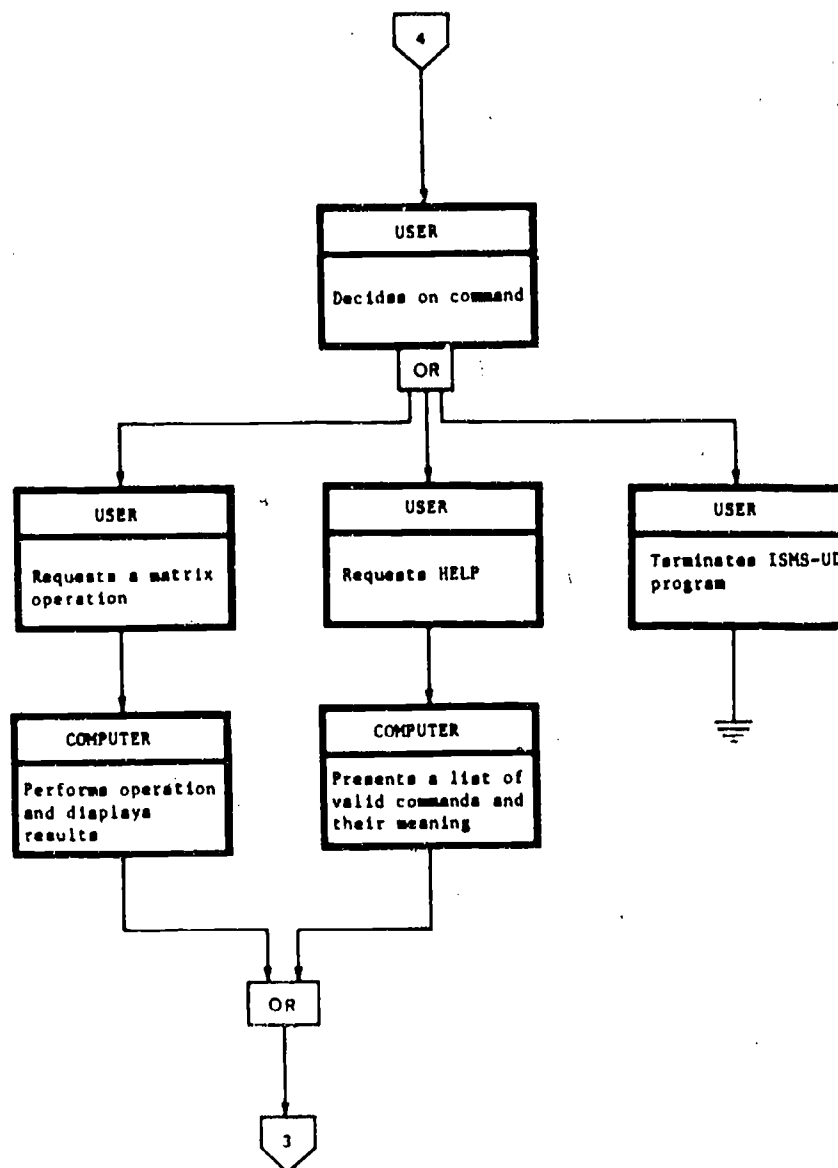




252

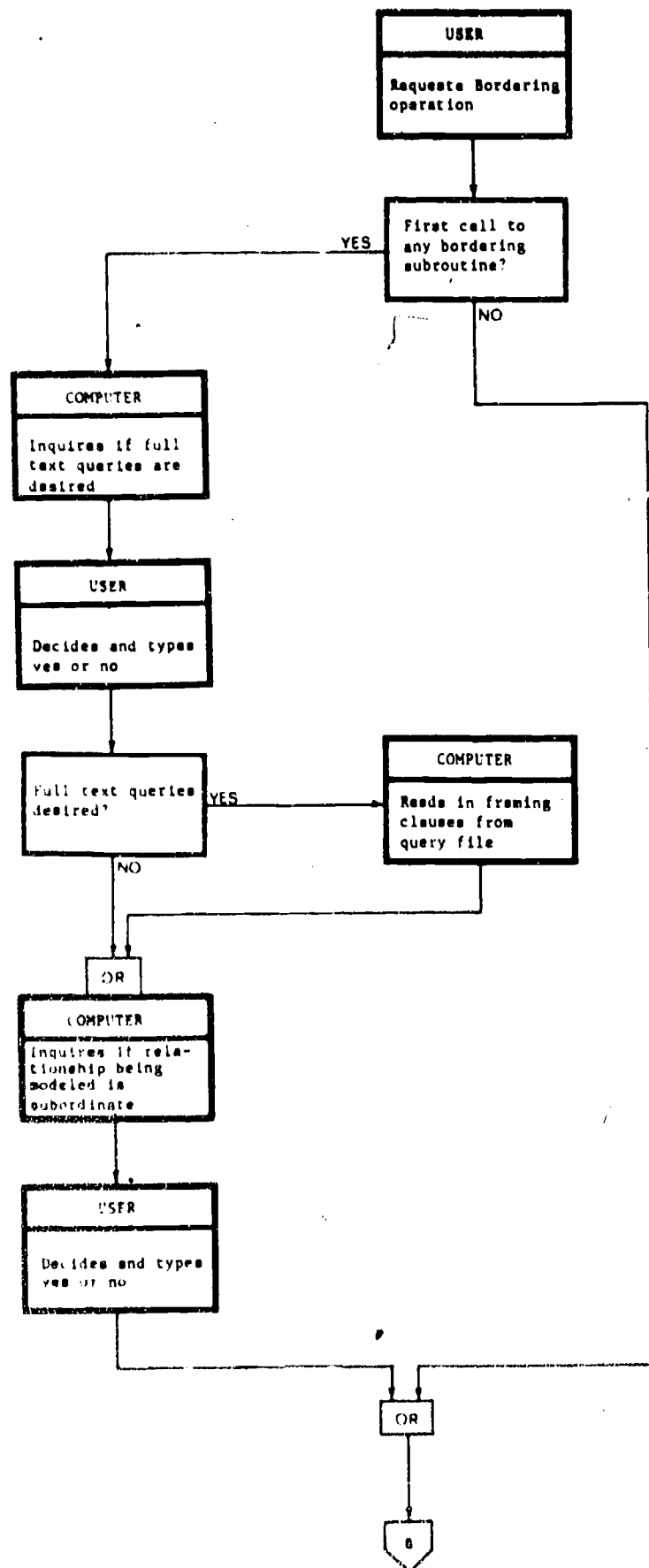
USING PROGRAM ISMS-UD

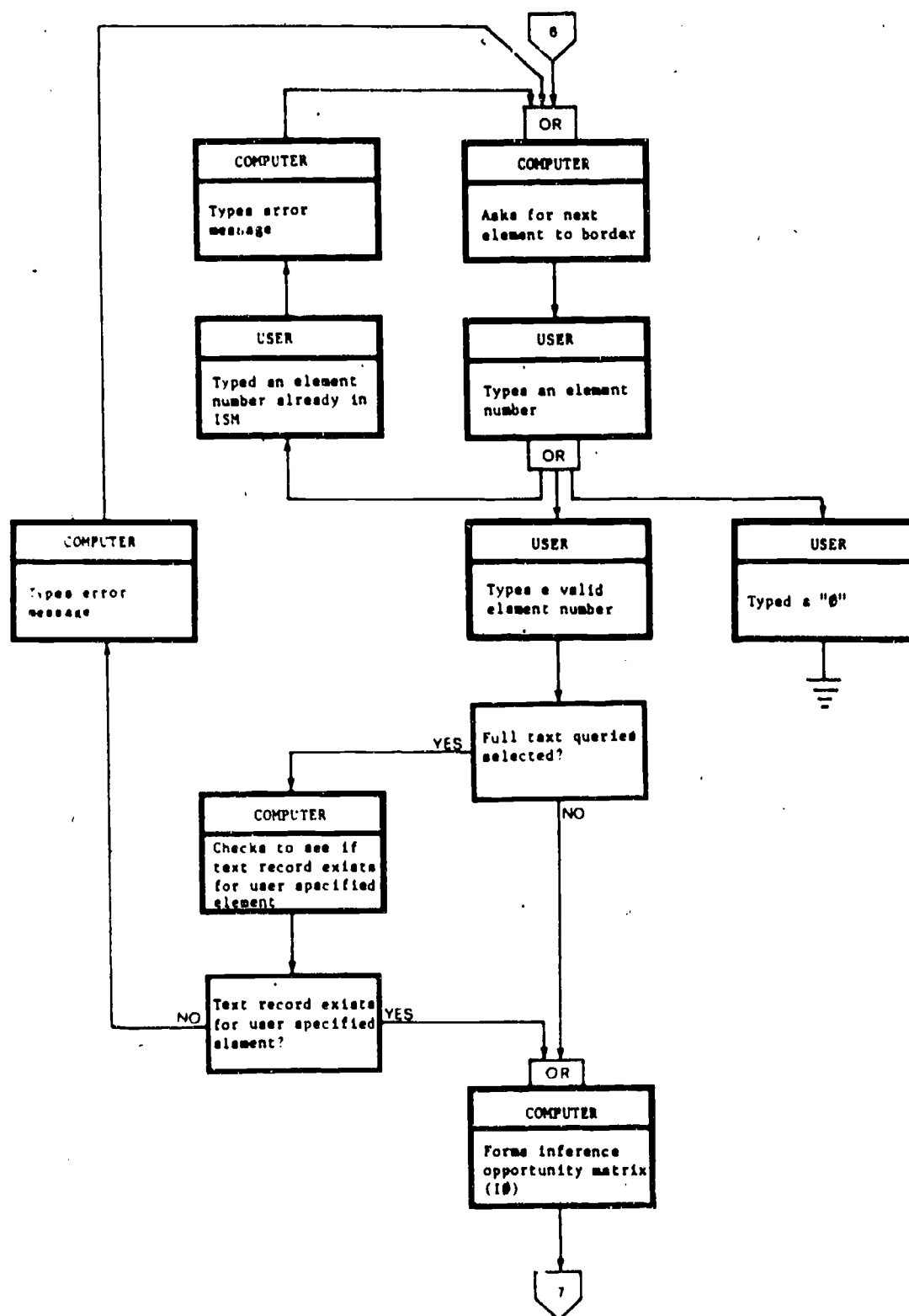




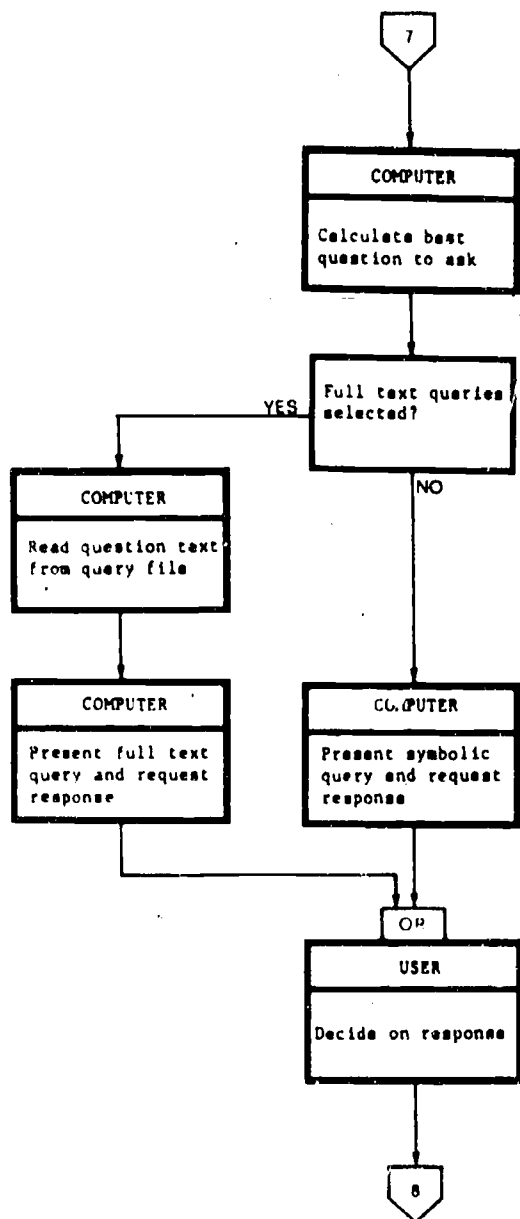
254

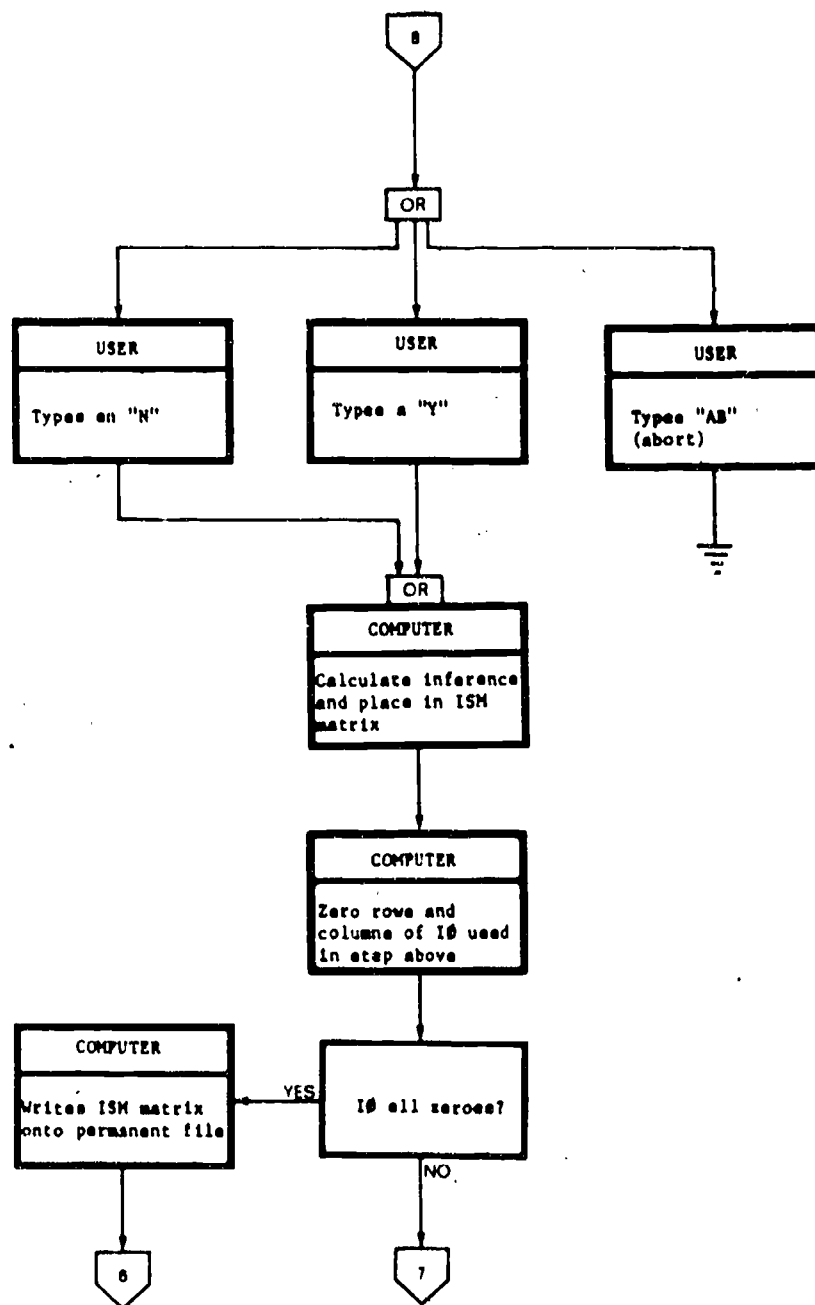
BO COMMAND





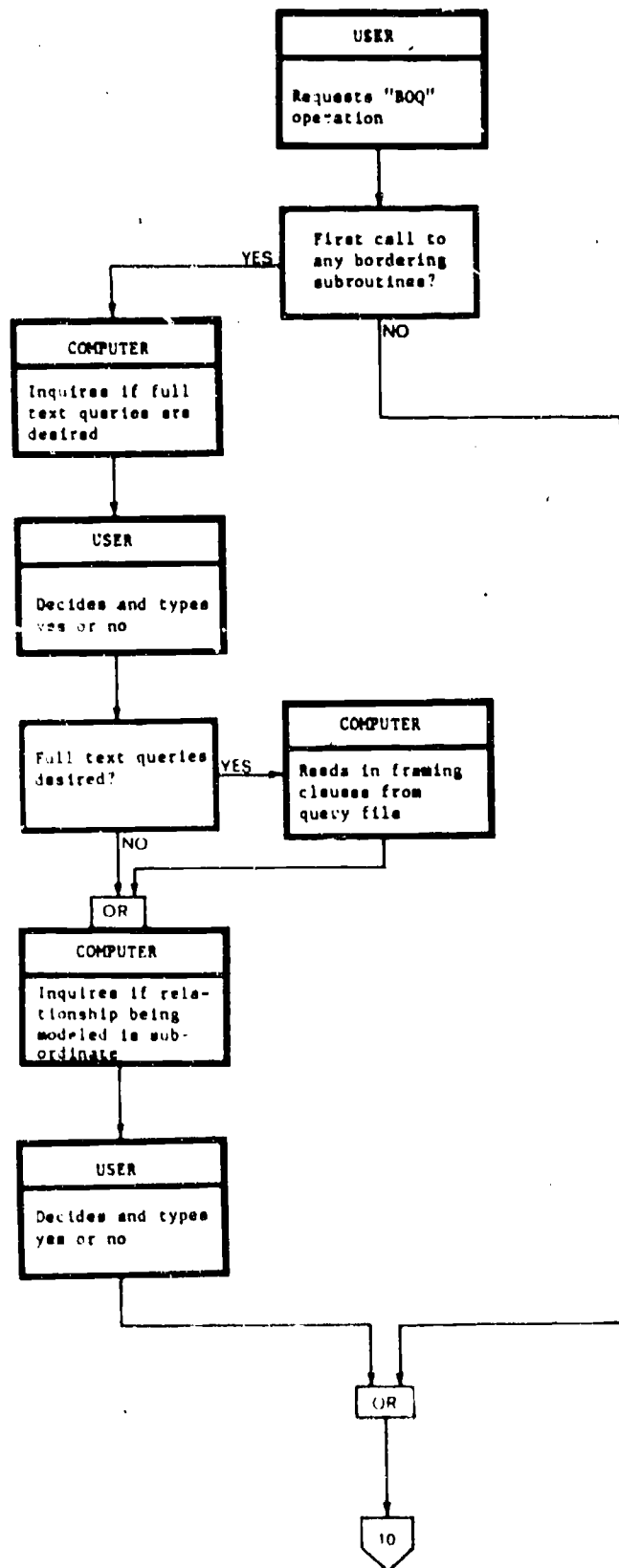
250

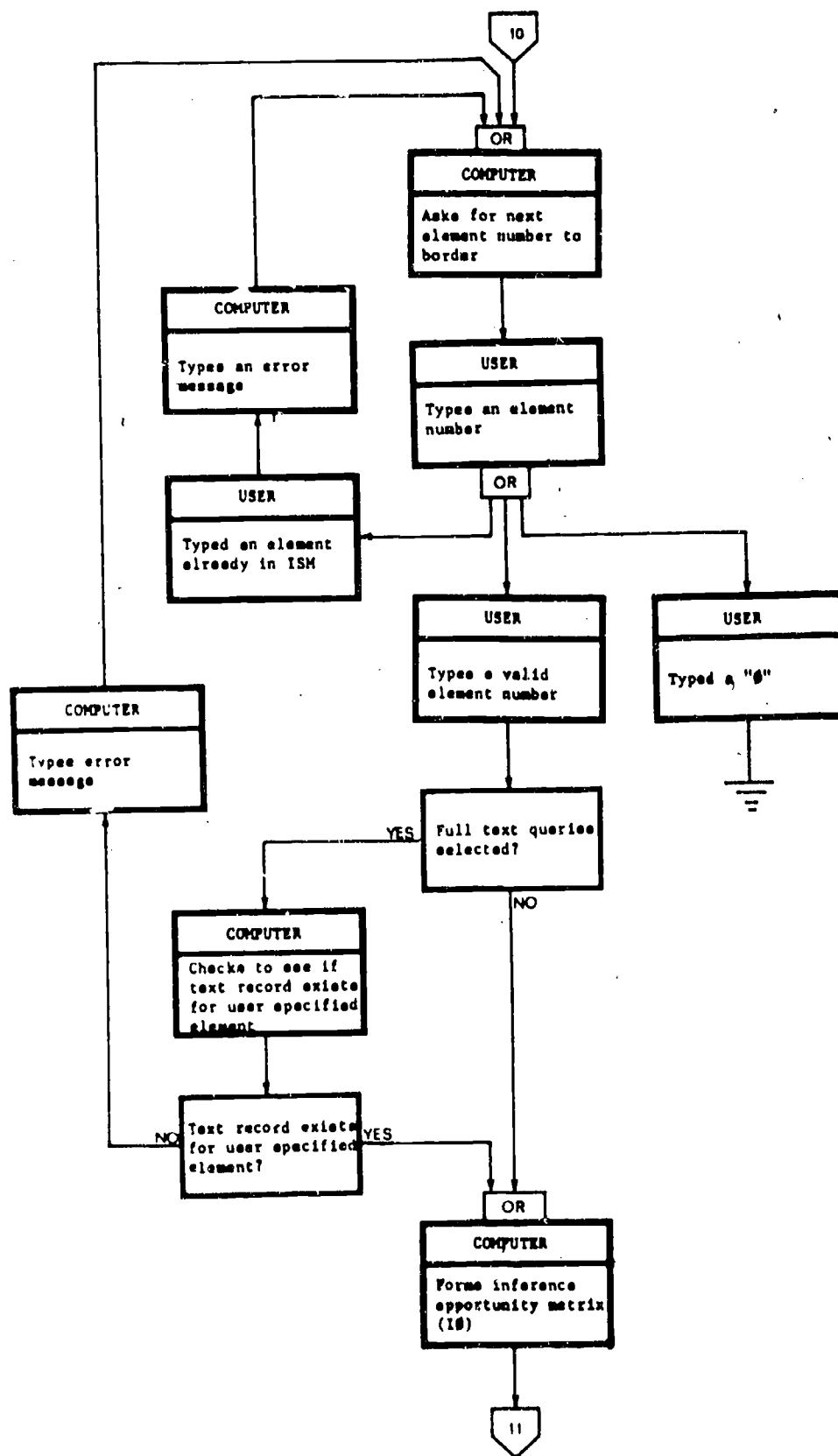




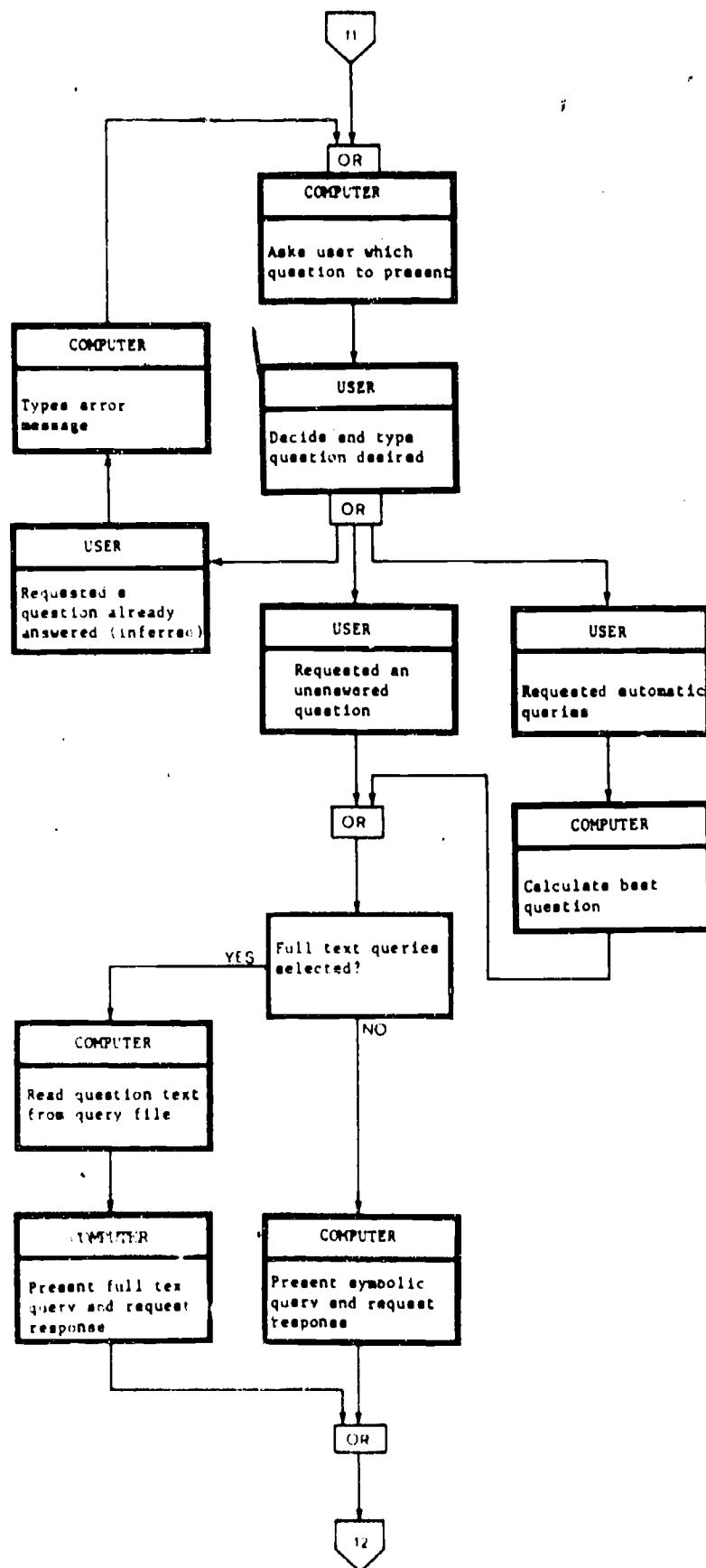
258

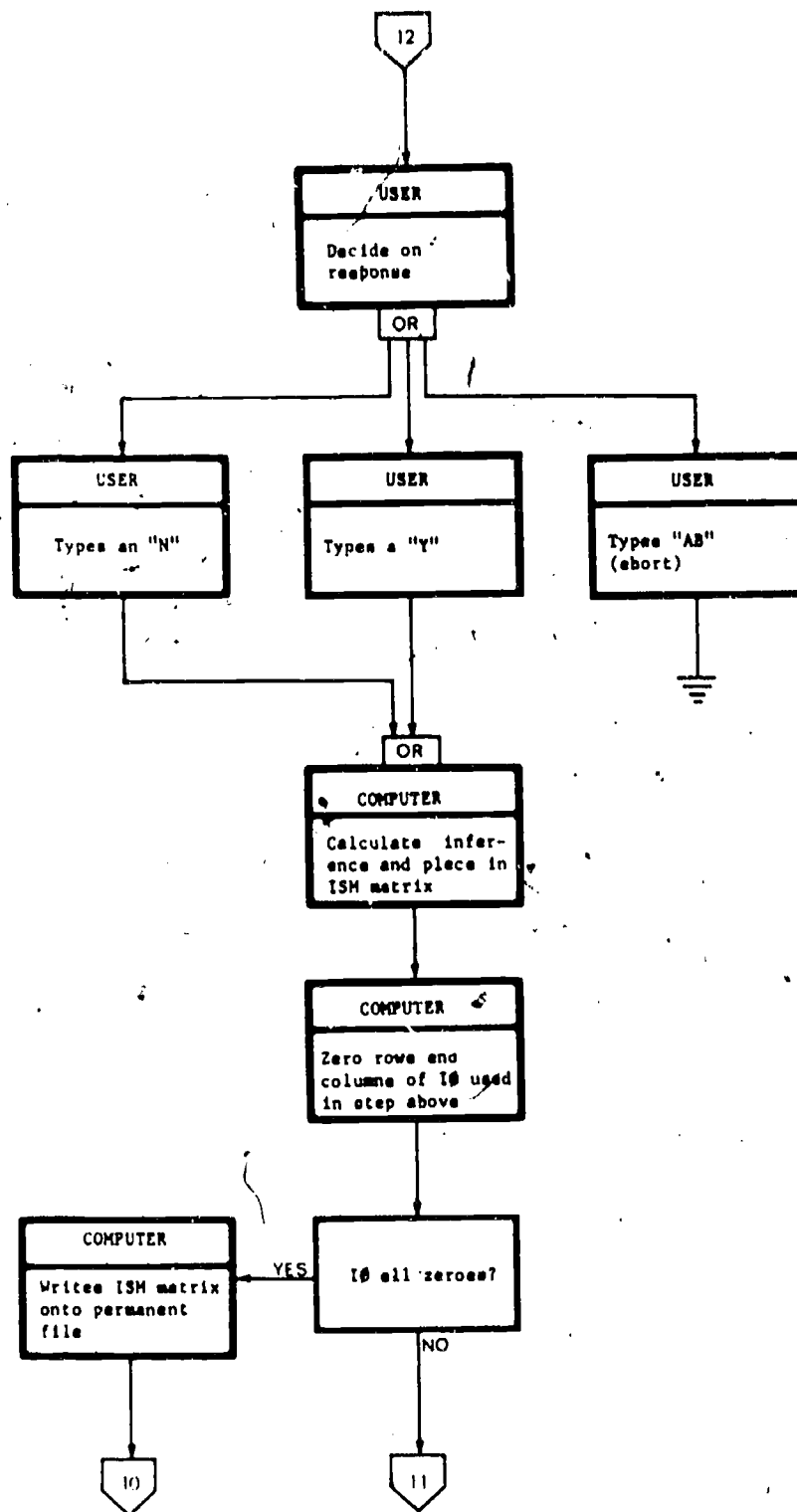
BOQ COMMAND





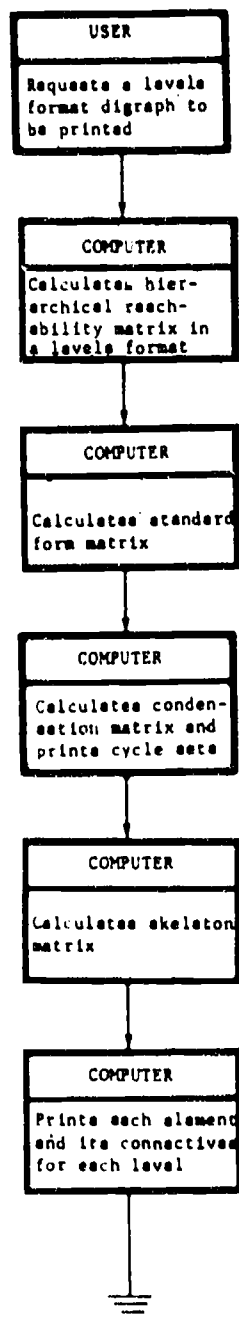
200



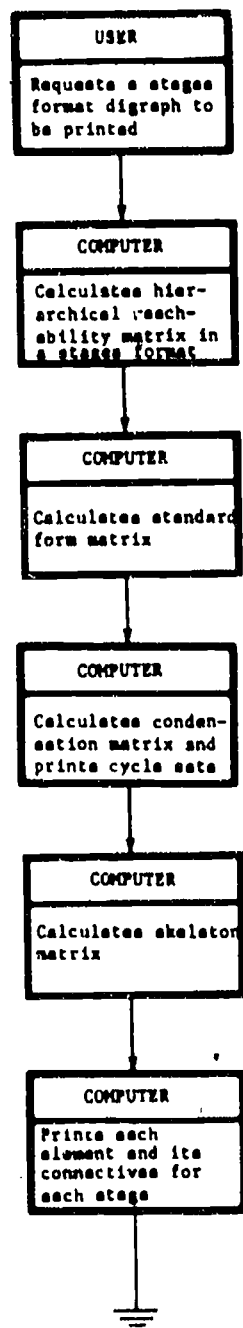


26

DI COMMAND

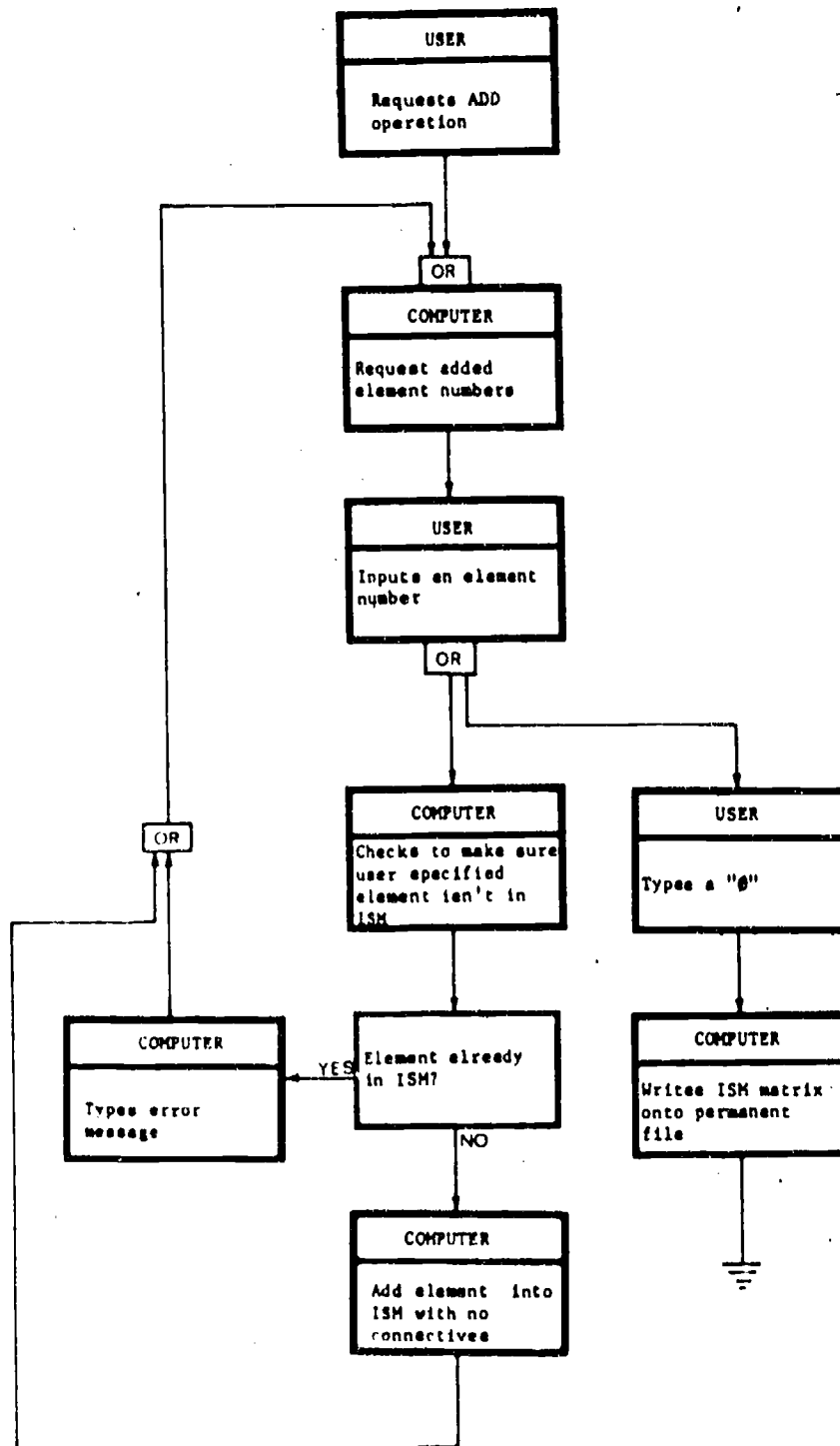


DIS COMMAND

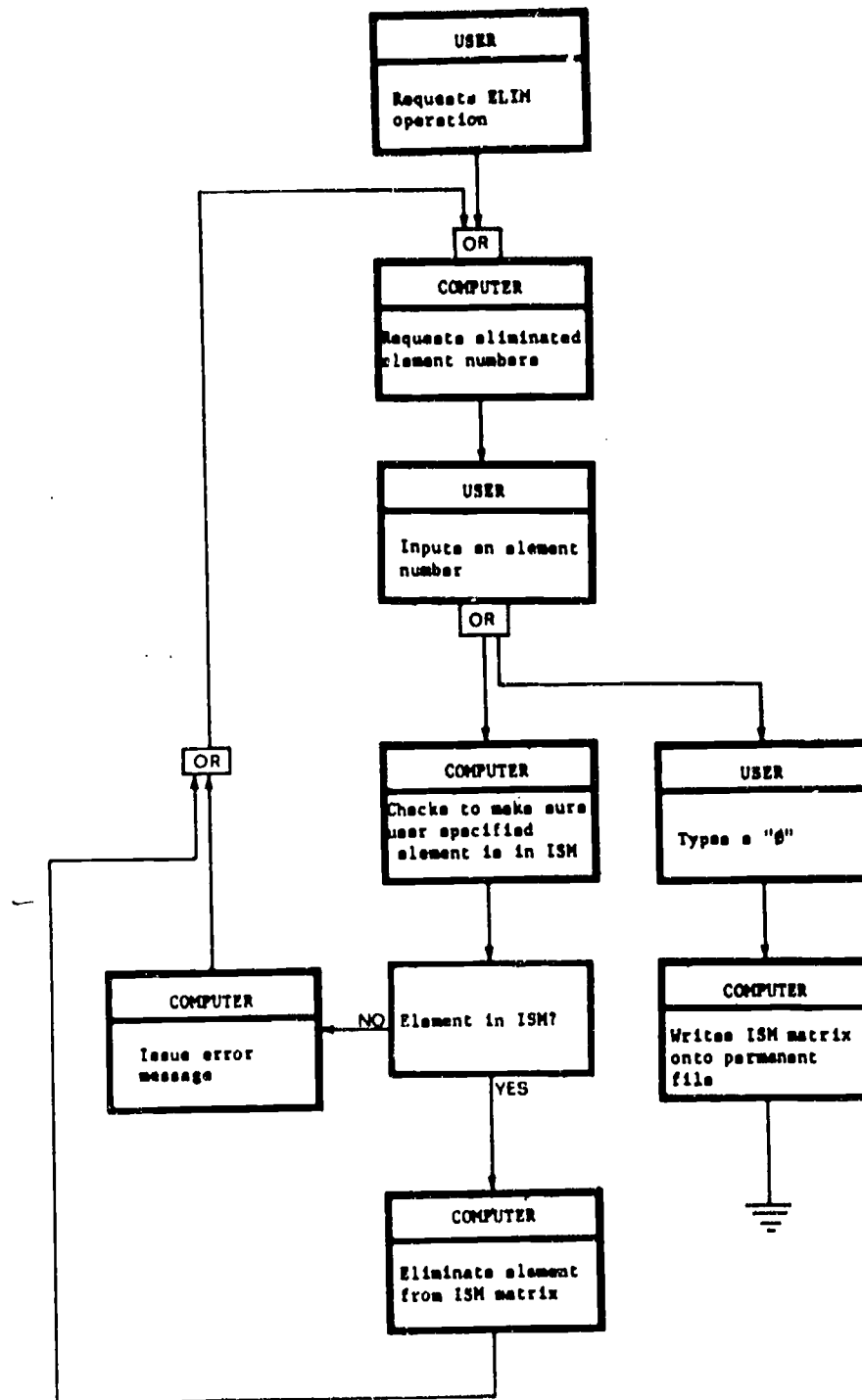


284

ADD COMMAND

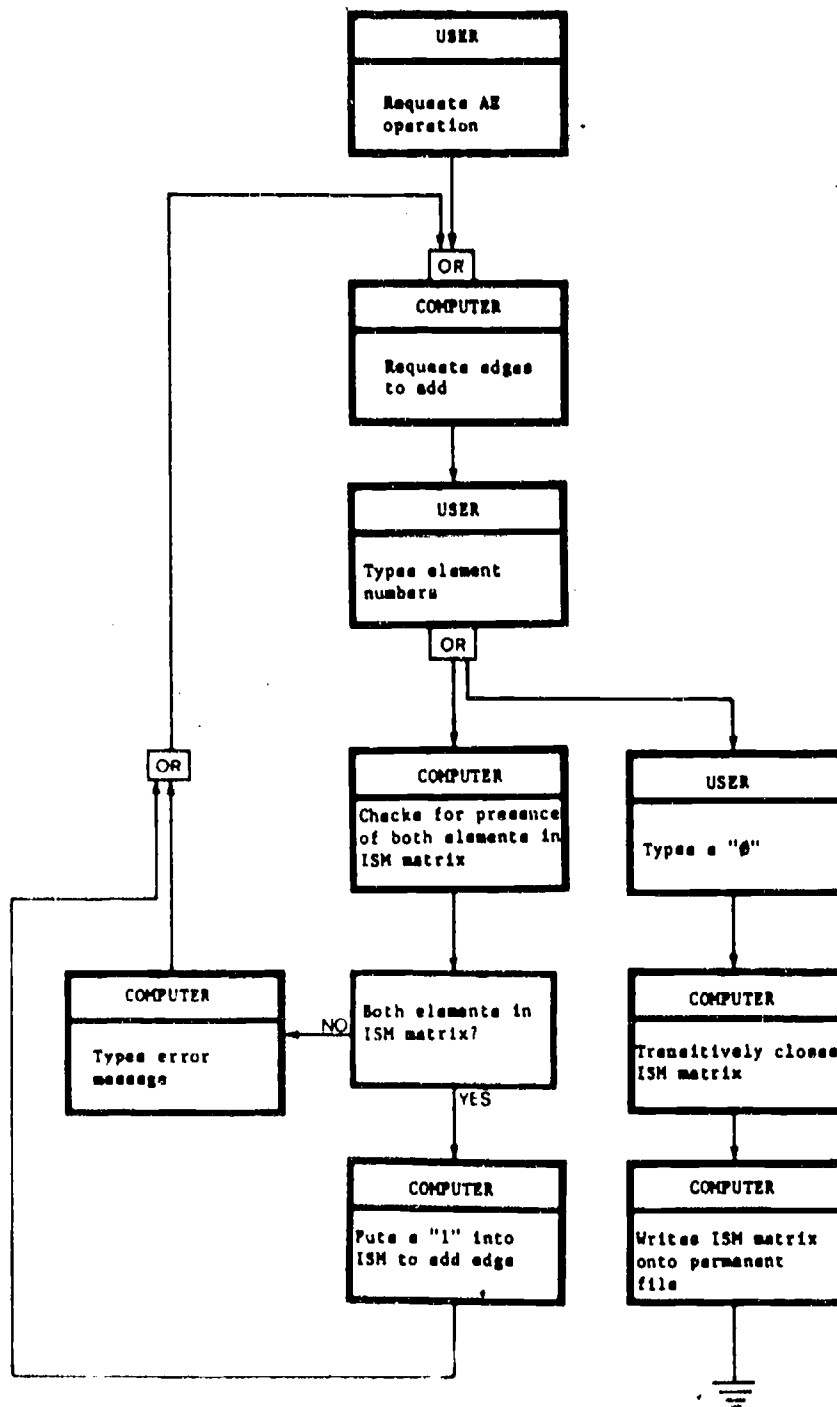


ELIM COMMAND

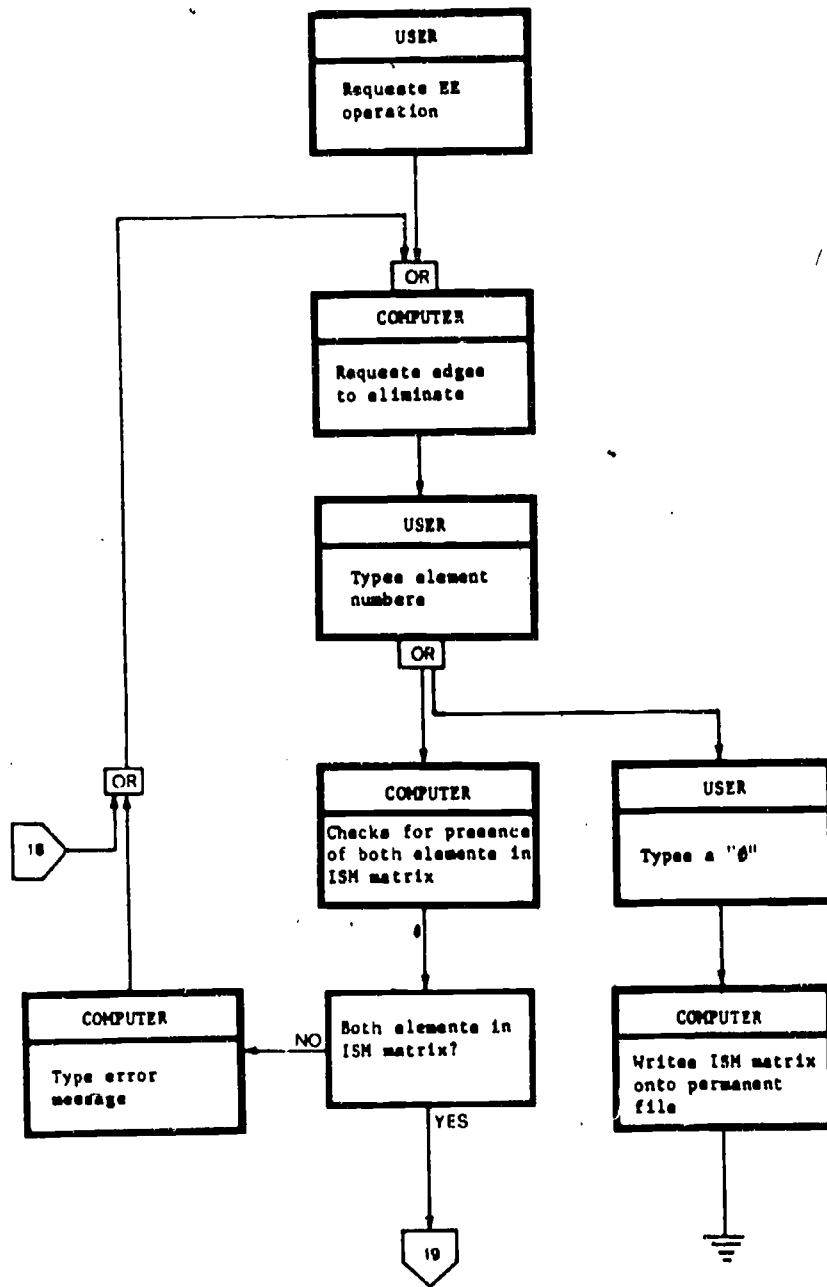


260

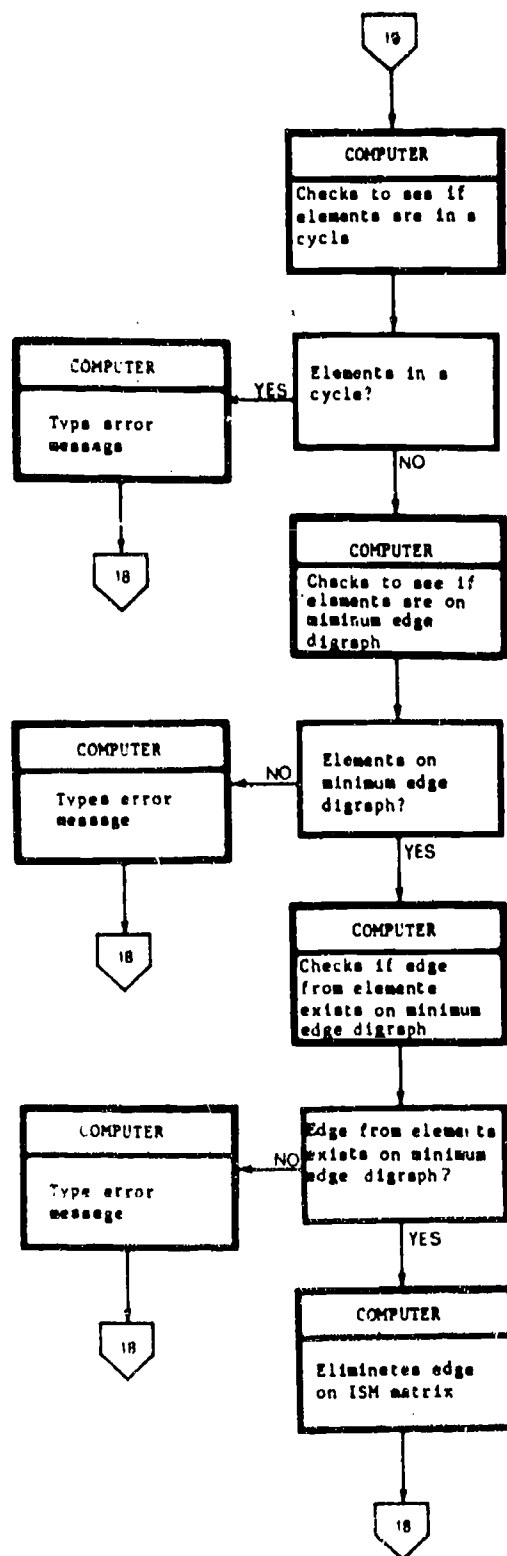
AE COMMAND



EE COMMAND

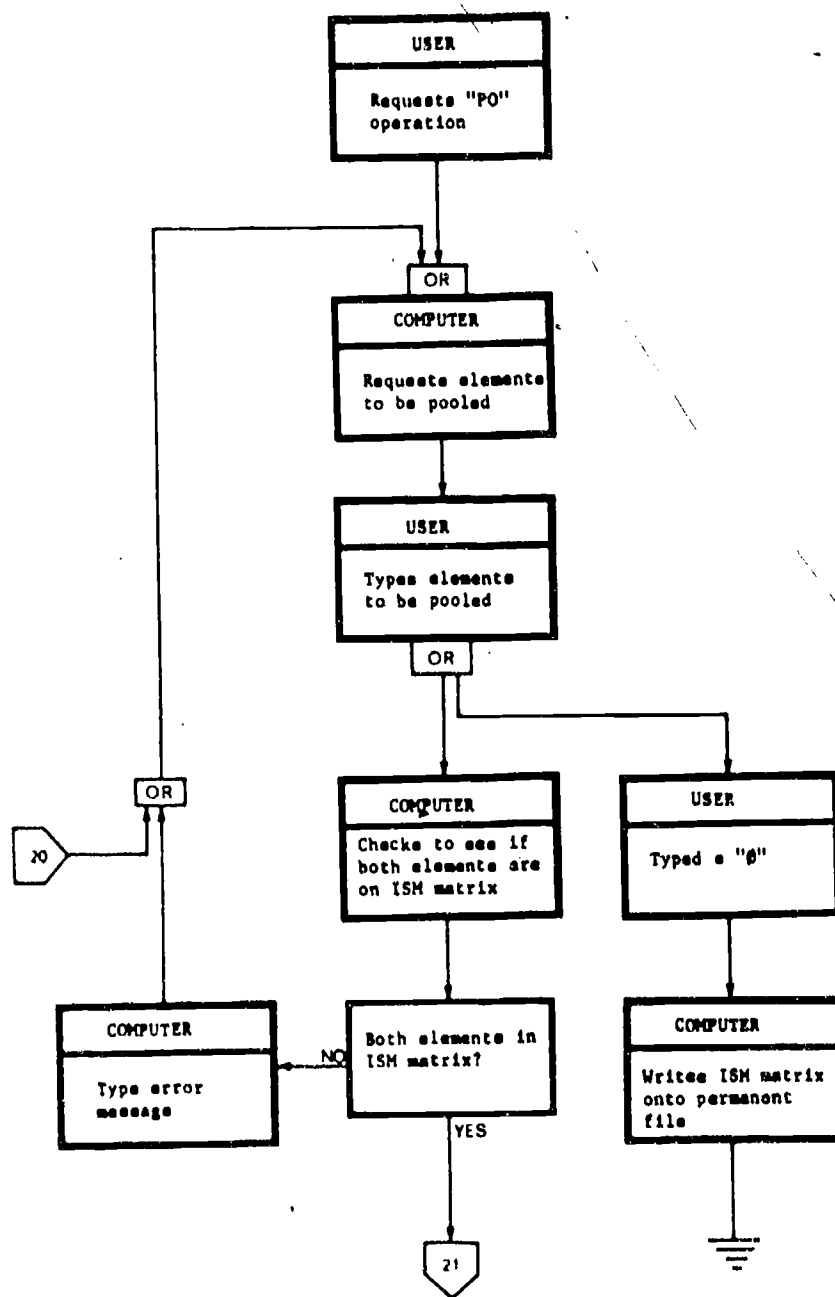


286

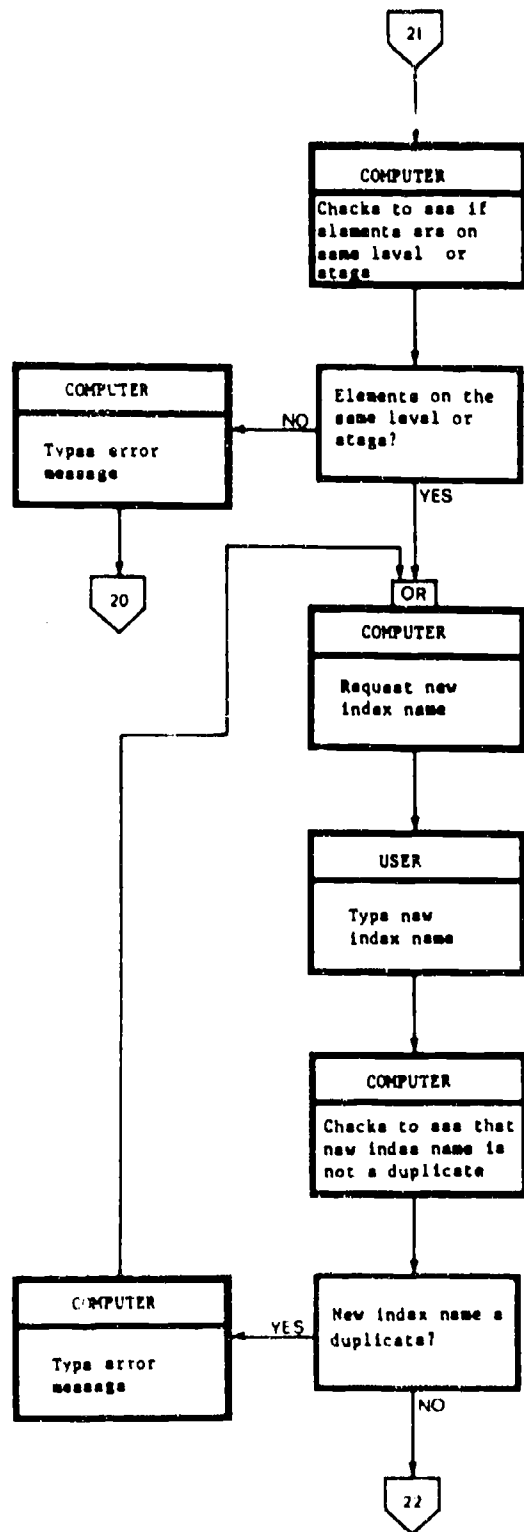


26J

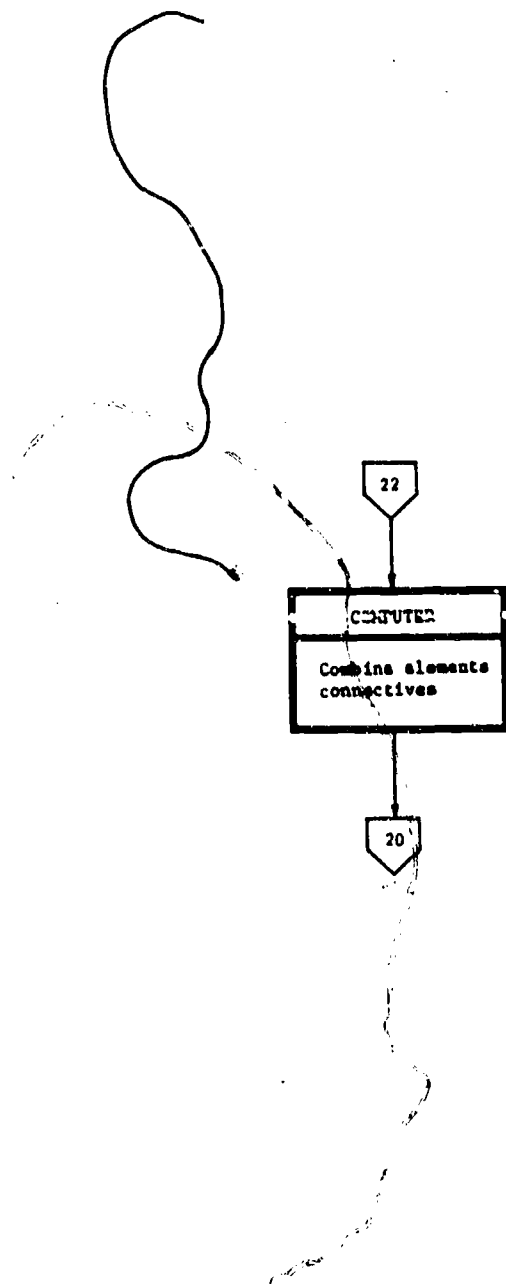
PO COMMAND



210

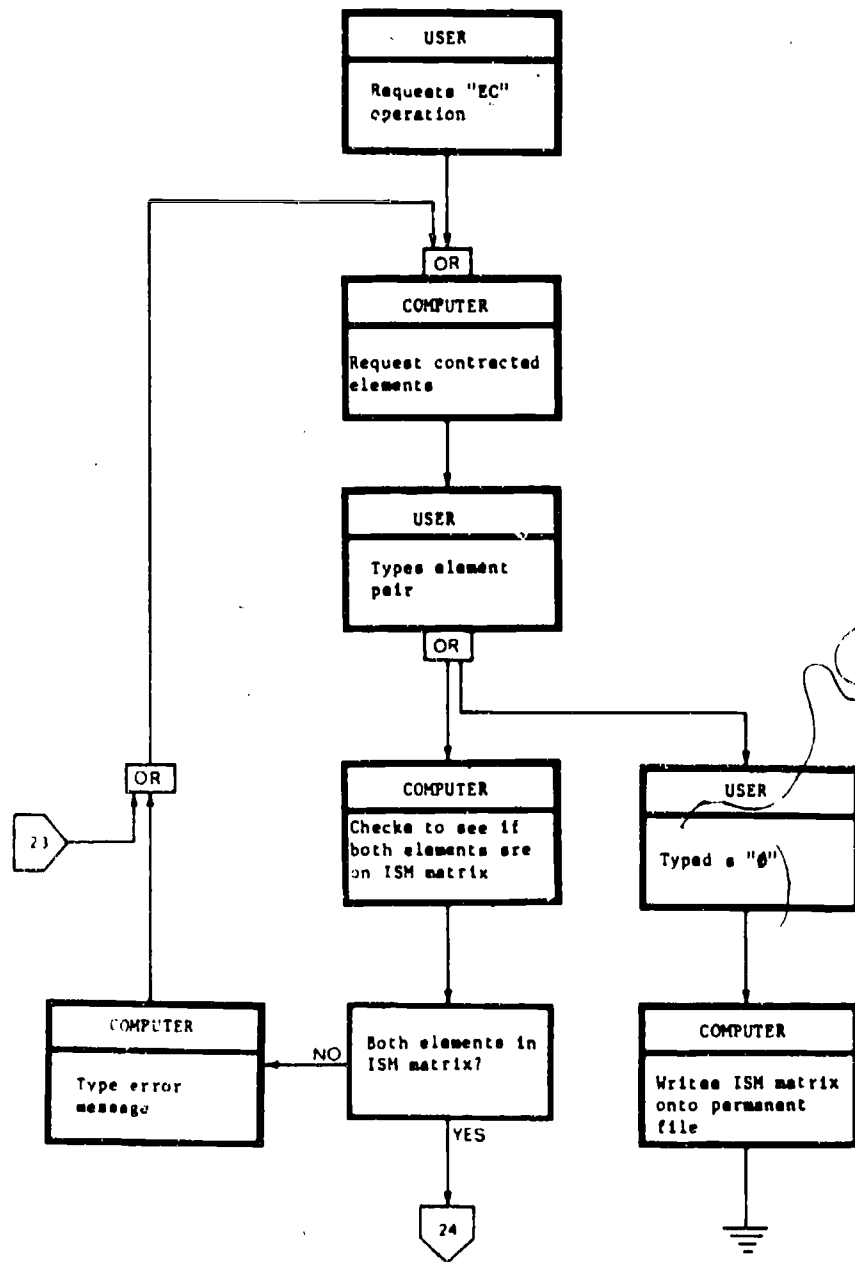


271

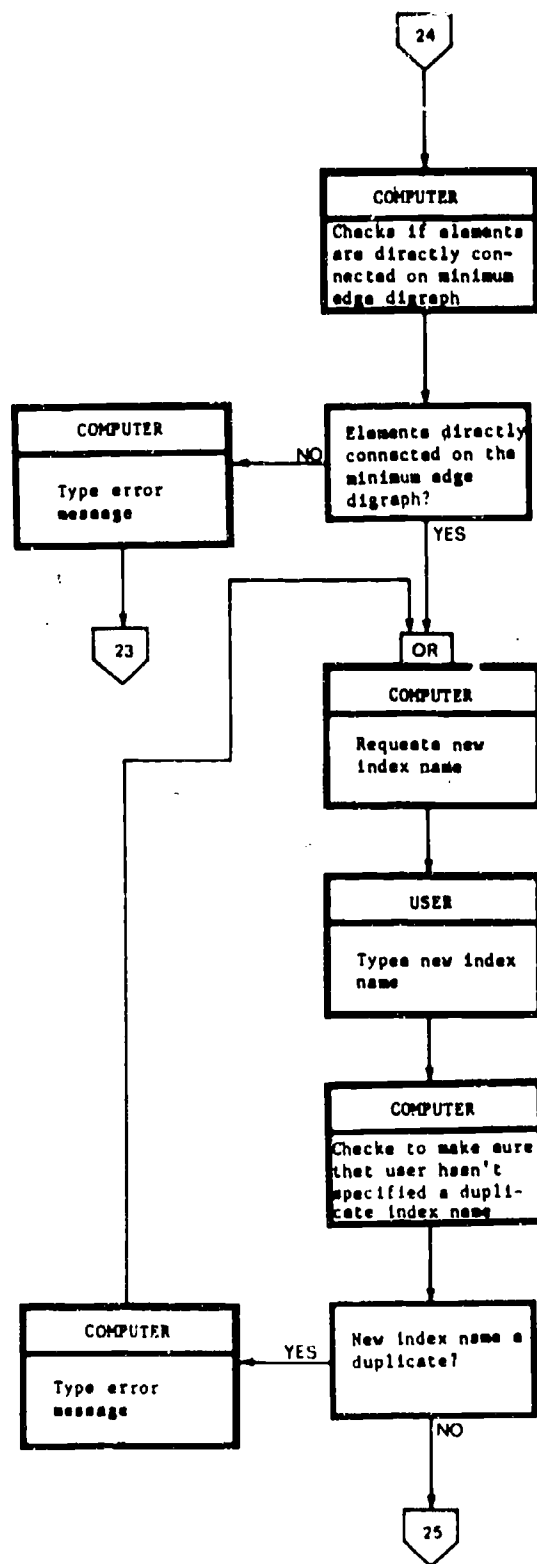


272

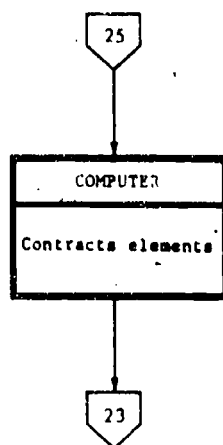
EC COMMAND



273

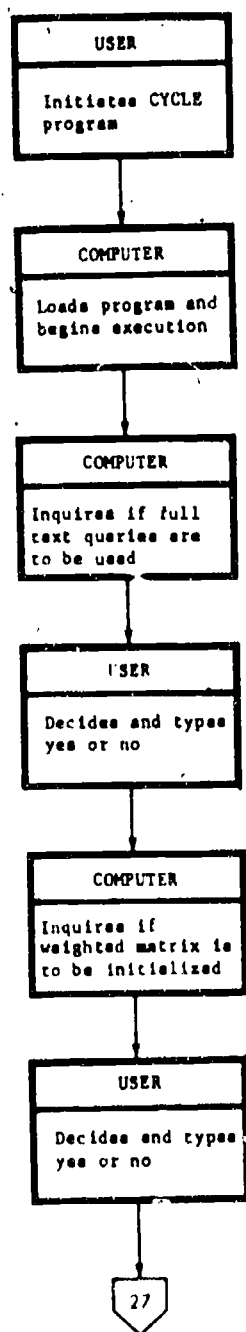


274

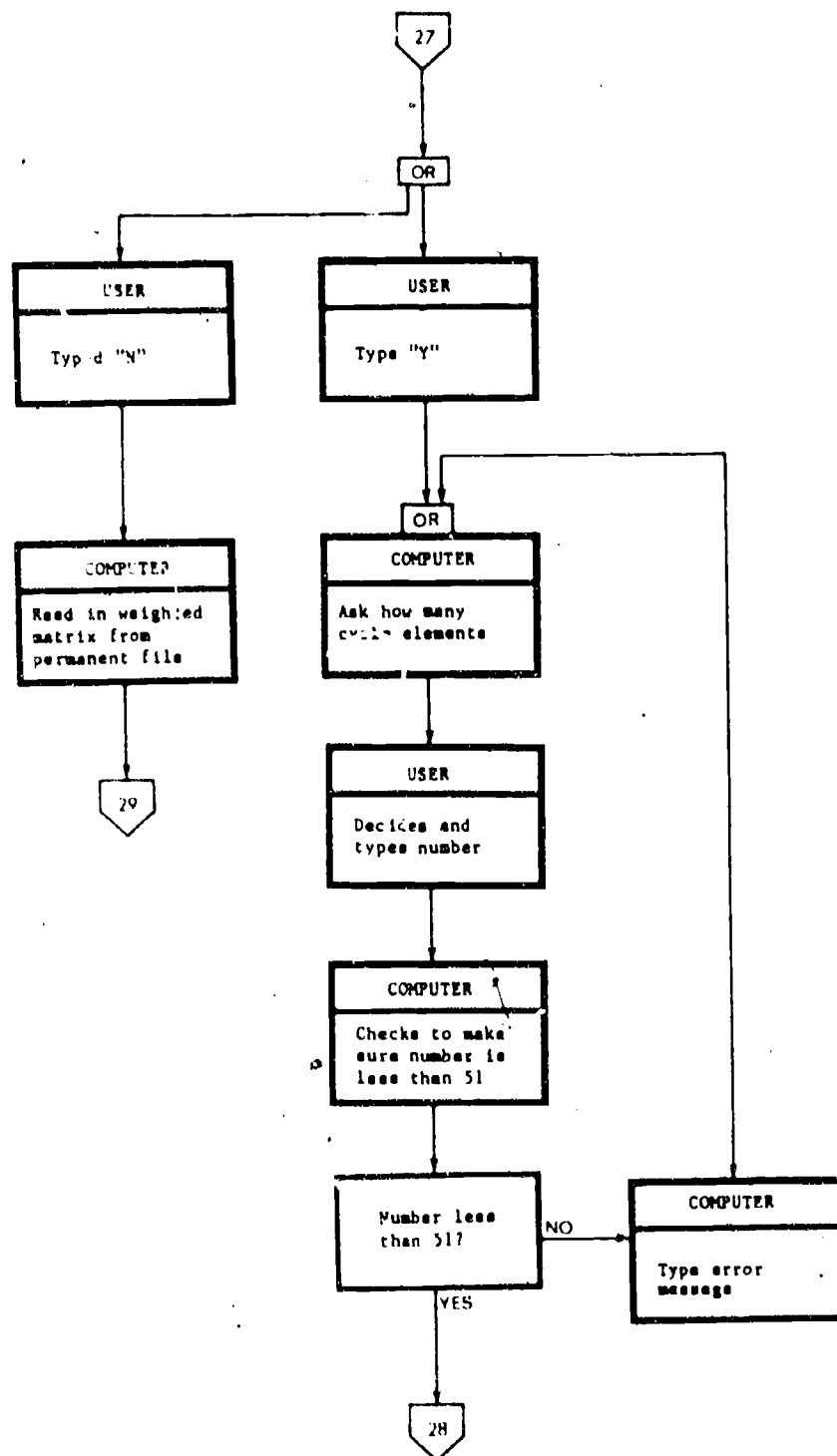


275

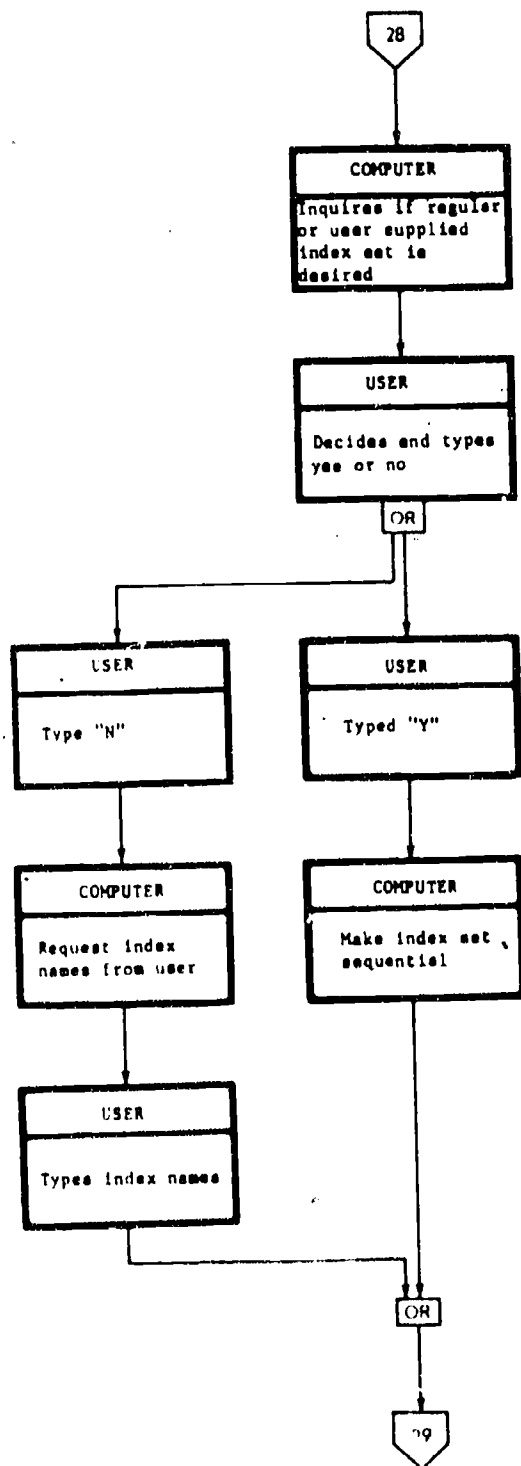
CYCLE PROGRAM



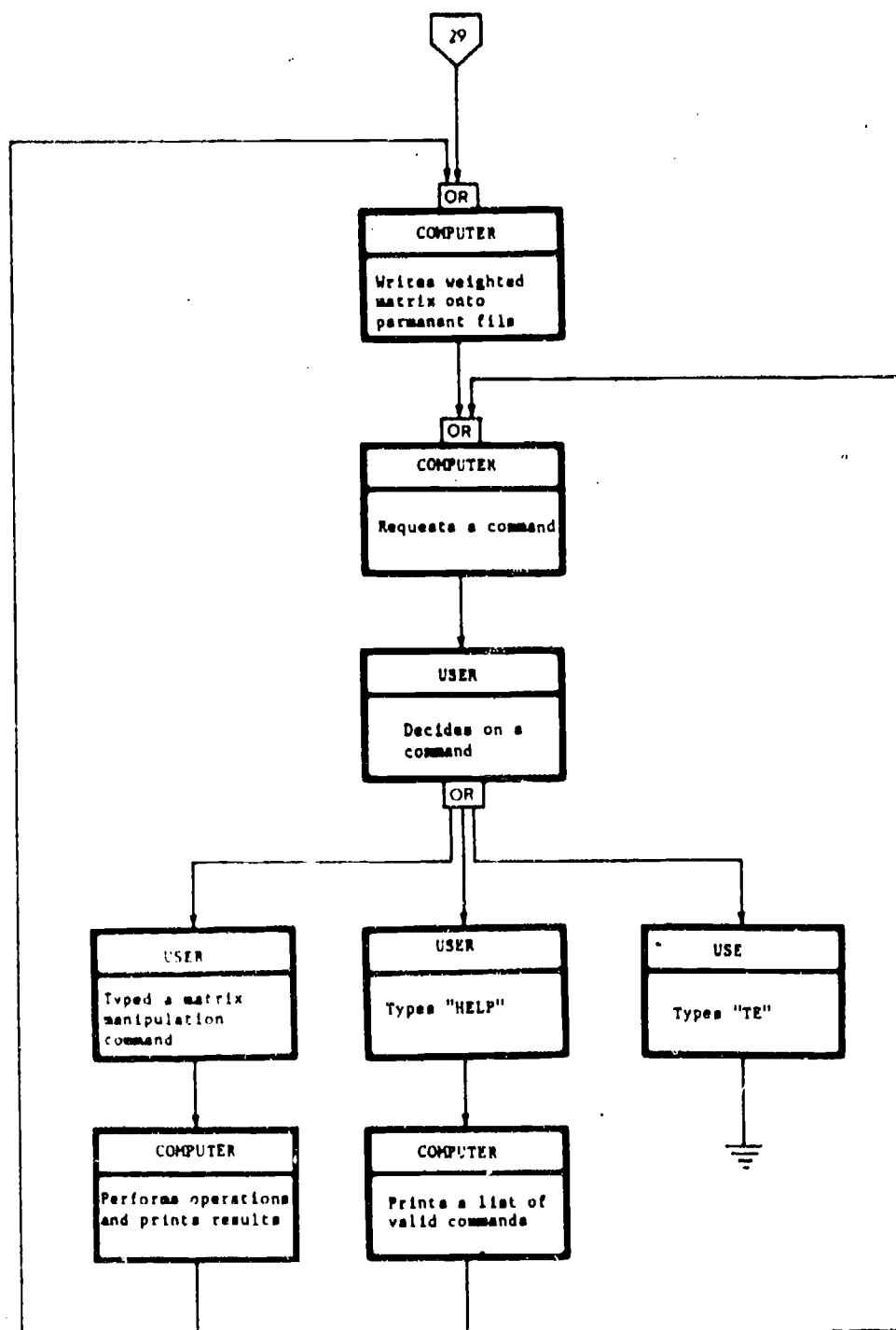
270



27.



270



270

Appendix to Volume 4:
Conducting Collective Inquiry
COMPUTER IMPLEMENTATION OF
INTERPRETIVE STRUCTURAL MODELING

(Part Two)

Written by:

Raymond L. Fitz, S.M.
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Joanne B. Troha
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University of Dayton

September 1979

Attachment No. 4

FORTTRAN PROGRAMS FOR ISMS

Prepared by:

David R. Yingling, Jr.

This attachment contains the listing of the FORTRAN IV programs entitled Interpretive Structural Modeling Software (ISMS). Tapes containing the listing of these programs can be obtained for a service fee by writing:

Director
The Office of Computing Activities
University of Dayton
Dayton, Ohio 45469

```

1 PROGRAM ISMSUD
2 C
3 C .....
4 C *
5 C *
6 C * ALL RIGHTS RESERVED. NO PART OF THIS PROGRAM MAY BE SOLD,
7 C * REPRODUCED, STORED IN A RETRIEVAL SYSTEM, OR TRANSMITTED
8 C * IN ANY FORM OR BY ANY MEANS, ELECTRONIC, MECHANICAL,
9 C * PHOTOCOPYING, RECORDING, OR OTHERWISE, WITHOUT THE
10 C * PRIOR PERMISSION OF THE
11 C * UNIVERSITY OF DAYTON RESEARCH INSTITUTE.
12 C *
13 C .....
14 C
15 C I S M S - U D
16 C INTERPRETIVE STRUCTURAL MODELING SOFTWARE - UNIVERSITY OF DAYTON
17 C
18 C WRITTEN BY: DAVID R. YINGLING, JR
19 C ENGINEERING AND PUBLIC POLICY GROUP
20 C UNIVERSITY OF DAYTON
21 C DAYTON, OHIO 45469
22 C
23 C VARIABLE NAME DESCRIPTION
24 C
25 C N THE NUMBER OF ELEMENTS IN THE
26 C REACHABILITY MATRIX "RM".
27 C
28 C RM THE CURRENT REACHABILITY MATRIX
29 C
30 C INDEX THE INDEX SET OF "RM".
31 C
32 C NUMS A VECTOR USED TO COMMUNICATE WITH
33 C SUBROUTINE GETNUM
34 C
35 C QTYPE LOGICAL VARIABLE WHICH WHEN .FALSE.,
36 C DENOTES FULL TEXT QUERIES, WHEN
37 C .TRUE., DENOTES SYMBOLIC QUERIES.
38 C
39 C SUBREL LOGICAL VARIABLE WHICH WHEN .TRUE.,
40 C DENOTES AN OPTIMIZATION OF THE
41 C TRANSITIVE BORDERING PROCESS IS TO
42 C BE DONE, WHEN .FALSE., THE
43 C OPTMIZATION IS NOT ABLE TO BE DONE.
44 C
45 C TTYIN FORTRAN READ UNIT NUMBER FOR TELETYPE
46 C
47 C TTYOUT FORTRAN WRITE UNIT NUMBER FOR TELETYPE
48 C
49 C
50 C TXTWDS THE NUMBER OF MACHINE WORDS REQUIRED
51 C TO HOLD ONE OF THE * PHRASES OF FULL
52 C TEXT QUERIES
53 C
54 C SYS DIMENSION SIZES OF SYSTEM MATRICES. IT
55 C IS THE THE MAXIMUM NUMBER OF ELEMENTS
56 C THAT CAN BE HANDELED BY THE PROGRAM.
57 C THIS VALUE IS SET BY THE PROGRAMMER
58 C DEPENDING ON THE AMOUNT OF MEMORY

```

```

59 C                                     AVAILABLE.
60 C
61 C ****COMMON BLOCK /FTEXT/
62 C
63 C     COMMON BLOCK /FTEXT/ IS USED WHEN FULL TEXT QUERIES
64 C     ARE SELECTED BY THE USER. IT CONSISTS OF 3 VECTORS
65 C     DIMENSIONED BY "TXTWDS". THE MAIN PROGRAM IS RESPONSIBLE
66 C     FOR READING IN THE "FRAMING" CLAUSES INTO THE BLOCK IN
67 C     THE POSITIONS R1, R2, AND R3.
68 C
69 C     IMPLICIT INTEGER*2 (A-Z)
70 C     LOGICAL QTYPE, SUBREL, FIRST, QST
71 C     LOGICAL*1 RM(128,128)
72 C     DIMENSION INDEX(128)
73 C     INTEGER BO,DI,ELIMM,HE,TE,YUP,DIS,AE,EE,ADD,PO,EC,ANSWER,R1,R2,R3
74 C     INTEGER L1,L2,UNUSED,BDQ,N,INDEX
75 C     DATA BO/2HBO/, DI/2HDI/, HE/4HHELP/, TE/2HTE/, YUP/1HY/
76 C     DATA DIS/3HDIS/, AE/2HAE/, EE/2HEE/, ELIMM/4HELIM/, ADD/3HADD/
77 C     DATA PO/2HPO/, EC/2HEC/, BDQ/3HBDQ/
78 C
79 C     COMMON /FTEXT/ R1(160), L1(160), R2(160), L2(160), R3(160)
80 C
81 C     THE BLANK COMMON IS FOR THE U.D. SYSTEM ONLY....
82 C     THE FORTRAN RUN TIME SUBROUTINES REQUIRE THE ASSOCIATED
83 C     VARIABLE OF THE DEFINE FILE STATEMENT IN BLANK COMMON
84 C
85 C     COMMON UNUSED
86 C     SYS      = 128
87 C     NCHAR    = 4
88 C     TXTWDS   = 160
89 C     TTYOUT   = 2
90 C     TTYIN    = 1
91 C     QTYPE    = .TRUE.
92 C     SUBREL   = .FALSE.
93 C     FIRST    = .TRUE.
94 C     QST      = .FALSE.
95 C
96 C     NEW SYSTEM ?
97 C
98 C     WRITE(TTYOUT,101)
99 C     READ (TTYIN, 200) ANSWER
100 C     IF(ANSWER .EQ. YUP) GOTO 3
101 C
102 C     SYSTEM MARTIX RESIDES ON A PERMFILE, READ IT IN
103 C
104 C     CALL IO(N,RM,INDEX,10,.TRUE.,SYS)
105 C     GOTO 3
106 C
107 C     UPDATE PERMFILE EVERY TIME IN CASE OF TIME
108 C     SHARING SYSTEM FAILURE
109 C
110 C     4 CALL IO(N,RM,INDEX,10,.FALSE.,SYS)
111 C
112 C     ASK THE USER FOR A ISMSUD COMMAND
113 C
114 C     3 WRITE(TTYOUT,103)
115 C     READ (TTYIN, 200) ANSWER
116 C

```

UDVER2

```

117 IF(ANSWER .EQ. BU) GOTO 5
118 IF(ANSWER .EQ. DI) GOTO 6
119 IF(ANSWER .EQ. DIS) GOTO 9
120 IF(ANSWER .EQ. ELIMM) GOTO 7
121 IF(ANSWER .EQ. HE) GOTO 8
122 IF(ANSWER .EQ. AE) GOTO 10
123 IF(ANSWER .EQ. EE) GOTO 11
124 IF(ANSWER .EQ. ADD) GOTO 12
125 IF(ANSWER .EQ. PD) GOTO 13
126 IF(ANSWER .EQ. EC) GOTO 14
127 IF(ANSWER .EQ. BOQ) GOTO 16
128 IF(ANSWER .EQ. TE) STOP
129 C
130 C INVALID COMMAND
131 C
132 WRITE(TTYOUT,104) ANSWER
133 GOTO 3
134 C
135 C *****
136 C * TRANSITIVE BORD RING *
137 C *****
138 C
139 5 IF(.NOT. FIRST) GOTO 15
140 FIRST = .FALSE.
141 C
142 C FULL TEXT QUERIES ?
143 C
144 WRITE(TTYOUT,100)
145 READ (TTYIN, 200) ANSWER
146 IF(ANSWER .NE. YUP) GOTO 1
147 C
148 C FULL TEXT QUERIES DESIRED
149 C
150 DEFINE FILE 8(999,160,U,UNUSED)
151 C
152 C *****NOTE*****
153 C THE RECORD SIZE FOR THE ABOVE DEFINE FILE STATEMENT SHOULD
154 C BE EQUAL TO "TXTWDS".
155 C
156 FIND(8'2)
157 QTYPE = .FALSE.
158 READ(8'2) (R1(1),I=1,TXTWDS)
159 READ(8'3) (R2(1),I=1,TXTWDS)
160 READ(8'4) (R3(1),I=1,TXTWDS)
161 C
162 C SUBORDINATION RELATION ?
163 C
164 1 WRITE(TTYOUT,106)
165 READ (TTYIN, 200) ANSWER
166 IF(ANSWER .EQ. YUP) SUBREL = .TRUE.
167 C
168 15 CALL BORDER(N,RH,INDEX,TTYIN,TTYOUT,QTYPE,SUBREL,QST,TXTWDS,SYS)
169 QST = .FALSE.
170 GOTO 3
171 C
172 C *****
173 C * DISPLAY DIGRAPH IN LEVELS *
174 C *****

```

```

175 C
176 6 CALL DIGLEV(N,RM,INDEX,TTYOUT,SYS)
177 GOTO 3
178 C
179 C *****
180 C *          DISPLAY      DIGRAPH      IN      STAGES          *
181 C *****
182 C
183 9 CALL DIGSTG(N,RM,INDEX,TTYOUT,SYS)
184 C
185 C RE-READ N,RM, AND INDEX BECAUSE STAGES ROUTINE
186 C HAS USED THEM FOR CRATCH AREAS
187 C
188 CALL IO(N,RM,INDEX,IO,.TRUE.,SYS)
189 GOTO 3
190 C
191 C *****
192 C *          ADD      AN      EDGE          *
193 C *****
194 C
195 10 CALL AEDGE(N,RM,INDEX,TTYIN,TTYOUT,SYS)
196 GOTO 4
197 C
198 C *****
199 C *          ERASE      AN      EDGE          *
200 C *****
201 C
202 11 CALL EEDGE(N,RM,INDEX,TTYIN,TTYOUT,SYS)
203 GOTO 3
204 C
205 C *****
206 C *          ELIMINATE      AN      ELEMENT          *
207 C *****
208 C
209 7 CALL ERASE(N,RM,INDEX,TTYIN,TTYOUT,SYS)
210 GOTO 4
211 C
212 C *****
213 C *          ADD      AN      ELEMENT          *
214 C *****
215 C
216 12 CALL ADDEL(N,RM,INDEX,TTYIN,TTYOUT,SYS)
217 GOTO 4
218 C
219 C *****
220 C *          POOL      ELEMENTS          *
221 C *****
222 C
223 13 CALL POOL(N,RM,INDEX,TTYIN,TTYOUT,SYS)
224 GOTO 4
225 C
226 C *****
227 C *          CONTRACT      ELEMENTS          *
228 C *****
229 C
230 14 CALL ELCONT(N,RM,INDEX,TTYIN,TTYOUT,SYS)
231 GOTO 4
232 C

```

```

233 .....
234 C * BORDER WITH YOUR OWN QUERIES *
235 C .....
236 16 QST = .TRUE.
237 GOTO 5
238 C
239 C HELP MESSAGE
240 C
241 8 WRITE(TTYOUT,105)
242 GOTO 3
243 C
244 C FORMATS
245 C
246 100 FORMAT(34H-FULL TEXT QUERIES DESIRED ? (Y/N))
247 101 FORMAT(19H-NEW SYSTEM ? (Y/N))
248 103 FORMAT(31H-TYPE ISMS-UD CMD ? (OR "HELP"))
249 104 FORMAT(14H-***ERROR*** ,A4,27H <--INVALID ISMS-UD COMMAND/45H IF
250 +IN DOUBT TYPE "HELP" FOR LIST OF COMMANDS)
251 105 FORMAT(1H-/1H-,5X,34HI S M S - U D C O M M A N D S/1H-,3X,
252 +8HEMBEDDING/1H0,4X,32HBD - TRANSITIVE BORDERING METHOD/
253 +4X,50HBDQ - TRANSITIVE BORDERING WITH SELECTABLE QUERIES/1H-,3X,
254 +10HDISPLAYING/1H0,4X,33HDI - DISPLAY MINIMUM EDGE DIGRAPH/10X,
255 +18HIN A LEVELS FORMAT/1H0,3X,34HDI - DISPLAY MINIMUM EDGE DIGRAPH
256 +/10X,18HIN A STAGES FORMAT/1H-,3X,20HSUBSTANTIVE AMENDING/1H0,3X,
257 +18HADD - ADD ELEMENTS/1H0,2X,29HELIM - ELIMINATE ELEMENTS/1H0,4X,
258 +30HAE - ADD EDGES (RELATIONSHIPS)/10X,27HON THE MINIMUM EDGE DIGRA
259 +PH/1H0,4X,32HEE - ERASE EDGES (RELATIONSHIPS)/10X,27HON THE MINIMU
260 +M EDGE DIGRAPH/1H-,3X,15HFORMAT AMENDING/1H0,4X,18HPO - POOL ELEME
261 +NTS/1H0,4X,27HEC - ELEMENTARY CONTRACTION/1H-,3X,11HINTERMINATION/
262 +1H0,4X,30HTE - TERMINATE ISMS-UD PROGRAM/11H-***NOTE***/1H0,2X,
263 +26HHELP - REPRINTS ABOVE LIST)
264 106 FORMAT(32H-SUBORDINATION RELATION ? (Y/N) )
265 200 FORMAT(A4)
266 END

```



```

1  SUBROUTINE BURDER(N,MAT,INDEX,TTYIN,TTYOUT,QTYPE,SUBREL,QST,
2  +TXTWDS,SYS)
3  C
4  C
5  C  THIS SUBROUTINE WILL EMBED AN ELEMENT INTO "MAT"
6  C  UTILIZING THE TRANSITIVE BORDERING ALGORITHM DEVELOPED
7  C  BY DR. JOHN WARFIELD.
8  C
9  C
10 C  WRITTEN BY:  DAVID R. YINGLING, JR.
11 C              ENGINEERING AND PUBLIC POLICY GROUP
12 C              UNIVERSITY OF DAYTON
13 C              DAYTON, OHIO 45469
14 C
15 C  VARIABLE NAME      DESCRIPTION
16 C
17 C  N                  THE NUMBER OF ELEMENTS IN "MAT". UPON
18 C                    COMPLETION OF THIS SUBROUTINE,
19 C                    "N" = "N" + 1 .
20 C
21 C  MAT                THE CURRENT REACHABILITY MATRIX MODEL
22 C
23 C  INDEX              INDEX SET OF "MAT"
24 C
25 C  TTYIN              FORTRAN READ UNIT NUMBER FOR TELETYPE
26 C
27 C  TTYOUT             FORTRAN WRITE UNIT NUMBER FOR TELETYPE
28 C
29 C  QTYPE              THE QUERY TYPE SWITCH. IF QTYPE =
30 C                    .TRUE., SYMBOLIC QUERIES ARE USED,
31 C                    IF QTYPE = .FALSE., FULL TEXT QUERIES
32 C                    ARE USED.
33 C
34 C  SUBREL             THE SUBORDINATION RELATIONSHIP SWITCH.
35 C                    IF THE EMBEDDING RELATIONSHIP IS
36 C                    SUBORDINATE, AN OPTMIZATION CAN BE
37 C                    MADE THAT REDUCES THE NUMBER OF
38 C                    QUESTIONS REQUIRED TO EMBEDD
39 C                    ELEMENTS. IF .FALSE., THIS
40 C                    OPTMIZATION IS NOT DONE.
41 C
42 C  QST                LOGICAL VARIABLE IF WHEN .TRUE.,
43 C                    THE USER IS ABLE TO SELECT HIS
44 C                    OWN QUESTIONS. WHEN .FALSE.,
45 C                    QUESTIONING IS AUTOMATIC
46 C
47 C
48 C  TXTWDS             THE NUMBER OF WORDS REQUIRED TO
49 C                    HOLD ONE OF THE 5 PHRASES OF
50 C                    FULL TEXT QUERIES.
51 C
52 C  SYS                DIMENSION SIZES OF SYSTEM MATRICES
53 C
54 C  -----
55 C
56 C  FLAG              DENOTES ROWS/COLS OF "PHI" WHICH HAVE
57 C                    BEEN ZEROED OUT AS A RESULT OF THE
58 C                    SEARCH/ZERO OUT ROUTINES

```

```

59 C
60 C      PHI                      THE TRANSITIVE BOUNDING INFERENCE
61 C                                     OPPORTUNITY MATRIX
62 C
63 C      Z                        SCRATCH VECTOR USED FOR PHI ZERO OUT
64 C                                     ROUTINE
65 C
66 C      DIMPHI                   THE DIMENSION OF THE CURRENT
67 C                                     PHI MATRIX
68 C
69 C      ZZZ                      LOGICAL SWITCH VARIABLE
70 C
71 C
72 C      ****NOTE****
73 C      THE DIMENSION SIZES OF "PHI", "Z", AND "FLAG" SHOULD BE
74 C      SET EQUAL TO 2*"SYS".
75 C
76 C      IMPLICIT INTEGER*2 (A-Z)
77 C      INTEGER  NOPE,YUP,ABORT,ANSWER,R1,R2,R3,L1,L2,RECORD,UNUSED,N
78 C      INTEGER INDEX, K
79 C      LOGICAL QTYPE, SUBREL, FOUND1, FOUND2, QST, QST1
80 C      LOGICAL*1 FLAG(256), Z(256), PHI(256,256), MAT(SYS,SYS), ZZZ
81 C      DIMENSION INDEX(SYS), NUMS(2)
82 C      DATA NOPE/1HN/, YUP/1HY/, ABORT/2HAB/
83 C
84 C      COMMON /FTEXT/ R1(160), L1(160), R2(160), L2(160), R3(160)
85 C
86 C      THIS BLANK COMMON IS FOR THE U.D. SYSTEM ONLY
87 C
88 C      COMMON UNUSED
89 C
90 C      INITIALIZE EVERYTHING
91 C
92 C      QST1 = QST
93 C      32 JIMPHI = N * 2
94 C      QST = QST1
95 C
96 C      DO 1 I=1,DIMPHI
97 C      FLAG(I) = .FALSE.
98 C      Z(I) = .FALSE.
99 C      DO 1 J=1,DIMPHI
100 C      PHI(I,J) = .FALSE.
101 C      1 CONTINUE
102 C
103 C      MAKE SURE SIZE OF ""SYS"" ELEMENTS IS NOT EXCEEDED
104 C
105 C      IF(N + 1 .LE. SYS) GOTO 2
106 C      WRITE(TTYOUT,103) SYS
107 C      GOTO 3
108 C
109 C      ASK INTERACTIVE USER FOR NEW ELEMENT NUMBER
110 C
111 C      2 WRITE(TTYOUT,100)
112 C      CALL GETNUM(NUMS,1,TTYIN,TTYOUT)
113 C      INDEX(N + 1) = NUMS(1)
114 C      IF(INDEX(N + 1) .EQ. 0) GOTO 3
115 C      IF(N .EQ. 0) GOTO 35
116 C

```

UDVER2

```

117 C      CHECK TO SEE IF USER SPECIFIED AN ELEMENT ALREADY
118 C      IN THE SYSTEM MATRIX
119 C
120      CALL FINDIT(N,INDEX(N + 1),J,1,J,INDEX,FOUND1,FOUND2,SYS)
121      IF(FOUND1) GOTO 5
122 35 IF(QTYPE) GOTO 6
123 C
124 C      FULL TEXT OPTION ON. MAKE SURE THIS
125 C      TEXT RECORD IS ON RANDOM TEXT FILE
126 C
127      RECORD = INDEX(N + 1) + 4
128      READ(8'RECORD,ERR=36) (L1(I),I=1,XTWDS)
129      GOTO 6
130 36 WRITE(TTYOUT,102) INDEX(N + 1)
131      GOTO 2
132 C
133 C      ALL OK, BEGIN XITIVE BORDERING
134 C      1. FORM PHI MATRIX
135 C
136 C
137 C          A          0
138 C
139 C      PHI =
140 C          -T
141 C          AA A          A
142 C
143 6 CONTINUE
144      IF(N .EQ. 0) GOTO 33
145      SYS2 = 2 * SYS
146      CALL TBPHI(N,MAT,PHI,DIMPHI,SUBREL,SYS2,SYS)
147 C
148 C      WE NOW HAVE AN INFERENCE OPPORTUNITY MATRIX
149 C
150 C      2. DO MIN/MAX PROCEDURE TO FIND THE Z (ZPOINT) WITH
151 C      THE MOST INFERENCE.
152 C
153 7 CONTINUE
154      IF(QST) GOTO 37
155      CALL FINDZ(DIMPHI,PHI,ZPOINT,FLAG,SYS2)
156 43 CONTINUE
157 C
158 C      SINCE WE NOW HAVE THE Z (ZPOINT) WITH THE MAXIMUM INFERENCE
159 C      POTENTIAL, ASK THE USER THE RELATION OF THIS ELEMENT WITH THE
160 C      ADDED ONE.
161 C
162 C      IF ZPOINT IS GREATER THEN N, A "Y" QUESTION IS TO BE ASKED
163 C
164 C      IF ZPOINT IS LESS THAN OR EQUAL TO N, AN "X" QUESTION
165 C      IS TO BE ASKED
166 C
167      RELATE = ZPOINT
168      IF(ZPOINT .GT. N) GOTO 8
169 C
170 C      ASK THE "X" QUESTION AND CHECK
171 C
172 9 CALL QUEST(INDEX(RELATE),INDEX(N + 1),TTYOUT,QTYPE,XTWDS)
173      READ(TTYIN,200) ANSWER
174      WRITE(TTYOUT,106) ANSWER

```

```

175      IF(ANSWER .EQ. ABORT) GOTO 31
176      IF(ANSWER .NE. YUP .AND. ANSWER .NE. NOPE) GOTO 10
177      ZZZ = .TRUE.
178      IF(ANSWER .EQ. YUP) GOTO 11
179      GOTO 12
180      11 ZZZ = .FALSE.
181      GOTO 13
182 C
183 C      ERROR MESSAGE
184 C
185      10 WRITE(TTYOUT,104)
186      GOTO 9
187 C
188 C      ASK THE "Y" QUESTION AND CHECK
189 C
190      8 RELATE = ZPOINT - N
191      14 CALL QUEST(INDEX(N + 1),INDEX(RELATE),TTYOUT,QTYPE,XTWDS)
192      READ(TTYIN,200) ANSWER
193      WRITE(TTYOUT,106) ANSWER
194      IF(ANSWER .EQ. ABORT) GOTO 31
195      IF(ANSWER .NE. YUP .AND. ANSWER .NE. NOPE) GOTO 15
196      ZZZ = .FALSE.
197      IF(ANSWER .EQ. YUP) GOTO 16
198      GOTO 12
199      16 ZZZ = .TRUE.
200      GOTO 13
201 C
202 C      ERROR MESSAGE
203 C
204      15 WRITE(TTYOUT,104)
205      GOTO 8
206 C
207 C      COME HERE WHEN THE QUESTION WAS ANSWERED NO.
208 C
209 C      IF ZZZ = .TRUE., "X" QUESTION ASKED-SEARCH ROW OF PHI
210 C
211 C      IF ZZZ = .FALSE., "Y" QUESTION ASKED-SEARCH COL OF PHI
212 C
213      12 IF(.NOT. ZZZ) GOTO 17
214      GOTO 18
215 C
216 C      COME HERE WHEN THE QUESTION WAS ANSWERED YES
217 C
218 C      IF ZZZ = .FALSE., "X" QUESTION ASKED-SEARCH COL OF PHI
219 C
220 C      IF ZZZ = .TRUE., "Y" QUESTION ASKED-SEARCH COL OF PHI
221 C
222      13 IF(.NOT. ZZZ) GOTO 17
223 C
224 C      SEARCH ROW OF PHI FOR INFERENCE TO BE ENTERED INTO "MAT"
225 C
226      18 DO 19 J=1,DIMPHI
227      IF(.NOT. PHI(ZPOINT,J)) GOTO 19
228      ICTR2 = J
229      IF(ICTR2 .GT. N) ICTR2 = ICTR2 - N
230 C
231      IF(J .LE. N) GOTO 20
232 C

```

```

233      MAT(N + 1,ICTR2) = .TRUE.
234      GOTO 21
235 C
236      20 MAT(ICTR2,N + 1) = .FALSE.
237 C
238      21 FLAG(J) = .TRUE.
239      Z(J) = .TRUE.
240      19 CONTINUE
241      GOTO 25
242 C
243 C      SEARCH COL OF PHI FOR INFERENCE TO BE ENTERED INTO "MAT"
244 C
245      17 DO 22 I=1,DIMPHI
246          IF(.NOT. PHI(I,ZPOINT)) GOTO 22
247          ICTR2 = I
248          IF(ICTR2 .GT. N) ICTR2 = ICTR2 - N
249 C
250          IF(I .LE. N) GOTO 23
251 C
252          MAT(N + 1,ICTR2) = .FALSE.
253          GOTO 24
254 C
255      23 MAT(ICTR2,N + 1) = .TRUE.
256 C
257      24 FLAG(I) = .TRUE.
258      Z(I) = .TRUE.
259      22 CONTINUE
260 C
261 C      PHI MATRIX ZERO OUT ROUTINE
262 C
263      25 DO 26 I=1,DIMPHI
264          IF(.NOT. Z(I)) GOTO 26
265 C
266          DO 27 J=1,DIMPHI
267              PHI(I,J) = .FALSE.
268              PHI(J,I) = .FALSE.
269      27 CONTINUE
270      26 CONTINUE
271 C
272 C      TEST FLAG VECTOR FOR ALL .TRUE.
273 C      IF SO, ALL ROWS/COLS OF PHI ARE ZERO-ALGORITHM
274 C      IS COMPLETE-RETURN
275 C
276 C      IF NOT ALL .TRUE., ZERO OUT NECESSARY VECTORS
277 C      AND GO AGAIN.
278 C
279 C
280      DO 28 I=1,DIMPHI
281          IF(.NOT. FLAG(I)) GOTO 29
282      28 CONTINUE
283 C
284 C      COME HERE WHEN PHI IS NULL
285 C
286      33 N = N + 1
287      MAT(N,N) = .TRUE.
288 C
289 C      SAVE NEW MATRIX ON PERMFILE IN CASE OF TIMESHARING FAILURE
290 C

```

```

291      CALL ID(1,MAT,INDEX,10,.FALSE.,SYS)
292 C
293 C      KEEP BORDERING UNTIL USER TYPES A "0" FOR NEW ELEMENT NUMBER
294 C
295      GOTO 32
296 3 RETURN
297 C
298 C      ERROR MESSAGE
299 C
300 5 WRITE(TTYOUT,101) INDEX(N + 1)
301      GOTO 2
302 C
303 C      PHI STILL HAS SOME UNES IN IT
304 C
305 29 DO 30 I=1,DIMPHI
306     Z(I) = .FALSE.
307 30 CONTINUE
308 C
309 C      GO AGAIN !
310 C
311      GOTO 7
312 C
313 C      COME HERE WHEN USER HAS ABORTED BORDERING SEQUENCE
314 C
315 31 WRITE(TTYOUT,105) INDEX(N + 1)
316      GOTO 3
317 C
318 C      *****
319 C      ASK YOUR OWN QUESTIONS CODE
320 C      *****
321 C
322 37 WRITE(TTYOUT,107)
323      CALL GETNUM(NUMS,2,TTYIN,TTYOUT)
324 C
325 C      RESET TO AUTOMATIC QUERIES ?
326 C
327      IF(NUMS(1) .LE. 0 .OR. NUMS(2) .LE. 0) GOTO 41
328 C
329 C      ONE OF NUMBERS TYPES MUST BE EQUAL TO INDEX(N + 1)
330 C
331      IF(NUMS(1) .NE. INDEX(N + 1) .AND. NUMS(2) .NE. INDEX(N + 1))
332 +GOTO 38
333 C
334 C      NEXT CHECK IS JUST IN CASE USER TRIES SOME FUNNY STUFF
335 C
336      IF(NUMS(1) .EQ. INDEX(N + 1) .AND. NUMS(2) .EQ. INDEX(N + 1))
337 +GOTO 37
338 C
339 C      WE KNOW THAT ONE OF THE NUMBERS READ IN IS THE NEW ELEMENT
340 C      BEING ADDED. FIND OUT WHICH ONE IT IS AND SEE IF OTHER
341 C      NUMBER TYPED IS IN SYSTEM INDEX
342 C
343      IF(NUMS(1) .EQ. INDEX(N + 1)) GOTO 39
344 C
345 C      IF WE MAKE IT HERE, SECOND NUMBER TYPED IS NEW ELEMENT.
346 C      SEE IF FIRST ONE IS IN SYSTEM INDEX
347 C
348      K = 1

```

```

349 42 CALL FINDIT(N,NUMS(K),J,I,J,INDEX,FOUND1,FOUND2,SYS)
350 IF(FOUND1) GOTO 40
351 C
352 C ERROR MESSAGE, ONE OF NUMBERS TYPED IS NOT IN SYSTEM INDEX
353 C
354 WRITE(TTYOUT,108) NUMS(K)
355 GOTO 37
356 C
357 C RESET BACK TO AUTOMATIC QUERIES
358 C
359 41 QST = .FALSE.
360 WRITE(TTYOUT,110)
361 GOTO 7
362 C
363 C ERROR MESSAGE, ONE OR THE OTHER OF THE NUMBERS TYPED MUST BE THE
364 C NEW ELEMENT BEING INTRODUCED INTO SYSTEM
365 C
366 38 WRITE(TTYOUT,109) INDEX(N + 1)
367 GOTO 37
368 C
369 C IF WE MAKE IT HERE, FIRST NUMBER IS NEW ELEMENT.
370 C SEE IF SECOND TYPED IS IN SYSTEM INDEX
371 C
372 39 K = 2
373 GOTO 42
374 C
375 C NOW CHECK TO MAKE SURE THAT USER HASN'T EITHER ALREADY
376 C ANSWERED THIS QUESTION OR INFERED THE ANSWER
377 C
378 40 ZPOINT = 1
379 IF(K .EQ. 2) ZPOINT = ZPOINT + N
380 IF(.NOT. FLAG(ZPOINT)) GOTO 43
381 C
382 K = N + 1
383 CALL FINDIT(K,NUMS(1),NUMS(2),I,J,INDEX,FOUND1,FOUND2,SYS)
384 IF(MAT(I,J)) GOTO 44
385 WRITE(TTYOUT,111) NUMS(1), NUMS(2)
386 GOTO 37
387 44 WRITE(TTYOUT,112) NUMS(1), NUMS(2)
388 GOTO 37
389 C
390 C FORMATS
391 C
392 100 FORMAT(42H-TYPE NEXT ELEMENT NUMBER OR 0 FOR BREAK ?)
393 101 FORMAT(12H-***ERROR***,15,24H ALREADY IN SYSTEM INDEX)
394 102 FORMAT(29H-***ERROR*** TEXT FOR ELEMENT,15,17H NOT IN TEXT FILE)
395 103 FORMAT(62H-***NOTE*** NUMBER OF ELEMENTS HAS REACHED COMPUTER'S LI
396 +HIT OF,15/24H "80" COMMAND TERMINATED)
397 104 FORMAT(46H-***ERROR*** INVALID RESPONSE TO LAST QUESTION)
398 105 FORMAT(42H-***NOTE*** BORDERING SEQUENCE FOR ELEMENT,15,17H HAS BE
399 +EN ABORTED)
400 106 FORMAT(1H+,2X,A2)
401 107 FORMAT(59H WHICH ELEMENTS TO BE COMPARED ? (TYPE 0 FOR AUTOMATIC))
402 108 FORMAT(12H-***ERROR***,15,20H NOT IN SYSTEM INDEX)
403 109 FORMAT(12H-***ERROR***,15,28H MUST BE ONE OF THE ELEMENTS)
404 110 FORMAT(43H 0 ACKNOWLEDGED, BEGINING AUTOMATIC QUERIES)
405 111 FORMAT(64H-SORRY, THIS QUESTION HAS BEEN EITHER ASKED ALREADY OR F
406 +ILLED BY/10H INFERENCE,15,2H R,15,5H = NO)

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407 112 FORMAT(64H-SORRY, THIS QUESTION HAS BEEN EITHER ASKED ALREADY OR F
408 FILLED BY/10H INFERENCE,15,2H R,15,6H = YES)
409 200 FORMAT(A2)
410 END

295


```

1 SUBROUTINE FINDZ(DIMPHI,PHI,ZPOINT,FLAG,SYS2)
2 C
3 C THIS SUBROUTINE RETURNS THE Z (ZPOINT) WITH THE MAXIMUM
4 C INFERENCE POTENTIAL GIVEN A PHI MATRIX
5 C
6 C WRITTEN BY: DAVID R. YINGLING, JR.
7 C ENGINEERING AND PUBLIC POLICY GROUP
8 C UNIVERSITY OF DAYTON
9 C DAYTON, OHIO 45469
10 C
11 C VARIABLE NAME DESCRIPTION
12 C
13 C DIMPHI THE CURRENT DIMENSION OF "PHI".
14 C
15 C PHI THE TRANSITIVE BORDERING INFERENCE
16 C OPPORTUNITY MATRIX.
17 C
18 C ZPOINT THE ROW/COL OF "PHI" WHICH
19 C HAS MAXIMUM INFERENCE POTENTIAL
20 C
21 C FLAG VECTOR DENOTING ROWS/COLS OF "PHI"
22 C WHICH HAVE BEEN SET TO ALL ZEROS
23 C
24 C SYS2 DIMENSION SIZES OF SYSTEM MATRICES
25 C
26 C -----
27 C
28 C MIN SCRATCH VECTOR CONTAINING THE
29 C MIN2(V1,V2) SET
30 C
31 C V1 VECTOR CONTAINING THE NUMBER OF ONES
32 C OF THE COLS OF "PHI".
33 C
34 C V1 VECTOR CONTAINING THE NUMBER OF ONES
35 C OF THE ROWS OF "PHI".
36 C
37 C Z DENOTES MEMBERS OF THE V1 SET
38 C
39 C
40 C *****NOTE*****
41 C THE DIMENSIONS OF "MIN", "V1", "V2", + "Z" SHOULD BE SET
42 C EQUAL TO "SYS2".
43 C
44 C
45 C IMPLICIT INTEGER*2 (A-Z)
46 C LOGICAL*1 PHI(SYS2,SYS2), Z(256), FLAG(SYS2)
47 C DIMENSION MIN(256), V1(256), V2(256)
48 C
49 C CALCULATE THE SET SET V
50 C
51 C DETERMINE V (1) AND V (1) FOR Z(1)
52 C 1 2
53 C WHERE: V (1) IS THE NUMBER OF 1'S IN COL 1 OF PHI
54 C 1
55 C V (1) IS THE NUMBER OF 1'S IN ROW 1 OF PHI
56 C 2
57 C
58 C DO 1 =1,DIMPHI

```

```

59      V1(I) = 0
60      V2(I) = 0
61      IF(FLAG(I)) GOTO 1
62      V1(I) = V1SET(PHI,I,DIMPHI,SYS2)
63      V2(I) = V2SET(PHI,I,DIMPHI,SYS2)
64      1 CONTINUE
65 C
66 C      NOW DETERMINE THE SET V'
67 C
68 C      V' = MAX2(MIN2(V1,V2))
69 C
70 C      FIND MIN2(V1,V2)
71 C
72      2 DO 3 I=1,DIMPHI
73          MIN(I) = 0
74          IF(FLAG(I)) GOTO 3
75          MIN(I) = MIN2(V1(I),V2(I))
76      3 CONTINUE
77 C
78 C      GET MAX2(MIN2(V1,V2))
79 C
80      4 BIGGER = MIN(1)
81      ZPOINT = 0
82      DO 5 I=1,DIMPHI
83          IF(FLAG(I)) GOTO 5
84          BIG = MAX2(BIGGER,MIN(I))
85          IF(BIG.GT. BIGGER) CALL ZER(Z,I,SYS2)
86          IF(BIG.LE. MIN(I)) Z(I) = .TRUE.
87          IF(BIG.GE. BIGGER) BIGGER = BIG
88      5 CONTINUE
89 C
90 C      NOW FIND SET OF V' FOR WHICH
91 C      V1 + V2 IS A MAXIMUM
92 C
93      9 BIGGER = 0
94      DO 10 I=1,DIMPHI
95          IF(.NOT. Z(I)) GOTO 10
96          BIG = MAX2(V1(I) + V2(I),BIGGER)
97          IF(BIG.LE. BIGGER) GOTO 10
98          BIGGER = BIG
99      ZPOINT = I
100      10 CONTINUE
101      11 RETURN
102      END

```

```

1 SUBROUTINE TBPHI(N,A,PHI,DIMPHI,SUBREL,SYS2,SYS)
2 C
3 C THIS SUBROUTINE FORMS THE TRANSITIVE BORDERING INFERENCE
4 C OPPORTUNITY MATRIX
5 C
6 C WRITTEN BY: DAVID R. YINGLING, JR.
7 C ENGINEERING AND PUBLIC POLICY GROUP
8 C UNIVERSITY OF DAYTON
9 C DAYTON, OHIO 45469
10 C
11 C VARIABLE NAME DESCRIPTION
12 C
13 C N THE NUMBER OF ELEMENTS IN "A".
14 C
15 C A THE CURRENT MODEL MATRIX
16 C
17 C PHI OUTPUT INFERENCE OPPORTUNITY MATRIX
18 C
19 C DIMPHI THE CURRENT DIMENSION OF "PHI".
20 C
21 C SUBREL SUBORDINATION RELATIONSHIP SWITCH.
22 C IF THIS VALUE IS .TRUE., A SPECIAL
23 C PROCEDURE IS PERFORMED ON THE "PHI"
24 C MATRIX. IF .FALSE., NO SPECIAL
25 C PROCEDURE IS DONE.
26 C
27 C SYS2 DIMENSION SIZES OF "PHI" AND
28 C ASSOCIATED MATRICES
29 C
30 C SYS DIMENSION SIZES OF SYSTEM MATRICES.
31 C
32 C
33 C IMPLICIT INTEGER*1 (A-Z)
34 C INTEGER N
35 C LOGICAL SUBREL
36 C LOGICAL*1 A(SYS,SYS), PHI(SYS2,SYS2)
37 C
38 C
39 C FORM THE INFERENCE OPPORTUNITY MATRIX PHI IN TWO STEPS
40 C
41 C 1. FORM N1
42 C
43 C A 0
44 C
45 C "N1" =
46 C
47 C B A
48 C
49 C -T
50 C WHERE: B = A IF A TRANSITIVE RELATIONSHIP IS USED.
51 C
52 C -T
53 C B = A + I IF A TRANSITIVE AND SUBORDINATION RELATIONSHIP
54 C IS USED.
55 C
56 C 2. MULTIPLY ON "N1" TO OBTAIN PHI
57 C
58 C A 0

```

```

50 C
60 C "PHI"
61 C
62 C          ABA      A
63 C
64 C      DO STEP 1
65 C
66 C      DO 1 I=1,N
67 C      DO 1 J=1,N
68 C
69 C      PUT "A" ON UPPER LEFT OF "N1"
70 C
71 C      PHI(I,J) = A(I,J)
72 C
73 C      PUT "A" ON LOWER RIGHT OF "N1"
74 C
75 C      PHI(I + N,J + N) = A(I,J)
76 C
77 C      PUT "B" ON LOWER LEFT OF "N1"
78 C
79 C      PHI(I + N,J) = .NOT. A(J,I)
80 C      1 CONTINUE
81 C
82 C      SEE IF USER WANTS SUBORDINATION RELATION.
83 C      IF YES, ADD "I" TO "B" SECTION.
84 C      IF NO, SKIP AROUND AND CONTINUE PROLESSING.
85 C
86 C      IF(.NOT. SUBREL) GOTO 2
87 C      DO 3 I=1,N
88 C      PHI(I + N,I) = .TRUE.
89 C      3 CONTINUE
90 C
91 C      DO STEP 2
92 C
93 C      FIRST MULTIPLY "B" TIMES "A" AND STORE IN "O" AS FIRST OPERATION
94 C
95 C      2 DIM = N + 1
96 C
97 C      DO 4 I=1,N
98 C      DO 4 J=1,N
99 C
100 C      PHI(I,J + N) = .FALSE.
101 C
102 C      DO 4 K=1,N
103 C      PHI(I,J + N) = PHI(I,J + N) .OR. (PHI(I+N,K) .AND. PHI(K+N,J+N))
104 C      4 CONTINUE
105 C
106 C      NOW MULTIPLY "A" (LEFT OF "O") TIMES THE PRODUCT STORED IN
107 C      "O" AS SECOND OPERATION. STORE THAT PRODUCT INTO ITS CORRECT
108 C      LOCATION (ABA).
109 C
110 C      DO 5 I=1,N
111 C      DO 5 J=1,N
112 C
113 C      PHI(I + N,J) = .FALSE.
114 C
115 C      DO 5 K=1,N
116 C      PHI(I + N,J) = PHI(I + N,J) .OR. (PHI(I,K) .AND. PHI(K,J+N))

```

```

117      5 CONTINUE
118 C
119 C      AS FINAL OPERATION, ZAP "O" SECTION BACK TO ZEROS
120 C
121      DO 6 I=1,N
122      DO 6 J=DIM,DIMPHI
123      PHI(I,J) = .FALSE.
124      6 CONTINUE
125      RETURN
126      END

```

300

```

1  INTEGER FUNCTION VISET*2(PHI,J,DIMPHI,SY52)
2  C
3  C  THIS FUNCTION SUBPROGRAM COMPUTES THE NUMBER OF
4  C  UNES IN COL "J" OF PHI.
5  C
6  C  WRITTEN BY:  DAVID R. YINGLING, JR
7  C               ENGINEERING AND PUBLIC POLICY GROUP
8  C               UNIVERSITY OF DAYTON
9  C               DAYTON, OHIO  45469
10 C
11  IMPLICIT INTEGER*2 (A-Z)
12  LOGICAL*1 PHI(SY52,SY52)
13  VISET = 0
14  DO 1 I=1,DIMPHI
15  IF(PHI(I,J)) VISET = VISET + 1
16  1 CONTINUE
17  RETURN
18  END

```

30i

```

1      INTEGER FUNCTION V2SET*2(PHI,I,DIMPHI,SYS2):
2 C
3 C      THIS FUNCTION SUBPROGRAM COMPUTER THE NUMBER OF
4 C      ONES IN ROW "I" OF PHI.
5 C
6 C      WRITTEN BY:  DAVID R. YINGLING, JR
7 C                  ENGINEERING AND PUBLIC POLICY GROUP
8 C                  UNIVERSITY OF DAYTON
9 C                  OHIO  45469
10 C
11     IMPLICIT IN
12     LOGICAL*1
13     V2SET = .
14     DO 1 J=1
15     IF(PHI(I,J)) V2SET = 1
16     1 CONTINUE
17     RETURN
18     END

```

```

1      SUBROUTINE DIGLEV(N,MATRIX,INDEX,TTYOUT,SY)
2 C
3 C      THIS SUBROUTINE WILL DISPLAY THE DIGRAPH OF "MATRIX" IN A
4 C      LEVELS FORMAT
5 C
6 C      WRITTEN BY: DAVID R. YINGLING, JR.,
7 C                  ENGINEERING AND PUBLIC POLICY GROUP
8 C                  UNIVERSITY OF DAYTON
9 C                  DAYTON, OHIO 45469
10 C
11 C      VARIABLE NAME          DESCRIPTION
12 C
13 C      N                      THE NUMBER OF ELEMENTS IN "MATRIX"
14 C
15 C      MATRIX                 REACHABILITY MATRIX TO BE DISPLAYED
16 C
17 C      INDEX                  THE INDEX SET OF "MATRIX"
18 C
19 C      TTYOUT                 FORTRAN WRITE UNIT NUMBER FOR
20 C                           TELETYPE
21 C
22 C      SYS                    DIMENSION SIZES OF SYSTEM MATRICES
23 C
24 C -----
25 C
26 C      NC                     THE NUMBER OF ELEMENTS IN THE
27 C                           CONDENSATION MATRIX
28 C
29 C      MATRXX                 SCRATCH MATRIX FOR LEVELS ROUTINE
30 C
31 C      INDX                  SCRATCH INDEX VECTOR FOR LEVELS
32 C                           ROUTINE
33 C
34 C      LEVELS                 SCRATCH VECTOR DENOTING THE NUMBER
35 C                           OF ELEMENTS ON EACH LEVEL,
36 C                           LEVELS(I) = NUMBER OF ELEMENTS ON
37 C                           LEVEL #I.
38 C
39 C      NLEVEL                 THE TOTAL NUMBER OF LEVELS
40 C
41 C
42 C      *****NOTE*****
43 C      THE DIMENSIONS OF "MATRXX", "INDXX", AND "LEVELS" SHOULD BE
44 C      EQUAL TO "SYS".
45 C
46 C
47 C      IMPLICIT INTEGER*2 (A-Z)
48 C      INTEGER N, INDEX, INDX, NC
49 C      DIMENSION INDEX(SYS), INDX(128), LEVELS(128)
50 C      LOGICAL*1 MATRIX(SYS,SYS), MATRXX(128,128)
51 C
52 C      CHECK FOR ERROR
53 C
54 C      IF(N .LE. 0) GOTO 1
55 C
56 C      STEP 1---REARRANGE "MATRIX" INTO HIERARCHIAL FORM - PUT
57 C      RESULT IN "MATRXX" AND "INDXX".
58 C

```



```

59      CALL HIERCH(N,MATRIX,INDEX,MATRXX,INXXX,NLEVEL,LEVELS,SYS)
60 C
61 C      STEP 2---PUT "MATRXX" INTO STANDARD FORM
62 C
63      CALL STAN(N,MATRXX,INXXX,NLEVEL,LEVELS,SYS)
64 C
65 C      STEP 3---COMPUTE CONDENSATION MATRIX OF "MATRXX" - LEAVE
66 C          RESULT IN "MATRXX".
67 C
68      NC = N
69      CALL CONDE(NC,MATRXX,INXXX,LEVELS,TTYOUT,.TRUE.,SYS)
70 C
71 C      COMPUTE NONREDUNDANT ADJACENCY MATRIX (SKELETON MATRIX)
72 C
73      CALL SKLTN(NC,MATRXX,SYS)
74 C
75 C      PRINT LEVELS FORMATTED DIGRAPH
76 C
77      CALL DISPLV(NC,MATRXX,INXXX,LEVELS,TTYOUT,SYS)
78      RETURN
79 C
80 C      ERROR MESSAGE
81 C
82      1 WRITE(TTYOUT,100)
83      RETURN
84 C
85 C      FORMAT
86 C
87      100 FORMAT(42H-***ERROR*** NO STRUCTURE CURRENTLY EXISTS)
88      END

```

```

1 SUBROUTINE DISPLV(N,SKLTN,INDEX,LEVELS,TTYOUT,SYS)
2 C
3 C THIS SUBROUTINE PRINTS A LEVEL FORMATTED DIGRAPH FROM
4 C THE INPUT SKELETON MATRIX.
5 C
6 C WRITTEN BY: DAVID R. YINGLING, JR.
7 C ENGINEERING AND PUBLIC POLICY GROUP
8 C UNIVERSITY OF DAYTON
9 C DAYTON, OHIO 45469
10 C
11 C VARIABLE NAME DESCRIPTION
12 C
13 C N NUMBER OF ELEMENTS IN THE INPUT
14 C SKELETON MATRIX
15 C
16 C SKLTN INPUT SKELETON MATRIX
17 C
18 C INDEX INDEX SET OF "SKLTN".
19 C
20 C LEVELS INPUT VECTOR DENOTING THE NUMBER
21 C OF ELEMENTS ON EACH LEVEL,
22 C LEVELS(I) = NUMBER OF ELEMENTS ON
23 C LEVEL # I.
24 C
25 C TTYOUT FORTRAN WRITE UNIT NUMBER FOR
26 C TELETYPE.
27 C
28 C SYS DIMENSION SIZES OF SYSTEM MATRICES
29 C
30 C -----
31 C
32 C LIST SCRATCH VECTOR FOR LEVELS PRINTOUT.
33 C "LIST" CONTAINS THE ELEMENT NUMBERS
34 C FOR PRINTING ON THE TELETYPE.
35 C
36 C LEVEL THE CURRENT LEVEL NUMBER
37 C
38 C ROW THE ROW OF THE LAST ELEMENT THAT
39 C IS ON LEVEL # LEVEL.
40 C
41 C COUNT THE NUMBER OF ELEMENTS IN "LIST".
42 C
43 C
44 C *****NOTE*****
45 C THE DIMENSION OF "LIST" SHOULD BE EQUAL TO "SYS".
46 C
47 C
48 C IMPLICIT INTEGER*2 (A-Z)
49 C INTEGER N, INDEX
50 C DIMENSION INDEX(SYS), LEVELS(SYS), LIST(120)
51 C LOGICAL*1 SKLTN(SYS,SYS)
52 C
53 C INITIALIZE PROCEDURE
54 C
55 C I = 1
56 C LEVEL = 0
57 C ROW = 1
58 C

```

```

59 C      PRUCESS NEXT LEVEL
60 C
61 C      1 LEVEL = LEVEL + 1
62 C      ROW = ROW + LEVELS(LEVEL)
63 C      WRITE(TTYOUT,100) LEVEL
64 C
65 C      IF THIS IS THE FIRST LEVEL, DO SPECIAL PROCESSING
66 C
67 C      2 IF(LEVEL .EQ. 1) GOTO 3
68 C
69 C      PROCESS ELEMENT #1,.....
70 C      FIND ALL ELEMENTS THAT ELEMENT #1 REACHES TO
71 C
72 C      COUNT = 0
73 C      IMINUS= 1 - 1
74 C      DO 4 J=1,IMINUS
75 C      IF(.NOT. SKLTN(1,J)) GOTO 4
76 C
77 C      FOUND ONE, PUT INTO "LIST" FOR PRINTOUT
78 C
79 C      COUNT = COUNT + 1
80 C      LIST(COUNT) = INDEX(J)
81 C      4 CONTINUE
82 C
83 C      ALL DONE PROCESSING ELEMENT #1, PRINT OUT LINE
84 C
85 C      WRITE(TTYOUT,101) INDEX(1), (LIST(II),II=1,COUNT)
86 C
87 C      POINT TO NEXT ELEMENT
88 C
89 C      5 I = I + 1
90 C
91 C      ARE WE DONE PRINTING THE DIGRAPH ???
92 C
93 C      IF(I .GT. N) RETURN
94 C
95 C      ARE WE DONE WITH THIS LEVEL ON DIGRAPH ???
96 C
97 C      IF(I .EQ. ROW) GOTO 1
98 C      GOTO 2
99 C
100 C     SPECIAL PROCESSING FOR FIRST LEVEL
101 C
102 C     3 WRITE(TTYOUT,102) INDEX(1)
103 C     GOTO 5
104 C
105 C     FORMATS
106 C
107 C     100 FORMAT(1H-,10X,11HLEVEL NO. ,13/1H0)
108 C     101 FORMAT(11X,15,3H =>,18(7(13,1H,1/20X))
109 C     102 FORMAT(14X,15)
110 C     END

```

```

1 SUBROUTINE HIERCHIN(INREA,INDXIN,REAH,INDXH,NL,LEVELS,SYS)
2 C
3 C THIS SUBROUTINE REARRANGES A REACHABILITY MATRIX INTO A
4 C LEVEL ORIENTED HIERARCHIAL REACHABILITY MATRIX.
5 C
6 C
7 C VARIABLE NAME          DESCRIPTION
8 C
9 C INDXIN                VECTOR CONTAINING THE INDEX SET OF
10 C                     THE INPUT REACHABILITY MATRIX,
11 C
12 C INDXH                 VECTOR CONTAINING THE INDEX SET OF THE
13 C                     OUTPUT HIERARCHIAL REACHABILITY MATRIX
14 C
15 C FLAG                  LOGICAL VECTOR WHICH DENOTES ELEMENTS
16 C                     THAT HAVE ALREADY BEEN PROCESSED. IF
17 C                     FLAG(I) = .TRUE., ELEMENT #I HAS BEEN
18 C                     PROCESSED.
19 C
20 C LEVELS                VECTOR CONTAINING THE NUMBER OF
21 C                     ELEMENTS ON EACH LEVEL. LEVELS(I) =
22 C                     NUMBER OF ELEMENTS ON LEVEL #I.
23 C
24 C TEMP                  SCRATCH VECTOR USED BY LEVELS
25 C                     PARTITION ALGORITHM.. IT HOLDS THE
26 C                     ELEMENTS THAT ARE ON THE CURRENT LEVEL
27 C
28 C INREA                 INPUT (ARGUMENT) REACHABILITY MATRIX
29 C
30 C REAH                  OUTPUT (RESULTANT) HIERARCHIAL
31 C                     REACHABILITY MATRIX
32 C
33 C N                     NUMBER OF ELEMENTS IN BOTH INPUT
34 C                     REACHABILITY MATRIX AND OUTPUT
35 C                     HIERARCHIAL REACHABILITY MATRIX
36 C
37 C NEAP                  NUMBER OF ELEMENTS ALREADY PROCESSED
38 C
39 C NEL                   NUMBER OF ELEMENTS ON CURRENT LEVEL
40 C
41 C NL                    NUMBER OF CURRENT LEVEL
42 C
43 C SYS                   DIMENSION SIZES OF SYSTEM MATRICES
44 C
45 C *****NOTE*****
46 C     THE DIMENSIONS OF "TEMP" AND "FLAG" SHOULD BE EQUAL TO
47 C     "SYS".
48 C
49 C IMPLICIT INTEGER*2 (A-Z)
50 C INTEGER N, INDXIN, INDXH
51 C DIMENSION INDXIN(SYS), INDXH(SYS), LEVELS(SYS), TEMP(128)
52 C LOGICAL*1 INREA(SYS,SYS), REAH(SYS,SYS), FLAG(128)
53 C
54 C COPY INREA INTO REAH : INITIALIZE FLAG
55 C
56 C DO 1 I=1,N
57 C   FLAG(I) = .FALSE.
58 C DO 1 J=1,N

```

```

59      REAH(I,J) = .INREA(I,J)
60      1 CONTINUE
61 C
62 C      INITIALIZE LEVELS PARTITION ALGORITHM
63 C
64      NL = 0
65      NEAP = 0
66 C
67 C      BEGIN LEVELS PARTITION ALGORITHM
68 C
69 C      THIS ALGORITHM REARRANGES INDXIN ACCORDING TO A LEVELS
70 C      PARTITION. THE RESULT IS IN INDXH.
71 C
72      2 NL = NL + 1
73      NEL = 0
74 C
75 C      FIND AN ELEMENT TO PROCESS
76 C
77      DO 3 I=1,N
78      IF(FLAG(I)) GOTO 3
79 C
80 C      TEST TO SEE IF THE REACHABILITY SET (R) IS A
81 C      SUBSET OF THE ANTECEDENT SET (A) FOR THIS ELEMENT
82 C
83      DO 4 J=1,N
84      IF(REAH(I,J) .AND. .NOT. REAH(J,I)) GOTO 3
85      4 CONTINUE
86 C
87 C      COME HERE IF R WAS A SUBSET OF A
88 C
89      NEL = NEL + 1
90      INDXH(NEAP + NEL) = INDXIN(I)
91      TEMP(NEL) = I
92      3 CONTINUE
93      NEAP = NEAP + NEL
94      LEVELS(NL) = NEL
95 C
96 C      FOUND ALL ELEMENTS ON CURRENT LEVEL (NL).
97 C      BLANK ROW AND COL ON REAH FOR ALL ELEMENTS ON THIS LEVEL.
98 C
99      DO 5 I=1,NEL
100     TEMP1 = TEMP(I)
101     FLAG(TEMP1) = .TRUE.
102     DO 5 J=1,N
103     REAH(TEMP1,J) = .FALSE.
104     REAH(J,TEMP1) = .FALSE.
105     5 CONTINUE
106 C
107 C      CHECK TO SEE IF ALL ELEMENTS HAVE BEEN PROCESSED
108 C
109     IF(NEAP .LT. N) GOTO 2
110 C
111 C      COME HERE WHEN ALL ELEMENTS HAVE BEEN PROCESSED.
112 C      LEVELS PARTITION ALGORITHM IS NOW COMPLETE, BEGIN TO
113 C      CONSTRUCT A LEVELS PARTITIONED HIERARCHIAL REACHABILITY
114 C      MATRIX BASED ON INDXH.
115 C
116 C      EXCHANGE ALL ROWS FIRST ACCORDING TO INDXH

```

```

117 C
118 DO 6 I=1,N
119 DO 7 J=1,N
120 IF (INDXH(I) .EQ. INDXIN(J)) GOTO 8
121 7 CONTINUE
122 C
123 8 DO 6 K=1,N
124 REAH(I,K) = INREA(J,K)
125 6 CONTINUE
126 C
127 C COPY REAH INTO INREA. THIS IS A NECESSARY STEP, DO NOT
128 C TAKE OUT !!!!
129 C
130 DO 9 I=1,N
131 DO 9 J=1,N
132 INREA(I,J) = REAH(I,J)
133 9 CONTINUE
134 C
135 C WE PRESENTLY HAVE A MATRIX (INREA) WHICH IS INDEXED ON THE
136 C TOP BY INDXH AND DOWN THE SIDE BY INDXIN. REARRANGE THE
137 C COLS SO THEY ALSO ARE INDEXED BY INDXH. LEAVE RESULT IN
138 C REAH.
139 C
140 DO 10 I=1,N
141 DO 11 J=1,N
142 IF (INDXH(I) .EQ. INDXIN(J)) GOTO 13
143 11 CONTINUE
144 C
145 13 DO 10 K=1,N
146 REAH(K,I) = INREA(K,J)
147 10 CONTINUE
148 C
149 C IN ITS PRESENT FORM, INREA, FOR ALL PURPOSES OTHER THAN
150 C THIS SUBROUTINE, IS SCRAMBLED. COPY REAH INTO INREA TO
151 C SOLVE THIS PROBLEM.
152 C
153 DO 14 I=1,N
154 INDXIN(I) = INDXH(I)
155 DO 14 J=1,N
156 INREA(I,J) = REAH(I,J)
157 14 CONTINUE
158 C
159 RETURN
160 ENO

```

```

1  SUBROUTINE STAN(N,MATRIX,INDEX,NLEVEL,LEVELS,SYS)
2  C
3  C  THIS SUBROUTINE CONVERTS AN INPUT HIERARCHIAL REACHABILITY
4  C  MATRIX (MATRIX) INTO ITS STANDARD FORM.
5  C
6  C  EDITED BY: DAVID R. YINGLING, JR.
7  C             ENGINEERING AND PUBLIC POLICY GROUP
8  C             UNIVERSITY OF DAYTON
9  C             DAYTON, OHIO 45469
10 C
11 C
12 C  VARIABLE NAME          DESCRIPTION
13 C
14 C  N                      NUMBER OF ELEMENTS IN INPUT MATRIX
15 C
16 C  MATRIX                INPUT/OUTPUT HIERARCHIAL REACHABILITY
17 C                        MATRIX.
18 C
19 C  INDEX                 INPUT/OUTPUT INDEX SET OF "MATRIX".
20 C
21 C  LEVELS                INPUT VECTOR DENOTING THE NUMBER OF
22 C                        ELEMENTS ON EACH LEVEL.  LEVELS(I) =
23 C                        THE NUMBER OF ELEMENTS ON LEVEL #I
24 C
25 C  NLEVEL                THE TOTAL NUMBER OF LEVELS
26 C
27 C  SYS                   DIMENSION SIZES OF SYSTEM MATRICES
28 C
29 C -----
30 C
31 C  NONES                 NUMBER OF ONES COUNTED
32 C
33 C  END                   ENDING SUBSCRIPT FOR LEVEL #I
34 C
35 C  START                 STARTING SUBSCRIPT FOR LEVEL #I
36 C
37 C
38 C  IMPLICIT INTEGER*2 (A-Z)
39 C  INTEGER N, INDEX
40 C  DIMENSION INDEX(SYS), LEVELS(SYS)
41 C  LOGICAL*1 MATRIX(SYS,SYS), SWIT
42 C
43 C  CHECK FIRST TO MAKE SURE ALL NON-CYCLE ELEMENTS
44 C  ARE UPPERMOST ON EACH LEVEL
45 C
46 C  END = 0
47 C  DO 1 I=1,NLEVEL
48 C  START = END + 1
49 C  END   = END + LEVELS(I)
50 C
51 C  IF THE NUMBER OF ELEMENTS ON LEVEL #I IS TWO OR
52 C  LESS, NO RE-ADJUSTMENT IS NECESSARY
53 C
54 C  IF(LEVELS(I) .LE. 2) GOTO 1
55 C
56 C  FIND AND MOVE NON-CYCLE ELEMENTS UP ON MATRIX IF NECESSARY
57 C
58 C  DO 1 II=START,END

```

310

```

59      LAST = J
60      DO 1 ROW=START,END
61      NONES = 0
62 C
63      DO 2 COL=START,END
64      IF(MATRIX(ROW,COL)) NONES = NONES + 1
65      2 CONTINUE
66 C
67 C      CHECK TO SEE IF ELEMENT NROW HAS LESS ONES THAN
68 C      LAST ELEMENT CHECKED
69 C
70      IF(NONES .LT. LAST) CALL SWITCH(N,MATRIX,INDEX,ROW,SYS)
71      IF(NONES .GE. LAST) LAST = NONES
72      1 CONTINUE
73 C
74 C      ALL NON-CYCLE ELEMENTS ARE AT THE BEGINING OF LEVELS
75 C      PARTITION. NOW GROUP THE CYCLES TOGETHER.
76 C
77      NMINUS = N - 1
78      3 SWIT = .FALSE.
79 C
80 C      CHECK FOR ONES ABOVE MAIN DIAGONAL
81 C
82      DO 4 I=1,NMINUS
83      IPLUS = I + 1
84      DO 5 J=IPLUS,N
85      IF(.NOT. MATRIX(I,J)) GOTO 5
86 C
87 C      COME HERE IF A ONE ABOVE THE MAIN DIAGONAL IS FOUND
88 C
89 C      CHECK TO SEE IF IT IS NEXT TO DIAGONAL ONE--IF SO,
90 C      DON'T SWITCH BECAUSE OF THE WAY THE "SWITCH" SUBROUTINE
91 C      WORKS. IF NOT, SWITCH THAT ELEMENT (J) WITH ELEMENT J-1.
92 C
93      6 JMINUS = J - 1
94      IF(I .EQ. JMINUS) GOTO 4
95 C
96      CALL SWITCH(N,MATRIX,INDEX,J,SYS)
97      SWIT = .TRUE.
98      J = J - 1
99      GOTO 6
100     5 CONTINUE
101     4 CONTINUE
102 C
103 C      IF ANY SWITCHING WAS DONE, WE NEED TO CHECK AGAIN
104 C      OTHERWISE RETURN.
105 C
106      IF(SWIT) GOTO 3
107      RETURN
108      END

```



```

1      SUBROUTINE SWITCH(N,MATRIX,INDEX,ROW,SYS)
2      C
3      C      THIS SUBROUTINE WILL SWITCH THE ROW AND COL OF "ROW" WITH
4      C      THE ROW AND COL OF "ROW" - 1.
5      C
6      C      WRITTEN BY: DAVID R. YINGLING, JR.
7      C      ENGINEERING AND PUBLIC POLICY GROUP
8      C      UNIVERSITY OF DAYTON
9      C      DAYTON, OHIO 45469
10     C
11     C
12     C      VARIABLE NAME      DESCRIPTION
13     C
14     C      N                  THE TOTAL NUMBER OF ELEMENTS IN
15     C                        "MATRIX".
16     C
17     C      MATRIX             INPUT MATRIX TO BE SWITCHED
18     C
19     C      INDEX              INDEX SET OF THE INPUT MATRIX
20     C
21     C      ROW                THE SUBSCRIPT OF THE MATRIX TO BE
22     C                        SWITCHED.
23     C
24     C      SYS                DIMENSION SIZE OF SYSTEM MATRICES
25     C
26     C      -----
27     C
28     C      OTHER              THE OTHER ROW TO BE SWITCHED. ALWAYS
29     C                        EQUAL TO "ROW" - 1
30     C
31     C      IMPLICIT INTEGER*2 (A-Z)
32     C      INTEGER N, INDEX, ITEMP
33     C      DIMENSION INDEX(SYS)
34     C      LOGICAL*1 MATRIX(SYS,SYS), TEMP
35     C
36     C      OTHER = ROW - 1
37     C
38     C      SWITCH THE ROWS
39     C
40     C      DO 1 I=1,N
41     C          TEMP          = MATRIX(ROW,I)
42     C          MATRIX(ROW,I)  = MATRIX(OTHER,I)
43     C          MATRIX(OTHER,I) = TEMP
44     C      1 CONTINUE
45     C
46     C      SWITCH THE COLS
47     C
48     C      DO 2 I=1,N
49     C          TEMP          = MATRIX(I,ROW)
50     C          MATRIX(I,ROW)  = MATRIX(I,OTHER)
51     C          MATRIX(I,OTHER) = TEMP
52     C      2 CONTINUE
53     C
54     C      SWITCH THE INDEX SET
55     C
56     C      ITEMP          = INDEX(ROW)
57     C      INDEX(ROW)     = INDEX(OTHER)
58     C      INDEX(OTHER)  = ITEMP

```

53
60
61

RETURN
END

313

```

1 SUBROUTINE CUNDE(N,MATRIX,INDEX,LEVELS,TTYOUT,TYPE,SYS)
2 C
3 C THIS SUBROUTINE TAKES THE INPUT HIERARCHIAL REACHABILITY MATRIX
4 C IN STANDARD FORM AND REDUCES EACH MAXIMAL CYCLE SET INTO A
5 C SINGLE PROXY ELEMENT -- THEREBY FORMING THE CONDENSATION MATRIX.
6 C
7 C EDITED BY: DAVID R. YINGLING, JR.
8 C ENGINEERING AND PUBLIC POLICY GROUP
9 C UNIVERSITY OF DAYTON
10 C DAYTON, OHIO 45469
11 C
12 C VARIABLE NAME DESCRIPTION
13 C
14 C N THE NUMBER OF ELEMENTS IN THE INPUT
15 C MATRIX. UPON COMPLETION OF THIS
16 C SUBROUTINE, THE NEW VALUE OF N WILL
17 C REFLECT THE ELEMENTS DELETED.
18 C
19 C MATRIX THE INPUT STANDARD FORM MATRIX/OUTPUT
20 C CONDENSATION MATRIX
21 C
22 C INDEX THE INDEX SET OF "MATRIX".
23 C
24 C LEVELS INPUT VECTOR WHICH CONTAINS THE
25 C NUMBER OF ELEMENTS IN EACH LEVEL.
26 C LEVELS(I) = NUMBER OF ELEMENTS ON
27 C LEVEL #I.
28 C
29 C TTYOUT FORTRAN WRITE UNIT NUMBER FOR
30 C THE TELETYPE
31 C
32 C TYPE LOGICAL VARIABLE WHICH WHEN .FALSE.
33 C SUPPRESSES THE PRINTING OF CYCLES AT
34 C THE TELETYPE. WHEN.TRUE., CYCLES
35 C ARE PRINTED
36 C
37 C SYS DIMENSION OF SYSTEM MATRICES
38 C
39 C -----
40 C
41 C LIST A SCRATCH VECTOR OF DIMENSION "SYS"
42 C THAT HOLDS THE INDEX NUMBERS OF THE
43 C CYCLE SET BEING OPERATED ON.
44 C
45 C COUNT THE NUMBER OF ELEMENTS IN THE CURRENT
46 C CYCLE SET BEING OPERATED ON.
47 C
48 C POSITN THE POSITION TALLY IN THE "LEVELS"
49 C VECTOR.
50 C
51 C *****NOTE*****
52 C THE DIMENSION OF "LIST" SHOULD BE EQUAL TO "SYS".
53 C
54 C IMPLICIT INTEGER*2 (A-Z)
55 C INTEGER N, INDEX
56 C DIMENSION INDEX(SYS), LEVELS(SYS), LIST(128)
57 C LOGICAL TYPE
58 C LOGICAL*1 MATRIX(SYS,SYS)

```

```

59 C
60 C CHECK FOR A ONE ABOVE THE MAIN DIAGONAL (FIND A CYCLE)
61 C
62 C
63 C I = 1
64 C 1 COUNT = 1
65 C LIST(1) = INDEX(I)
66 C J = I + 1
67 C IF(.NOT. MATRIX(I,J)) GOTO 5
68 C
69 C A ONE WAS FOUND, PUT THAT ELEMENT INTO CYCLE PRINTOUT LIST
70 C AND THEN ELIMINATE FROM THE MATRIX.
71 C
72 C 2 COUNT = COUNT + 1
73 C LIST(COUNT) = INDEX(J)
74 C CALL ELIM(N,MATRIX,INDEX,J,SYS)
75 C
76 C NOW REDUCE NUMBER OF ELEMENTS ON LEVEL WHERE AN ELEMENT
77 C WAS JUST ELIMINATED.
78 C
79 C POSITN = 0
80 C DO 3 II=1,N
81 C POSITN = POSITN + LEVELS(II)
82 C IF(POSITN .GE. J) GOTO 4
83 C 3 CONTINUE
84 C
85 C 4 LEVELS(II) = LEVELS(II) - 1
86 C
87 C ANY MORE ELEMENTS IN THIS CYCLE SET???
88 C
89 C IF(MATRIX(I,J) .AND. J .LE. N) GOTO 2
90 C
91 C WRITE CYCLE OUT TO TELETYPE
92 C
93 C IF (.TYPE) WRITE(TTYOUT,100) (LIST(III), III=1,COUNT)
94 C 5 I = I + 1
95 C IF(I .LT. N) GOTO 1
96 C RETURN
97 C
98 C FORMAT
99 C
100 C 100 FORMAT(11H0 CYCLE ON,2X,13(10(14,1H,)/13X))
101 C END

```

```

1      SUBROUTINE ELIMIN,MATRIX,INDEX,DELETE,SYS)
2 C
3 C      THIS SUBROUTINE ELIMINATES AN ELEMENT FROM A GIVEN
4 C      INPUT MATRIX.
5 C
6 C      EDITED BY: DAVID R. YINGLING, JR.
7 C                  ENGINEERING AND PUBLIC POLICY GROUP
8 C                  UNIVERSITY OF DAYTON
9 C                  DAYTON, OHIO 45469
10 C
11 C      VARIABLE NAME      DESCRIPTION
12 C
13 C      N                    NUMBER OF ELEMENTS IN "MATRIX", UPON
14 C                        COMPLETION OF THIS ROUTINE, THE NEW
15 C                        VALUE OF "N" WILL REFLECT THE ELEMENT
16 C                        DELETED.
17 C
18 C      MATRIX              THE INPUT/OUTPUT MATRIX
19 C
20 C      INDEX              THE INDEX SET OF "MATRIX".
21 C
22 C      DELETE              THE SUBSCRIPT OF "MATRIX" TO BE
23 C                        ELIMINATED.
24 C
25 C      SYS                 DIMENSION SIZE OF SYSTEM MATRICES
26 C
27      IMPLICIT INTEGER*2 (A-Z)
28      INTEGER N, INDEX
29      DIMENSION INDEX(SYS)
30      LOGICAL*1 MATRIX(SYS,SYS)
31 C
32 C      CHECK FOR "DELETE" .EQ. TO LAST LOGICAL POSITION ON "MATRIX"
33 C
34      NMINUS = N - 1
35      IF(DELETE .EQ. N) GOTO 1
36 C
37 C      MOVE ALL COLUMNS BELOW "DELETE" OVER BY 1
38 C
39      DO 2 ROW1=DELETE,NMINUS
40      ROW2 = ROW1 + 1
41 C
42      DO 3 COL=1,N
43      MATRIX(ROW1,COL) = MATRIX(ROW2,COL)
44      3 CONTINUE
45 C
46 C      MOVE ALL ROWS BELOW "DELETE" UP BY 1
47 C
48      DO 2 ROW=1,N
49      MATRIX(ROW,ROW1) = MATRIX(ROW,ROW2)
50      2 CONTINUE
51 C
52 C      FIX UP INDEX SET
53 C
54      DO 4 ROW1=DELETE,NMINUS
55      ROW2 = ROW1 + 1
56      INDEX(ROW1) = INDEX(ROW2)
57      4 CONTINUE
58 C

```

59 C SUBTRACT ONE FROM "N" TO REFLECT DELETION
60 C
61 1 4 * NMENUS
62 RETURN
63 END

317

```

1 SUBROUTINE SKLTN(N,MATRIX,SYS)
2 C
3 C THIS SUBROUTINE CONVERTS THE INPUT MATRIX INTO A
4 C NONREDUNDANT ADJACENCY MATRIX (SKELETON MATRIX)
5 C
6 C THIS ALGORITHM IS SIMILAR TO THE ONE DESCRIBED BY
7 C R.K. SHYAMASUNDAR, "BOOLEAN MATRIX METHOD FOR THE
8 C CONSTRUCTION OF HIERARCHIAL GRAPHS", IEEE TRANSACTIONS
9 C ON SYSTEMS, MAN, AND CYBERNETICS, VOL.SMC-8, NO. 2,
10 C FEBRUARY, 1978.
11 C
12 C EDITED BY: DAVID R. YINGLING, JR.
13 C ENGINEERING AND PUBLIC POLICY GROUP
14 C UNIVERSITY OF DAYTON
15 C DAYTON, OHIO 45469
16 C
17 C VARIABLE NAME DESCRIPTION
18 C
19 C N NUMBER OF ELEMENTS IN "MATRIX".
20 C
21 C MATRIX INPUT/OUTPUT MATRIX TO BE CONVERTED
22 C
23 C SYS DIMENSION SIZES OF SYSTEM MATRICES
24 C
25 C IMPLICIT INTEGER*2 (A-Z)
26 C INTEGER N
27 C LOGICAL*1 MATRIX(SYS,SYS)
28 C
29 C NMINUS = N - 1
30 C DO 1 I=2,NMINUS
31 C IMINUS = I - 1
32 C DO 1 J=1,IMINUS
33 C
34 C CHECK REACHABILITY OF NODE J TO NODE I
35 C
36 C IF(.NOT. MATRIX(I,J)) GOTO 1
37 C
38 C ADD ALL NODES TO ROW J THAT CAN BE REACHED FROM NODE I
39 C
40 C IPLUS = I + 1
41 C DO 1 K=IPLUS,N
42 C MATRIX(K,J) = MATRIX(K,J) .AND. .NOT. MATRIX(K,I)
43 C 1 CONTINUE
44 C RETURN
45 C END

```

```

1 SUBROUTINE ZEROVECTOR(I, SYS2)
2
3 C THIS SUBROUTINE ZEROS OUT ALL PREVIOUSLY
4 C FLAGED MAXIMUMS
5 C
6 C WRITTEN BY: DAVID R. YINGLING, JR
7 C             ENGINEERING AND PUBLIC POLICY GROUP
8 C             UNIVERSITY OF DAYTON
9 C             DAYTON, OHIO 45469
10 C
11 C IMPLICIT INTEGER*2 (A-Z)
12 C LOGICAL*1 VECTOR(SYS2)
13 C
14 C II = 1 - 1
15 C DO 1 J=1, II
16 C VECTOR(J) = .FALSE.
17 C 1 CONTINUE
18 C RETURN
19 C END

```

✓


```

1      SUBROUTINE QUEST(EL1,EL2,TTYOUT,QTYPE,XTWDS)
2 C
3 C      THIS SUBROUTINE DISPLAYS EITHER FULL TEXT
4 C      OR SYMBOLIC QUERIES
5 C
6 C      WRITTEN BY:  DAVID R. YINGLING, JR
7 C                  ENGINEERING AND PUBLIC POLICY GROUP
8 C                  UNIVERSITY OF DAYTON
9 C                  DAYTON, OHIO  45469
10 C
11     IMPLICIT INTEGER (A-Z)
12     INTEGER*2 EL1,EL2,TTYOUT,XTWDS
13     LOGICAL QTYPE
14     DIMENSION BLOCK(800),BUFFER(256)
15     EQUIVALENCE (BLOCK,R1)
16     COMMON /FTEXT/  R1(160), L1(160), R2(160), L2(160), R3(160)
17     COMMON UNUSED
18     DATA CRLF/215152551/, INIT/21525250C/
19 C
20 C      *****NOTE*****
21 C      THE DIMENSION OF "R1", "L1", "R2", "L2", "R3" SHOULD BE
22 C      EQUAL TO "XTWDS".
23 C
24 C      THE DIMENSION OF "BLOCK" SHOULD BE EQUAL TO "XTWDS * 5.0".
25 C
26 C
27 C      SYMBOLIC QUERIES ?
28 C
29 C      IF(QTYPE)  GOTO 1
30 C
31 C      NOPE, FULL TEXT - READ IN ELEMENTS
32 C
33     I1 = EL1 + 4
34     FIND(8'I1)
35     I2 = EL2 + 4
36     READ(8'I1) (L1(I),I=1,XTWDS)
37     READ(8'I2) (L2(I),I=1,XTWDS)
38     OFFSET = 0
39 C
40 C      PRESENT FIVE LINES OF I/O
41 C      1) INTRODUCTORY CLAUSE   (RELATIONAL CLAUSE 1)
42 C      2) ELEMENT A
43 C      3) CORRELATION CLAUSE   (RELATIONAL CLAUSE 2)
44 C      4) ELEMENT B
45 C      5) QUALIFYING CLAUSE    (RELATIONAL CLAUSE 3)
46 C
47     COUNT = 0
48     I4 = 2
49     BUFFER(14) = INIT
50     DO 2 I = 1,5
51     I1 = 0
52     I2 = 0
53 C
54 C      PRINT UP TO TEN LINES FOR EACH OF THE ABOVE PHRASES
55 C
56     DO 3 J = 1,10
57 C
58 C      IF LENGTH INDICATOR IS ZERO, DON'T PRINT

```

```

57 C
58 IF(BLOCK(J+OFFSET) .EQ. 0) GOTO 3
59 C
60 NOT ZERO, COMPUTE LENGTH AND LOCATION OF LINE
61 C
62 LENGTH = BLOCK(J+OFFSET)
63 C
64 I1 = I2 + 1
65 I2 = I1 + LENGTH - 1
66 C
67 PUT TEXT INTO BUFFER
68 C
69 C
70 I3 = I4 + 1
71 I4 = I3 + LENGTH - 1
72 C = I1
73 C
74 DO 4 I11=I3,I4
75 BUFFER(I11) = BLOCK(C + OFFSET + 10)
76 C = C + 1
77 4 CONTINUE
78 C
79 I4 = I4 + 1
80 BUFFER(I4) = CRLF
81 COUNT = COUNT + (LENGTH * 4) + 4
82 3 CONTINUE
83 OFFSET = OFFSET + TXTWDS
84 2 CONTINUE
85 COUNT = COUNT - 4
86 CALL ZAP(COUNT,BUFFER)
87 GOTO 6
88 C
89 SYMBOLIC TEXT QUERIES
90 C
91 1 WRITE(TTYOUT,101) EL1,EL2
92 6 RETURN
93 C
94 CHECK 4 FORMAT
95 C
96 100 FORMAT(1X,17A4)
97 101 FORMAT(1H-,15,2H R,15,2H? )
98 END

```

32i

```

1 SUBROUTINE FINDIT(N,N1,N2,S1,S2,INDEX,FOUND1,FOUND2,SYS)
2 C
3 C THIS SUBROUTINE FINDS ELEMENTS N1, N2 IN THE INDEX SET
4 C
5 C THE VALUES S1, S2 ARE THE POSITIONS OF N1 AND N2 IN THE
6 C INDEX SET.
7 C
8 C FOUND1 AND FOUND2 ARE LOGICAL VALUES AND ARE SET EQUAL
9 C TO .TRUE. IF N1 OR N2 (RESPECTIVELY) ARE FOUND IN THE
10 C INDEX SET.
11 C
12 C WRITTEN BY: DAVID R. YINGLING, JR.
13 C ENGINEERING AND PUBLIC POLICY GROUP
14 C UNIVERSITY OF DAYTON
15 C DAYTON, OHIO 45469
16 C
17 C IMPLICIT INTEGER*2 (A-Z)
18 C INTEGER N, INDEX
19 C DIMENSION INDEX(SYS)
20 C LOGICAL FOUND1, FOUND2
21 C
22 C FOUND1 = .FALSE.
23 C FOUND2 = .FALSE.
24 C S1 = 0
25 C S2 = 0
26 C
27 C DO 1 I=1,N
28 C IF(N1 .EQ. INDEX(I)) S1 = I
29 C IF(N2 .EQ. INDEX(I)) S2 = I
30 C 1 CONTINUE
31 C
32 C IF(S1 .GT. 0) FOUND1 = .TRUE.
33 C IF(S2 .GT. 0) FOUND2 = .TRUE.
34 C RETURN
35 C END

```

```

1 SUBROUTINE GETNUM(ARRAY,N,TTYIN,TTYOUT)
2 C
3 C THIS SUBROUTINE WILL READ "N" UNSIGNED INTEGERS FROM
4 C THE TERMINAL TYPED IN A FREE FORMAT AND STORE THEM IN
5 C "ARRAY"
6 C
7 C WRITTEN BY: DAVID R. YINGLING, JR.
8 C ENGINEERING AND PUBLIC POLICY GROUP
9 C UNIVERSITY OF DAYTON
10 C DAYTON, OHIO 45469
11 C
12 C IMPLICIT INTEGER*2 (A-Z)
13 C INTEGER NUMS,BUFFER,BLANK,COMMA
14 C DIMENSION ARRAY(1), BUFFER(80), NUMS(10)
15 C DATA NUMS/1H0,1H1,1H2,1H3,1H4,1H5,1H6,1H7,1H8,1H9/
16 C DATA BLANK/1H /, COMMA/1H,/
17 C
18 C
19 C 1 READ(TTYIN,200) BUFFER
20 C
21 C L = N
22 C ARRAY(L) = 0
23 C POWER = 0
24 C
25 C DO 2 I=1,80
26 C K = 81 - I
27 C IF(BUFFER(K) .EQ. BLANK .OR. BUFFER(K) .EQ. COMMA) GOTO 3
28 C
29 C FOUND A CHARACTER, SEE IF IT'S A VALID NUMERIC
30 C
31 C DO 4 J=1,10
32 C II = J - 1
33 C IF(BUFFER(K) .NE. NUMS(J)) GOTO 4
34 C
35 C ITS A NUMBER, ADD IT TO PRESENT SUM
36 C
37 C ARRAY(L) = ARRAY(L) + (II * (10**POWER))
38 C POWER = POWER + 1
39 C GOTO 2
40 C 4 CONTINUE
41 C
42 C COME HERE IF CHARACTER FOUND WAS NOT NUMERIC
43 C
44 C WRITE(TTYOUT,100)
45 C GOTO 1
46 C
47 C FOUND A DELIMITER, SEE IF END OF A NUMBER
48 C
49 C 3 IF(ARRAY(L) .EQ. 0) GOTO 2
50 C L = L - 1
51 C IF(L .EQ. 0) GOTO 5
52 C ARRAY(L) = 0
53 C POWER = 0
54 C 2 CONTINUE
55 C
56 C MAKE SURE NUMBER(S) IS/ARE LESS THAN 99999 SIN
57 C WE DON'T EXCEED 15 FORMATS
58 C

```

```

52      5 DO 6 I=1,N
60      IF (ARRAY(I) .GT. 99999) GOTO 7
61      6 CONTINUE
62      RETURN
63 C
64 C      ERROR MESSAGE
65 C
66      7 WRITE(11,OUT,101)
67      GOTO 1
68 C
69 C      FORMATS
70 C
71      100 FORMAT(37H-***ERROR*** INPUT NOT NUMERIC--RETRY)
72      101 FORMAT(39H-***ERROR*** NUMBER(S) TOO LARGE--RETRY)
73      200 FORMAT(80A1)
74      END

```

```

1 SUBROUTINE IU(N,MAT,INDEX,UNITNO,READ,SYS)
2 C
3 C THIS SUBROUTINE HANDLES ALL PERMFILE I/O
4 C
5 C IT WRITES/READS N (THE NUMBER OF ELEMENTS IN THE MATRIX),
6 C THE SYSTEM MATRIX (MAT), AND THE INDEX SET (INDEX) TO/FROM
7 C A PERMFILE.
8 C
9 C WRITTEN BY: DAVID R. YINGLING, JR
10 C ENGINEERING AND PUBLIC POLICY GROUP
11 C UNIVERSITY OF DAYTON
12 C DAYTON, OHIO 45469
13 C
14 C IMPLICIT INTEGER*2 (A-Z)
15 C INTEGER N, INDEX
16 C LOGICAL READ
17 C LOGICAL*1 MAT(SYS,SYS)
18 C DIMENSION INDEX(SYS)
19 C
20 C CHECK TO SEE IF THIS IS A READ REQUEST
21 C
22 C IF(READ) GOTO 1
23 C
24 C NUP, IT'S A WRITE REQUEST
25 C
26 C REWIND UNITNO
27 C WRITE(UNITNO) N
28 C WRITE(UNITNO) MAT
29 C WRITE(UNITNO) INDEX
30 C GOTO 2
31 C
32 C YUP, IT WAS A READ REQUEST
33 C
34 C 1 REWIND UNITNO
35 C READ (UNITNO) N
36 C READ (UNITNO) MAT
37 C READ (UNITNO) INDEX
38 C
39 C 2 RETURN
40 C END

```

```

1      SUBROUTINE DIGSTG(N,MATRIX,INDEX,TTYOUT,SY)
2 C
3 C      THIS SUBROUTINE DISPLAYS THE DIGRAPH IN STAGES
4 C
5 C      WRITTEN BY:  DAVID R. YINGLING, JR.
6 C                  ENGINEERING AND PUBLIC POLICY GROUP
7 C                  UNIVERSITY OF DAYTON
8 C                  DAYTON, OHIO 45469
9 C
10 C
11 C      VARIABLE NAME          DESCRIPTION
12 C
13 C      N                      THE NUMBER OF ELEMENTS IN THE INPUT
14 C                          MATRIX.
15 C
16 C      MATRIX                INPUT REACHABILITY MATRIX
17 C
18 C      INDEX                 INDEX SET OF THE INPUT MATRIX
19 C
20 C      TTYOUT                FORTRAN WRITE UNIT NUMBER FOR TELETYPE
21 C
22 C      SYS                   DIMENSION SIZES OF SYSTEM MATRICES
23 C
24 C -----
25 C
26 C      MATRXX                SCRATCH MATRIX FOR STAGES ROUTINE
27 C
28 C      INDXX                 INDEX SET OF SCRATCH MATRIX
29 C
30 C      STAGES                VECTOR DENOTING THE NUMBER OF ELEMENTS
31 C                          ON EACH STAGE.  STAGES(I) = THE
32 C                          NUMBER OF ELEMENTS ON STAGE #I.
33 C
34 C      NS                     THE TOTAL NUMBER OF STAGES
35 C
36 C      *****NOTE*****
37 C          THE DIMENSIONS OF "MATRXX", "INDXX", + "STAGES" SHOULD BE
38 C          EQUAL TO "SYS".
39 C
40 C      *****NOTE*****
41 C          IN ORDER TO CONSERVE CORE STORAGE, THE STATE OF THE INPUT
42 C          REACHABILITY MATRIX HAS BEEN DESTROYED.  LET THE PROGRAMMER
43 C          BEWARE!!!!!!
44 C
45 C      IMPLICIT INTEGER*2 (A-Z)
46 C      INTEGER N, INDEX, INDXX
47 C      DIMENSION INDEX(SYS), INDXX(128), STAGES(128)
48 C      LOGICAL*1 MATRIX(SYS,SYS), MATRXX(128,128)
49 C
50 C      CHECK FOR ERROR CONDITION
51 C
52 C      IF(N .LE. 0) GOTO 4
53 C
54 C      STEP 1---TRANPOSE INPUT ORIGINAL REACHABILITY MATRIX
55 C
56 C      DO 1 I=1,N
57 C      DO 1 J=1,N
58 C      MATRXX(I,J) = MATRIX(J,I)

```

```

59      1 CONTINUE
60 C
61      DO 2 I=1,N
62      DO 2 J=1,N
63      MATRIX(I,J) = MATRXX(I,J)
64      2 CONTINUE
65 C
66 C      STEP 2---LEVELS PARTITION "MATRIX", LEAVE RESULT IN "MATRXX"
67 C
68      CALL HIERCH(N,MATRIX,INDEX,MATRXX,INDXX,NS,STAGES,SYS)
69 C
70 C      STEP 3---PUT "MATRXX" INTO STANDARD FORM
71 C
72      CALL STAN(N,MATRXX,INDXX,NS,STAGES,SYS)
73 C
74 C      STEP 4-- CALCULATE CONDENSATION MATRIX
75 C
76      CALL CONDE(N,MATRXX,INDXX,STAGES,TTYOUT,.TRUE.,SYS)
77 C
78 C      STEP 5---CALCULATE SKELETON MATRIX
79 C
80      CALL SKLTN(N,MATRXX,SYS)
81 C
82 C      STEP 6---TRANSPose SKELETON MATRIX TO PUT INTO UPPER
83 C      TRIANGULAR FORM
84 C
85      DO 3 I=1,N
86      DO 3 J=1,N
87      MATRIX(I,J) = MATRXX(J,I)
88      3 CONTINUE
89 C
90 C      STEP 7---PRINT OUT STAGE DIGRAPH
91 C
92      CALL DISPST(N,MATRIX,INDXX,STAGES,NS,TTYOUT,SYS)
93      RETURN
94 C
95 C      ERROR MESSAGE
96 C
97      4 WRITE(TTYOUT,100)
98      RETURN
99 C
100 C      FORMAT
101 C
102      100 FORMAT(42H-***ERROR*** NO STRUCTURE CURRENTLY EXISTS)
103      END

```



```

1      SUBROUTINE DISPST(N,MATRIX,INDEX,STAGES,NSTAGE,TTYOUT,SYN)
2
3      THIS SUBROUTINE TAKES AN UPPER TRIANGULAR SKELETON MATRIX
4      AND PRINTS OUT A STAGES DIGRAPH
5
6      WRITTEN BY:  DAVID R. YINGLING, JR.
7                  ENGINEERING AND PUBLIC POLICY GROUP
8                  UNIVERSITY OF DAYTON
9                  DAYTON, OHIO  45469
10
11
12      VARIABLE NAME          DESCRIPTION
13
14      N                      NUMBER OF ELEMENTS IN INPUT SKELETON
15                              MATRIX.
16
17      MATRIX                 INPUT SKELETON MATRIX
18
19      INDEX                  INDEX SET OF SKELETON MATRIX
20
21      STAGES                 INPUT VECTOR DENOTING NUMBER OF
22                              ELEMENTS ON EACH STAGE.  STAGES(I) =
23                              THE NUMBER OF ELEMENTS ON STAGE #I
24
25      TTYOUT                 FORTRAN WRITE UNIT NUMBER FOR TELETYPE
26
27      SYN                    DIMENSION SIZES OF SYSTEM MATRICES
28
29      -----
30
31      LIST                   SCRATCH VECTOR WHICH CONTAINS THE
32                              INDEX NUMBER OF ELEMENTS RELATED TO
33                              ELEMENT #I
34
35      COUNT                  NUMBER OF ELEMENTS IN "LIST".
36
37      ROW                    CURRENT ROW BEING PROCESSED
38
39      STAGE                  NUMBER OF THE CURRENT STAGE BEING
40                              PROCESSED
41
42      NSTAGE                 THE TOTAL NUMBER OF STAGES
43
44      ISTART                 STARTING SEARCH INDEX FOR UPPER
45                              TRIANGULAR MATRIX
46
47      *****NOTE*****
48      THE DIMENSIC          "LIST" SHOULD BE EQUAL TO "SYN".
49
50
51      IMPLICIT INTEGER*2 (A-Z)
52      INTEGER N, INDEX
53      DIMENSION INDEX(SYN), STAGES(SYN), LIST(128)
54      LOGICAL*1 MATRIX(SYN,SYN)
55
56      BEGIN PROCESSING: INITIALIZE
57
58      ROW      = 1

```

```

59      STAGE = 0
60      I      = 1
61 C
62 C      BEGIN FINDING ELEMENTS
63 C
64      1 STAGE = STAGE + 1
65      ROW   = ROW + STAGES(STAGE)
66      WRITE(TTYOUT,100) STAGE
67 C
68 C      SEE IF WE ARE PROCESSING LAST STAGE
69 C
70      4 IF(STAGE .EQ. NSTAGE) GOTO 2
71 C
72 C      PROCESS ELEMENTS FOR STAGE# STAGE
73 C
74      COUNT = 0
75      ISTART = I + 1
76      DO 3 J=ISTART,N
77 C
78 C      FIND ALL ELEMENTS THAT ELEMENT# I REACHES TO
79 C
80      IF(.NOT. MATRIX(I,J)) GOTO 3
81 C
82 C      FOUND ONE, KEEP A RECORD OF IT
83 C
84      COUNT      = COUNT + 1
85      LIST(COUNT) = INDEX(J)
86      3 CONTINUE
87      IF(COUNT .EQ. 0) GOTO 2
88 C
89 C      ALL DONE WITH ELEMENT #I, PRINT OUT
90 C
91      WRITE(TTYOUT,101) INDEX(I), (LIST(I), I=1,COUNT)
92 C
93      5 I = I + 1
94 C
95 C      ALL DONE WITH STAGES PRINTOUT????
96 C
97      IF(I .GT. N) RETURN
98 C
99 C      ALL DONE WITH THIS STAGE????
100 C
101      IF(I .EQ. ROW) GOTO 1
102 C
103 C      CONTINUE PROCESSING THIS STAGE
104 C
105      GOTO 4
106 C
107 C      PROCESS LAST STAGE
108 C
109      2 WRITE(TTYOUT,102) INDEX(I)
110      GOTO 5
111 C
112 C      FORMATS
113 C
114      100 FORMAT(14-,10X,11HSTAGE NO. ,13/1H0)
115      101 FORMAT(11X,15,3H =>,18(7(15,1H,1/20X))
116      102 FORMAT(11X,15)

```

```

1 SUBROUTINE ELCONT(N,REA,INDEX,TTYIN,TTYOUT,SYS),
2 C
3 C THIS SUBROUTINE PERFORMS THE ELEMENTARY CONTRACTION PROCESS
4 C
5 C EDITED BY: DAVID R. YINGLING, JR.
6 C ENGINEERING AND PUBLIC POLICY GROUP
7 C UNIVERSITY OF DAYTON
8 C DAYTON, OHIO 45469
9 C
10 C IMPLICIT INTEGER*2 (A-Z)
11 C INTEGER N, INDEX, INDH, NN
12 C LOGICAL FOUND1, FOUND2
13 C LOGICAL*1 REA(SYS, SYS), REAH(128,128)
14 C DIMENSION INUEX(SYS), INDH(128), NUMS(2), LO(128)
15 C
16 C VARIABLE NAME DESCRIPTION
17 C
18 C N THE NUMBER OF ELEMENTS IN "REA"
19 C
20 C REA ARGUMENT REACHABILITY MATRIX
21 C
22 C INDEX INDEX SET OF "REA"
23 C
24 C TTYIN FORTRAN READ UNIT NUMBER FOR TELETYPE
25 C
26 C TTYOUT FORTRAN WRITE UNIT NUMBER FOR TELETYPE
27 C
28 C SYS DIMENSION SIZES OF SYSTEM MATRICES
29 C
30 C -----
31 C
32 C REAH SCRATCH REACHABILITY MATRIX
33 C
34 C INDH INDEX SET OF "REAH"
35 C
36 C LO SCRATCH VECTOR
37 C
38 C IU POSITION OF ELEMENT U ON "REA"
39 C
40 C IV POSITION OF ELEMENT V ON "REA"
41 C
42 C IIU POSITION OF ELEMENT U ON "REAH"
43 C
44 C IIV POSITION OF ELEMENT V ON "REAH"
45 C
46 C NEWNAM THE NEW INDEX VALUE FOR THE CONTRACTED
47 C ELEMENTS.
48 C
49 C U ELEMENT # U TO BE CONTRACTED
50 C
51 C V ELEMENT #V TO BE CONTRACTED
52 C
53 C
54 C *****NOTE*****
55 C THE DIMENSIONS OF "INDH", "REAH",+"LO" SHOULD BE EQUAL TO "SYS"
56 C
57 C
58 C ASK FOR TWO ELEMENTS

```

```

53 C
60 1 WRITE(TTYOUT,100)
61 CALL GETNUM(NUMS,2,TTYIN,TTYOUT)
62 U = NUMS(1)
63 V = NUMS(2)
64 IF(U .EQ. 0 .OR. V .EQ. 0) GOTO 2
65 C
66 C CHECK TO SEE THAT U AND V ARE IN SYSTEM INDEX
67 C
68 CALL FINDIT(N,U,V,IU,IV,INDEX,FOUND1,FOUND2,SYS)
69 IF(FOUND1 .AND. FOUND2) GOTO 3
70 C
71 C U AND/OR V DOESN'T EXIST
72 C
73 IF(.NOT. FOUND1) WRITE(TTYOUT,101) U
74 IF(.NOT. FOUND2) WRITE(TTYOUT,101) V
75 GOTO 1
76 C
77 C DOES U REACH TO V ?
78 C
79 3 NN = N
80 CALL HIERCH(NN,REA,INDEX,REAH,INDH,NL,LO,SYS)
81 CALL STAH(NN,REAH,INDH,NL,LO,SYS)
82 CALL CONDE(NN,REAH,INDH,LO,TTYOUT,.FALSE.,SYS)
83 CALL SKLTN(NN,REAH,SYS)
84 CALL FINDIT(NN,U,V,IU,IV,INDH,FOUND1,FOUND2,SYS)
85 C
86 IF(FOUND1 .AND. FOUND2 .AND. REAH(IU,IV)) GOTO 4
87 C
88 C U IS NOT ADJACENT TO V, ISSUE ERROR MESSAGE
89 C
90 WRITE(TTYOUT,102) U, V
91 GOTO 1
92 C
93 C OK, GET NEW NAME
94 C
95 4 WRITE(TTYOUT,103)
96 CALL GETNUM(NUMS,1,TTYIN,TTYOUT)
97 NEWNAM = NUMS(1)
98 IF(NEWNAM .EQ. 0) GOTO 2
99 C
100 C CHECK TO SEE THAT NEWNAM IS NOT IN SYSTEM INDEX
101 C
102 CALL FINDIT(N,NEWNAM,1,1,1,INDEX,FOUND1,FOUND2,SYS)
103 IF(.NOT. FOUND1) GOTO 5
104 C
105 C NEWNAM IS IN THE INDEX SET, ISSUE ERROR MSG
106 C
107 WRITE(TTYOUT,104) NEWNAM
108 GOTO 4
109 C
110 C OK, NOW CHANGE MATRIX.
111 C
112 5 CALL COMBIN(N,REA,INDEX,IU,IV,NEWNAM,SYS)
113 CALL TRNCLS(N,REA,SYS)
114 CALL ID(N,REA,INDEX,10,.FALSE.,SYS)
115 GOTO 1
116 C

```

117 2 RETURN
118 C
119 C FORMATS
120 C
121 100 FORMAT(38H-TYPE TWO ELEMENTS TO BE CONTRACTED ?)
122 101 FORMAT(12H-***ERROR***,15,20H NOT IN SYSTEM INDEX)
123 102 FORMAT(12H-***ERROR***,15,19H IS NOT ADJACENT TO,15)
124 103 FORMAT(20H NEW INDEX NUMBER ?)
125 104 FORMAT(12H-***ERROR***,15,24H ALREADY IN SYSTEM INDEX)
125 END

332

```

1 SUBROUTINE POOLIN,REA,INDEX,TTYIN,TTYOUT,SYSD
2 C
3 C THIS SUBROUTINE PERFORMS THE POOLING OPERATION
4 C
5 C EDITED BY: DAVID R. YINGLING, JR.
6 C ENGINEERING AND PUBLIC POLICY GROUP
7 C UNIVERSITY OF DAYTON
8 C DAYTON, OHIO 45469
9 C
10 C
11 C VARIABLE NAME DESCRIPTION
12 C
13 C N NUMBER OF ELEMENTS IN "REA"
14 C
15 C REA ARGUMENT REACHABILITY MATRIX
16 C
17 C INDEX INDEX SET FOR "REA"
18 C
19 C TTYIN FORTRAN READ UNIT NUMBER FOR TELETYPE
20 C
21 C TTYOUT FORTRAN WRITE UNIT NUMBER FOR TELETYPE
22 C
23 C SYS DIMENSION SIZES OF SYSTEM MATRICES
24 C
25 C -----
26 C
27 C END ENDING SUBSCRIPT OF LEVEL OR STAGE #1
28 C
29 C IU POSITION OF ELEMENT U ON "REA"
30 C
31 C IV POSITION OF ELEMENT V ON "REA"
32 C
33 C INDH INDEX SET FOR "REAH"
34 C
35 C LO SCRATCH VECTOR, LO(I) = NUMBER
36 C OF ELEMENTS ON LEVEL # I (OR STAGE
37 C # I IFF "STAGES" = .TRUE.)
38 C
39 C NEWNAM THE NEW INDEX VALUE FOR THE
40 C POOLED ELEMENTS
41 C
42 C REAH SCRATCH REACHABILITY MATRIX
43 C
44 C STAGES LOGICAL VARIABLE IF WHEN .TRUE.,
45 C THE STAGES ARE BEING EXAMINED, WHEN
46 C .FALSE., THE LEVELS ARE BEING
47 C EXAMINED
48 C
49 C START STARTING SUBSCRIPT OF LEVEL OR STAGE #1
50 C
51 C U ELEMENT #1 TO BE POOLED
52 C
53 C V ELEMENT #2 TO BE POOLED
54 C
55 C IMPLICIT INTEGER*2 (A-Z)
56 C INTEGER N, INDEX, INDH
57 C LOGICAL FOUND1, FOUND2, STAGES
58 C LOGICAL*1 REA(SYS,SYS), REAH(128,128)

```

```

59 DIMENSION INDEX(SYS), INDH(128), LO(128), NUMS(2)
60 C
61 C *****NOTE*****
62 C THE DIMENSIONS OF "REAH", "INDH", + "LO" SHOULD BE EQUAL TO "SYS".
63 C
64 C
65 C ASK USER FOR ELEMENTS TO BE POOLED
66 C
67 1 WRITE(TTYOUT,100)
68 CALL GETNUM(NUMS,2,TTYIN,TTYOUT)
69 U = NUMS(1)
70 V = NUMS(2)
71 IF(U .EQ. 0 .OR. V .EQ. 0) GOTO 2
72 C
73 C SEE IF U AND V EXIST IN SYSTEM INDEX
74 C
75 CALL FINDIT(N,U,V,IU,IV,INDEX,FOUND1,FOUND2,SYS)
76 IF(FOUND1 .AND. FOUND2) GOTO 12
77 C
78 C U AND/OR V DOESN'T EXIST
79 C
80 IF(.NOT. FOUND1) WRITE(TTYOUT,101) U
81 IF(.NOT. FOUND2) WRITE(TTYOUT,101) V
82 GOTO 1
83 C
84 C CHECK FOR U AND V ON SAME LEVEL
85 C
86 12 STAGES = .FALSE.
87 CALL HIERCH(N,REA,INDEX,REAH,INDH,NL,LO,SYS)
88 C
89 C USE "LO" TO DETERMINE IF U AND V ARE ON SAME LEVEL OR STAGE.
90 C
91 11 FOUND2 = .FALSE.
92 FOUND1 = .FALSE.
93 START = 1
94 C
95 DO 5 I=1,NL
96 END = START + LO(I) - 1
97 C
98 DO 6 J=START,END
99 IF(INDH(J) .EQ. U) FOUND1 = .TRUE.
100 IF(INDH(J) .EQ. V) FOUND2 = .TRUE.
101 6 CONTINUE
102 C
103 IF(FOUND1 .AND. FOUND2) GOTO 3
104 C
105 IF(FOUND1 .OR. FOUND2) GOTO 7
106 C
107 START = END + 1
108 5 CONTINUE
109 C
110 C NOT ON SAME LEVEL, SEE IF ON SAME STAGE
111 C
112 7 IF(STAGES) GOTO 13
113 C
114 C TRANSPOSE MODEL MATRIX IN ORDER TO FOOL "HIERCH" SUBROUTINE
115 C
116 DO 8 I=1,N

```

```

117      DO 8 J=1,N
118      REAH(1,J) = REA(J,1)
119      8 CONTINUE
120 C
121      DO 9 I=1,N
122      DO 9 J=1,N
123      REA(1,J) = REAH(1,J)
124      9 CONTINUE
125 C
126      CALL HIERCHIN(REA,INDEX,REAH,INDH,HL,LD,SY3)
127 C
128 C      TRANSPOSE MATRIX BACK SO IT IS NOT SCRAMBLED
129 C
130      DO 10 I=1,N
131      DO 10 J=1,N
132      REA(1,J) = REAH(1,J)
133      10 CONTINUE
134      STAGES = .TRUE.
135      GOTO 11
136 C
137 C      POOLING ERROR MESSAGE
138 C
139      13 WRITE(TTYOUT,104) U, V
140      GOTO 1
141 C
142 C      GET PROPER "POSITIONS" FOR ELEMENTS U AND V ON "REA".
143 C
144      3 CALL FINDIT(N,U,V,IU,IV,INDEX,FOUND1,FOUND2,SY3)
145 C
146 C      ASK FOR NEW NAME
147 C
148      WRITE(TTYOUT,102)
149      CALL GETNUM(NUMS,1,TTYIN,TTYOUT)
150      NEWNAM = NUMS(1)
151      IF(NEWNAM .EQ. 0) GOTO 2
152 C
153 C      MAKE SURE NEWNAM DOESN'T EXIST IN SYSTEM INDEX
154 C
155      CALL FINDIT(N,NEWNAM,J,J,J,INDEX,FOUND1,FOUND2,SY3)
156      IF(.NOT. FOUND1) GOTO 4
157 C
158 C      NEWNAM IS IN THE INDEX SET, ISSUE ERROR MESSAGE
159 C
160      WRITE(TTYOUT,103) NEWNAM
161      GOTO 3
162 C
163 C      OK, CHANGE THE MATRIX
164 C
165      4 CALL COMBIN(N,REA,INDEX,IU,IV,NEWNAM,SY3)
166      CALL TRNCLS(N,REA,SY3)
167      GOTO 1
168 C
169      2 RETURN
170 C
171 C      -FORMATS
172 C
173      100 FORMAT(34H-TYPE TWO ELEMENTS TO BE POOLED ? )
174      101 FORMAT(12H-***ERROR***,15,20H NOT IN SYSTEM INDEX)

```


175 102 FORMAT(20H NEW INDEX NUMBER 7)
176 103 FORMAT(12H-***ERROR***,15,24H ALREADY IN SYSTEM INDEX)
177 104 FORMAT(12H-***ERROR***,15,4H AND,15,35H ARE NOT ON THE SAME LEVEL
178 +OR STAGE)
179 END

333

```

1 SUBROUTINE COMBIN(N,REA,INDEX,IU,IV,NEWMAM,SY)
2 C
3 C THIS SUBROUTINE COMBINES IU AND IV IN REA
4 C
5 C EDITED BY: DAVID R. YINGLING, JR.
6 C ENGINEERING AND PUBLIC POLICY GROUP
7 C UNIVERSITY OF DAYTON
8 C DAYTON, OHIO 45469
9 C 513-228-2238
10 C
11 C
12 C VARIABLE NAME DESCRIPTION
13 C
14 C N NUMBER OF ELEMENTS IN "REA"
15 C
16 C REA ARGUMENT REACHABILITY MATRIX
17 C
18 C INDEX INDEX SET FOR "REA"
19 C
20 C IU SUBSCRIPT #1 TO BE COMBINED
21 C
22 C IV SUBSCRIPT #2 TO BE COMBINED
23 C
24 C NEWMAM THE INDEX VALUE FOR THE COMBINED
25 C ELEMENTS
26 C
27 C SYS DIMENSION SIZES OF SYSTEM MATRICES
28 C
29 C IMPLICIT INTEGER*2 (A-Z)
30 C INTEGER N, INDEX
31 C LOGICAL*1 REA(SYS,SY)
32 C DIMENSION INDEX(SYS)
33 C
34 C REPLACE ROW V WITH THE BOOLEAN SUM OF U AND V
35 C
36 C DO 1 J=1,N
37 C REA(IV,J) = REA(IU,J) .OR. REA(IV,J)
38 C 1 CONTINUE
39 C
40 C REPLACE COL V WITH THE BOOLEAN SUM OF U AND V
41 C
42 C DO 2 I=1,N
43 C REA(I,IV) = REA(I,IU) .OR. REA(I,IV)
44 C 2 CONTINUE
45 C
46 C REPLACE V'S INDEX WITH NEWMAM
47 C
48 C INDEX(IV) = NEWMAM
49 C
50 C ERASE ROW, COL, AND INDEX FOR U
51 C
52 C CALL E1IM(N,REA,INDEX,IU,SY)
53 C RETURN
54 C END

```

```

1      SUBROUTINE TRNCLS(N,MATRIX,SYS)
2 C
3 C      THIS SUBROUTINE WILL TRANSITIVELY CLOSE THE INPUT MATRIX
4 C
5 C      WRITTEN BY: RAYMOND L. FITZ, S.M.
6 C      EDITED BY: DAVID R. YINGLING, JR.
7 C      ENGINEERING AND PUBLIC POLICY GROUP
8 C      UNIVERSITY OF DAYTON
9 C      DAYTON, OHIO 45469
10 C
11 C      VARIABLE NAME      DESCRIPTION
12 C
13 C      N                  THE NUMBER OF ELEMENTS IN "MATRIX"
14 C
15 C      MATRIX             INPUT ADJACENCY MATRIX/OUTPUT
16 C                        REACHABILITY MATRIX .
17 C
18 C      SYS                DIMENSION SIZES OF SYSTEM MATRICES
19 C
20 C -----
21 C
22 C      NONES              NUMBER OF ONES IN THE REACHABILITY
23 C                        SET OF ELEMENT #1
24 C
25 C      LAST               NUMBER OF ONES IN THE REACHABILITY
26 C                        SET OF ELEMENT #1 FROM LAST
27 C                        COMPUTATION
28 C
29 C      RECORD             VECTOR USED TO KEEP ACCOUNT
30 C                        OF ONES IN THE REACHABILITY SET
31 C                        OF ELEMENT #1
32 C
33 C      IMPLICIT INTEGER*2 (A-Z)
34 C      INTEGER N
35 C      LOGICAL*1 MATRIX(SYS,SYS)
36 C      DIMENSION RECORD(128)
37 C
38 C      *****NOTE*****
39 C      THE DIMENSION OF "RECORD" SHOULD BE EQUAL TO "SYS".
40 C
41 C      INITIALIZE PROCEDURE
42 C
43 C      DO 1 I=1,N
44 C
45 C      PROCESS ELEMENT #I
46 C
47 C      NONES = 0
48 C      LAST = 0
49 C
50 C      FIND ALL ONES IN REACHABILITY SET OF ELEMENT #I
51 C
52 C      2 DO 3 K=1,N
53 C      IF(.NOT. MATRIX(I,K)) GOTO 3
54 C
55 C      FOUND A ONE, KEEP A RECORD OF IT
56 C
57 C      NONES      = NONES + 1
58 C      RECORD(NONES) = K

```

```

53 C
60 3 CONTINUE
61 C
62 C CHECK TO SEE IF ANY NEW ONES WERE ADDED FROM LAST
63 C TIME THROUGH
64 C
65 IF(NONES .EQ. LAST) GOTO 1
66 C
67 C NO, COMPUTE NEW ELEMENTS IN REACHABILITY SET III
68 C ELEMENT #I BY TRANSITIVITY
69 C
70 LAST = NONES
71 C
72 C THE NEXT GROUP OF CODE PERFORMS THE PROCESS:
73 C IF MATRIX(I,K)=1 .AND. MATRIX(K,J)=1,
74 C THEN MATRIX(I,J)=1.
75 C
76 DO 4 L=1,NONES
77 K = RECORD(L)
78 DO 4 J=1,N
79 MATRIX(I,J) = MATRIX(I,J) .OR. MATRIX(K,J).
80 4 CONTINUE
81 C
82 C KEEP GOING UNTIL ALL REACHABILITY SETS HAVE BEEN EXAMINED
83 C
84 NONES = 0
85 GOTO 2
86 C
87 1 CONTINUE
88 RETURN
89 END

```

```

1 SUBROUTINE AUEGE(N,REA,INDEX,TTYIN,TTYOUT,SY)
2 C
3 C THIS SUBROUTINE ADDS AN EDGE ON THE MINIMUM EDGE DIGRAPH
4 C
5 C WRITTEN BY: DAVID R. YINGLING, JR.
6 C ENGINEERING AND PUBLIC POLICY GROUP
7 C UNIVERSITY OF DAYTON
8 C DAYTON, OHIO 45469
9 C 513-228-2238
10 C
11 C VARIABLE NAME DESCRIPTION
12 C
13 C N NUMBER OF ELEMENTS IN "REA"
14 C
15 C REA ARGUMENT REACHABILITY MATRIX
16 C
17 C INDEX INDEX SET FOR "REA"
18 C
19 C TTYIN FORTRAN READ UNIT NUMBER FOR TELETYPE
20 C
21 C TTYOUT FORTRAN WRITE UNIT NUMBER FOR TELETYPE
22 C
23 C SYS DIMENSION SIZES OF SYSTEM MATRICES
24 C
25 C -----
26 C
27 C N1 INDEX VALUE OF ORIGINATING EDGE ELEMENT
28 C
29 C N2 INDEX VALUE OF DESTINATION EDGE
30 C ELEMENT
31 C
32 C II SUBSCRIPT OF "N1" ON "REA"
33 C
34 C JJ SUBSCRIPT OF "N2" ON "REA"
35 C
36 C IMPLICIT INTEGER*2 (A-Z)
37 C INTEGER N, INDEX
38 C LOGICAL FOUND1, FOUND2
39 C LOGICAL*1 REA(SYS,SY)
40 C DIMENSION INDEX(SYS), NUMS(2)
41 C
42 C ASK USER FOR A REACHES TO B
43 C
44 C WRITE(TTYOUT,100)
45 C 1 CALL GETNUM(NUMS,2,TTYIN,TTYOUT)
46 C N1 = NUMS(1)
47 C N2 = NUMS(2)
48 C IF(N1 .EQ. 0 .OR. N2 .EQ. 0) GOTO 2
49 C
50 C CHECK TO SEE IF N1 AND N2 ARE IN SYSTEM INDEX
51 C
52 C CALL FINDIT(N,N1,N2,II,JJ,INDEX,FOUND1,FOUND2,SY)
53 C IF(FOUND1 .AND. FOUND2) GOTO 3
54 C
55 C N1 AND/OR N2 NOT IN SYSTEM INDEX
56 C
57 C IF(.NOT. FOUND1) WRITE(TTYOUT,101) N1
58 C IF(.NOT. FOUND2) WRITE(TTYOUT,101) N2

```

```

50      GOTO 2
60 C
61 C      PUT IN EDGE
62 C
63      3 REA(II,JJ) = .TRUE.
64      GOTO 1
65 C
66 C      PERFORM TRANSITIVE CLOSURE
67 C
68      2 CALL TRNCLSN,REA,SYS)
69      RETURN..
70 C
71 C      FORMATS
72 C
73      100 FORMAT(22H-TYPE <A REACHES TO B>)
74      101 FORMAT(12H-***ERROR***,15,20H NOT IN SYSTEM INDEX)
75      END

```

```

1 SUBROUTINE EREDGE(N,REA,INDEX,TTYIN,TTYOUT,SYN)
2 C
3 C THIS SUBROUTINE ERASES AND EDGE ON THE MINIMUM EDGE DIGRAPH
4 C
5 C WRITTEN BY: DAVID R. YINGLING, JR.
6 C DECISIONS SYSTEMS LAB
7 C ENGINEERING AND PUBLIC POLICY GROUP
8 C UNIVERSITY OF DAYTON
9 C DAYTON, OHIO 45469
10 C
11 C
12 C
13 C N NUMBER OF ELEMENTS IN "REA"
14 C
15 C REA ARGUMENT REACHABILITY MATRIX
16 C
17 C INDEX INDEX SET FOR "REA"
18 C
19 C TTYIN FORTRAN READ UNIT NUMBER FOR TELETYPE
20 C
21 C TTYOUT FORTRAN WRITE UNIT NUMBER FOR TELETYPE
22 C
23 C SYN DIMENSION SIZES OF SYSTEM MATRICES
24 C
25 C -----
26 C
27 C N1 INDEX VALUE OF ORGINATION OF EDGE TO
28 C BE ERASED
29 C
30 C N2 INDEX VALUE OF DESTINATION OF EDGE
31 C TO BE ERASED
32 C
33 IMPLICIT INTEGER*2 (A-Z)
34 INTEGER N, INDEX, INDH, NN
35 LOGICAL FOUND1, FOUND2
36 LOGICAL*1 REA(SYN,SYN), REAH(128,128)
37 DIMENSION INDEX(SYN), INDH(128), LO(128), NUMS(2)
38 C
39 C *****NOTE*****
40 C THE DIMENSIONS OF "LO", "INDH", + "REAH" SHOULD BE EQUAL TO "SYN".
41 C
42 C
43 C ASK USER FOR A REACHES TO B TO BE ERASED
44 C
45 C WRITE(TTYOUT,100)
46 C 1 CALL GETNUM(NUMS,2,TTYIN,TTYOUT)
47 C N1 = NUMS(1)
48 C N2 = NUMS(2)
49 C IF(N1 .EQ. 0 .OR. N2 .EQ. 0) GOTO 2
50 C
51 C CHECK TO SEE IF N1 AND N2 ARE IN SYSTEM INDEX
52 C
53 C CALL FINDIT(N,N1,N2,11,JJ,INDEX,FOUND1,FOUND2,SYN)
54 C IF(FOUND1 .AND. FOUND2) GOTO 3
55 C
56 C N1 AND/OR N2 NOT IN SYSTEM INDEX, ISSUE ERROR MSG
57 C
58 C IF(.NOT. FOUND1) WRITE(TTYOUT,101) N1

```

```

59 IF(.NOT. FOUND2) WRITE(TTYOUT,101) N2
60 RETURN
61 C
62 C OK, NOW CHECK FOR CYCLES
63 C
64 C VARIABLE NAME DESCRIPTION
65 3 IF(.NOT. (REA(II,JJ) .AND. REA(JJ,II))) GOTO 4
66 C
67 C WHOOPS, N1 AND N2 ARE IN A CYCLE, ISSUE ERROR MSG
68 C
69 WRITE(TTYOUT,102) N1, N2, N1
70 GOTO 2
71 C
72 C OK, CALCULATE SKELETON MATRIX
73 C
74 4 NN = N
75 CALL HIERCH(NN,REA,INDEX,REAH,INDH,NL,LO,SYS)
76 CALL STAH(NN,REAH,INDH,NL,LO,SYS)
77 CALL CONDE(NN,REAH,INDH,LO,TTYOUT,.FALSE.,SYS)
78 CALL SKLTN(NN,REAH,SYS)
79 C
80 C CHECK TO SEE IF N1 AND N2 ARE ON MINIMUM EDGE DIGRAPH
81 C
82 CALL FINDIT(NN,N1,N2,II,JJ,INDH,FOUND1,FOUND2,SYS)
83 IF(FOUND1 .AND. FOUND2) GOTO 5
84 C
85 C N1 AND/OR N2 WAS NOT ON MINIMUM EDGE DIGRAPH
86 C
87 IF(.NOT. FOUND1) WRITE(TTYOUT,103) N1
88 IF(.NOT. FOUND1) WRITE(TTYOUT,103) N2
89 GOTO 2
90 C
91 C CHECK TO SEE IF THE EDGE FROM N1 TO N2 EXISTS ON MINIMUM
92 C EDGE DIGRAPH
93 C
94 5 IF(REA(II,JJ)) GOTO 6
95 C
96 C THE EDGE FROM N1 TO N2 WAS NOT ON MINIMUM EDGE DIGRAPH
97 C ISSUE ERROR MESSAGE
98 C
99 WRITE(TTYOUT,104) N1, N2
100 GOTO 2
101 C
102 C OK, N1 REACHES TO N2 ON MINIMUM EDGE DIGRAPH
103 C
104 C ELIMINATE REACHABILITY BY DISCONNECTING ANTECEDENT SET
105 C OF N1 FROM THE REACHABILITY SET OF N2
106 C
107 6 CALL FINDIT(N,N1,N2,II,JJ,INDEX,FOUND1,FOUND2,SYS)
108 C
109 C REMOVE EDGE FROM L TO K IF:
110 C 1. L IS A MEMBER OF THE ANTECEDENT SET
111 C OF N1
112 C
113 C AND
114 C
115 C 2. K IS A MEMBER OF THE REACHABILITY
116 C SET OF N2

```



```

117 C
118 DO 7 L=1,N
119 IF(.NOT. REA(L,11)) GOTO 7
120 DO 8 K=1,N
121 IF(.NOT. REA(JJ,K)) GOTO 8
122 REA(L,K) = .FALSE.
123 8 CONTINUE
124 7 CONTINUE
125 C
126 C CALL TRANSITIVE CLOSURE
127 C
128 CALL TRNCLS(N,REA,SYS)
129 CALL IO(N,REA,INDEX,10,.FALSE.,SYS)
130 GOTO 1
131 C
132 C FORMATS
133 C
134 100 FORMAT(35H-TYPE <A REACHES TO B> TO BE ERASED)
135 101 FORMAT(12H-***ERROR***,15,20H NOT IN SYSTEM INDEX)
136 102 FORMAT(12H-***ERROR***,15,4H AND,15,16H ARE IN A CYCLE./7H "ELIM",
137 +15,36H AND RE-ENTER USING THE "80" COMMAND)
138 103 FORMAT(12H-***ERROR***,15,31H IS NOT ON MINIMUM EDGE DIGRAPH)
139 104 FORMAT(26H-***ERROR*** THE EDGE FROM,15,3H TO,15,9H DOES NOT/30H E
140 +XIST ON MINIMUM EDGE DIGRAPH)
141 END

```

```

1 SUBROUTINE ERASE(N,MATRIX,INDEX,TTYIN,TTYOUT,115)
2
3 C THIS SUBROUTINE ERASES AN ELEMENT FROM "MATRIX".
4 C
5 C WRITTEN BY: DAVID R. YINGLING, JR.
6 C ENGINEERING AND PUBLIC POLICY GROUP
7 C UNIVERSITY OF DAYTON
8 C DAYTON, OHIO 45469
9 C
10 C VARIABLE NAME DESCRIPTION
11 C
12 C N THE NUMBER OF ELEMENTS IN "MATRIX"
13 C
14 C MATRIX THE MATRIX TO ERASE THE ELEMENT FROM
15 C
16 C INDEX THE INDEX SET OF "MATRIX".
17 C
18 C TTYIN FORTRAN READ UNIT NUMBER FOR TELETYPE
19 C
20 C TTYOUT FORTRAN WRITE UNIT NUMBER FOR TELETYPE
21 C
22 C SYS DIMENSION SIZES OF SYSTEM MATRICES
23 C
24 C
25 C IMPLICIT INTEGER*2 (A-Z)
26 C INTEGER N, INDEX
27 C LOGICAL*1 MATRIX(SYS,SYS)
28 C LOGICAL FOUND1, FOUND2
29 C DIMENSION INDEX(SYS), NUMS(1)
30 C
31 C IF(N .LE. 0) GOTO 4
32 C 1 WRITE(TTYOUT,100)
33 C 2 CALL GETNUM(NUMS,1,TTYIN,TTYOUT)
34 C I = NUMS(1)
35 C IF(I .EQ. 0) RETURN
36 C
37 C CHECK TO SEE IF "I" IS A MEMBER OF THE INDEX SET
38 C
39 C CALL FINDIT(N,I,J,II,J,INDEX,FOUND1,FOUND2,SYS)
40 C IF(.NOT. FOUND1) GOTO 3
41 C
42 C NUMBER WAS VALID ELEMENT NUMBER, ERASE FROM MATRIX
43 C
44 C CALL ELIM(N,MATRIX,INDEX,II,SYS)
45 C IF(N .LE. 0) GOTO 4
46 C
47 C GO GET ANOTHER
48 C
49 C GOTO 2
50 C
51 C ERROR PRINTOUT
52 C
53 C 3 WRITE(TTYOUT,101) I
54 C GOTO 1
55 C
56 C USER HAS DELETED ALL ELEMENTS
57 C
58 C 4 WRITE(TTYOUT,102)

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59 RETURN
60 C
61 C FORMATS
62 C
63 100 FORMAT(36H-TYPE ELEMENT NUMBERS TO BE ERASED ?)
64 101 FORMAT(12H-***ERROR***,15,20H NOT IN SYSTEM INDEX)
65 102 FORMAT(53H-***NOTE*** THERE ARE NO ELEMENTS IN THE MODEL MATRIX/
66 +26H "ELIM" COMMAND TERMINATED)
67 END

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```

1 SUBROUTINE ADEL(N,MATRIX,INDEX,TTYIN,TTYOUT,SY)
2
3 THIS SUBROUTINE ADDS ELEMENTS TO THE MATRIX "MATRIX"
4 AND PUTS A "ONE" ON THE MAIN DIAGONAL.
5
6 WRITTEN BY: DAVID R. YINGLING, JR.
7 ENGINEERING AND PUBLIC POLICY GROUP
8 UNIVERSITY OF DAYTON
9 DAYTON, OHIO 45469
10
11 VARIABLE NAME DESCRIPTION
12
13 N THE NUMBER OF ELEMENTS IN "MATRIX"
14
15 MATRIX THE MATRIX TO ADD ELEMENTS TO
16
17 INDEX THE INDEX SET OF "MATRIX".
18
19 TTYIN FORTRAN READ UNIT NUMBER FOR TELETYPE
20
21 TTYOUT FORTRAN WRITE UNIT NUMBER FOR TELETYPE
22
23 SYS DIMENSION SIZES OF SYSTEM MATRICES
24
25
26 IMPLICIT INTEGER*2 (A-Z)
27 INTEGER N, INDEX
28 LOGICAL*1 MATRIX(SYS,SY)
29 LOGICAL FOUND1, FOUND2
30 DIMENSION INDEX(SYS), NUMS(1)
31
32 IF(N .GE. SYS) GOTO 5
33 1 WRITE(TTYOUT,100)
34 2 CALL GETNUM(NUMS,1,TTYIN,TTYOUT)
35 I = NUMS(1)
36 IF(I .EQ. 0) RETURN
37
38 MAKE SURE THAT WE DON'T ALREADY HAVE AN ELEMENT # I
39
40 CALL FINDIT(N,I,J,J,INDEX,FOUND1,FOUND2,SY)
41 IF(FOUND1) GOTO 3
42
43 ALL OK, ADD ELEMENT AND PUT A ONE ON MAIN DIAGONAL
44
45 N = N + 1
46 INDEX(N) = I
47
48 ZERO OUT ROW AND COL OF NEW ELEMENT
49
50 DO 4 J=1,N
51 MATRIX(J,N) = .FALSE.
52 MATRIX(N,J) = .FALSE.
53 4 CONTINUE
54
55 MATRIX(N,N) = .TRUE.
56
57 GO GET ANOTHER ELEMENT
58

```

```

59 IF(N .GE. SYS) GOTO 5
60 GOTO 2
61 C
62 C ERROR PRINTOUT
63 C
64 C 3 WRITE(ITYOUT,101) 1
65 GOTO 1
66 C
67 C USER HAS EXCEEDED SYSTEM MATRIX SIZE
68 C
69 C 5 WRITE(ITYOUT,102) SYS
70 RETURN
71 C
72 C FORMATS
73 C
74 100 FORMAT(28H-TYPE ELEMENTS TO BE ADDED ?)
75 101 FORMAT(12H-***ERROR***,15,24H ALREADY IN SYSTEM INDEX)
76 102 FORMAT(62H-***NOTE*** NUMBER OF ELEMENTS HAS REACHED COMPUTER'S LI
77 MIT OF,15/25H "ADD" COMMAND TERMINATED)
78 END

```

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PROGRAM CYCLE

NOTICE

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REPRODUCED, STORED IN A RETRIEVAL SYSTEM, OR TRANSMITTED
IN ANY FORM OR BY ANY MEANS, ELECTRONIC, MECHANICAL,
PHOTOCOPYING, RECORDING, OR OTHERWISE, WITHOUT THE
PRIOR PERMISSION OF THE
UNIVERSITY OF DAYTON RESEARCH INSTITUTE.

PROGRAM CYCLE (SYS04)

PROGRAM THAT INITIALIZES A WEIGHTED MATRIX
AND RESOLVES THE THRESHOLD WITH WITH GEODEDIC OUTPUT

WRITTEN BY: DAVID R. YINGLING, JR.
DECISION SYSTEMS LAB
ENGINEERING AND PUBLIC POLICY GROUP
UNIVERSITY OF DAYTON
DAYTON, OHIO 45469
513-229-2238

IMPLICIT INTEGER (A-Z)
COMMON /INFO/ QTYPE, TTYIN, TTYOUT
COMMON /BLK1/ L1(160), L2(160)

DIMENSION IM(50,50), INDEX(50), NUMS(3), BULL(5704)
LOGICAL QTYPE
DATA CH/2HCH/, AL/2HAL/, DL/2HDL/, FI/2HFI/
DATA PR/2HPR/, RE/2HRE/, TE/2HTE/, HELP/2HHE/
DATA YES/1HY/, NU/1HN/

INITIALIZATIONS

TTYIN = 5
TTYOUT = 2
MWEIGH = 9
QTYPE = .TRUE.
PERMFL = 20

MAIN CONTROL SECTION

FULL TEXT ?

WRITE(TTYOUT,2)
READ(TTYIN,3) ANSWER
IF(ANSWER.EQ. YES) QTYPE = .FALSE.

NEW SYSTEM ?

WRITE(TTYOUT,4)
READ(TTYIN,3) ANSWER

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```

59      IF (ANSWER .EQ. NO) GOTO 10
60 C
61 C      COME HERE WHEN NEW SYSTEM
62 C
63      11 WRITE (TTYOUT,5)
64      CALL GETNUM (NUMS,1)
65      N = NUMS(1)
66      IF (N .LE. 50) GOTO 12
67      WRITE (TTYOUT,7)
68      GOTO 11
69 C
70 C      READ IN FROM PERMFL OLD SYSTEM
71 C
72      10 REWIND 20
73      READ (20) N
74      READ (20) IM
75      READ (20) INDEX
76      GOTO 15
77 C
78 C      ASK IF USER WANTS REGULAR INDEXING
79 C
80      12 CONTINUE
81      WRITE (TTYOUT,8)
82      READ (TTYIN,3) ANSWER
83      IF (ANSWER .EQ. NO) GOTO 13
84 C
85 C      PUT IN REGULAR INDEXES
86 C
87      DO 14 I=1,N
88      INDEX(I) = 1
89      14 CONTINUE
90      GOTO 15
91 C
92 C      READ IRREGULAR INDEXES
93 C
94      13 CONTINUE
95      WRITE (TTYOUT,9)
96      DO 16 I=1,N
97      CALL GETNUM (NUMS,1)
98      INDEX(I) = NUMS(1)
99      IF (I .EQ. 1) GOTO 16
100     M = 1 - 1
101     DO 46 J = 1,M
102     IF (INDEX(J) .EQ. INDEX(I)) WRITE (TTYOUT,103) INDEX(I)
103     IF (INDEX(J) .EQ. INDEX(I)) I = I - 1
104     46 CONTINUE
105     16 CONTINUE
106 C
107     15 REWIND 20
108     WRITE (20) N
109     WRITE (20) IM
110     WRITE (20) INDEX
111 C
112 C      BEGIN FILLING THE ADJACENCY MATRIX
113 C      ASK USER FOR A COMMAND KEYWORD AND CHECK
114 C
115     WRITE (TTYOUT,102)
116     READ (TTYIN,100) CMD

```

```

117
118 IF(CMD .EQ. CH) GOTO 17
119 IF(CMD .EQ. AL) GOTO 18
120 IF(CMD .EQ. FI) GOTO 19
121 IF(CMD .EQ. DL) GOTO 20
122 IF(CMD .EQ. HELP) GOTO 21
123 IF(CMD .EQ. RE) GOTO 22
124 IF(CMD .EQ. TE) GOTO 23
125 IF(CMD .EQ. PR) GOTO 45
126 C
127 WRITE(ITYOUT,101) CMD
128 GOTO 15
129 C
130 C *****
131 C * CHANGE WEIGHT ROUTINE *
132 C *****
133 17 CALL CHANGE(N,IM,INDEX)
134 GOTO 15
135 C *****
136 C * ADD ELEMENT ROUTINE *
137 C *****
138 18 CALL ADD(N,IM,INDEX)
139 GOTO 15
140 C *****
141 C * DELETE ELEMENT ROUTINE *
142 C *****
143 20 CALL DELETE(N,IM,INDEX)
144 GOTO 15
145 C *****
146 C * FILL SYSTEM WITH WEIGHTS *
147 C *****
148 19 CALL FILL(N,IM,INDEX,R1,R2)
149 GOTO 15
150 C *****
151 C * PRINT SYSTEM OUT *
152 C *****
153 45 CALL PRNTMT(N,IM,INDEX,THRESH, FALSE, )
154 GOTO 15
155 C *****
156 C * RESOLVE SYSTEM *
157 C *****
158 22 CALL RESOLV(N,IM,INDEX,MWEIGH)
159 GOTO 15
160 C *****
161 C * TERMINATION *
162 C *****
163 23 CALL EXIT
164 C *****
165 C * HELP !!! *
166 C *****
167 21 WRITE(ITYOUT,104)
168 GOTO 15
169 C
170 C *****
171 C FORMATS
172 C *****
173 C
174 2 FORMAT(1H0,35H FULL TEXT QUERIES DESIRED ? (Y/N) )

```



```

175 3 FORMAT(A1)
176 4 FORMAT(11,20H NEW SYSTEM ? (Y/N) )
177 5 FORMAT(11,31H NUMBER OF ELEMENTS (50 MAX.) ?)
178 7 FORMAT(11,27H ***TOO LARGE, TRY AGAIN***)
179 8 FORMAT(11,46H REGULAR INDEXING OF ELEMENTS DESIRED ? (Y/N) )
180 9 FORMAT(11,28H ENTER INDEXES ONE AT A TIME)
181 100 FORMAT(A2)
182 101 FORMAT(11-,31H ***ERROR*** INVALID COMMAND=>,A2)
183 102 FORMAT(11-,32H TYPE CYCLE COMMAND (OR "HELP") )
184 103 FORMAT(11,34H ELEMENT HAS ALREADY BEEN ENTERED>,12)
185 104 FORMAT(11-,10X,18H***HELP MESSAGE***/1H0,30H AL - ADD AN ELEMENT T
186 +U SYSTEM/1H ,35H DL - DELETE AN ELEMENT FROM SYSTEM/1H ,34H FI - F
187 +ILL SYSTEM (ASSIGN WEIGHTS)/1H ,37H CH - CHANGE WEIGHT OF A RELATI
188 +ONSHIP/1H ,20H RE - RESOLVE SYSTEM/1H ,22H PR - PRINT SYSTEM OUT/1
189 +H ,23H TE - TERMINATE SESSION/1H ,27H HELP - REPRINTS ABOVE LIST)
190 END

```

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```

1 SUBROUTINE PRNTH(N,MAT,INDEX,THRESH,SELPNT)
2 C
3 C THIS SUBROUTINE PRINTS ALL RELATIONSHIPS
4 C   >= TO THE THRESHOLD
5 C
6 C
7 C WRITTEN BY: DAVID R. YINGLING, JR.
8 C             DECISION SYSTEMS LAB
9 C             ENGINEERING AND PUBLIC POLICY GROUP
10 C            UNIVERSITY OF DAYTON
11 C            DAYTON, OHIO 45469
12 C            513-229-2238
13 C
14 IMPLICIT INTEGER (A-Z)
15 COMMON /INFO/ QTYPE, TTYIN, TTYOUT
16 C
17 INTEGER INDEX(50), MAT(50,50), LIST(50,2)
18 LOGICAL QTYPE, SELPNT
19 C
20 IF(SELPNT) WRITE(TTYOUT,1)
21 C
22 DO 10 I=1,N
23   LTR = 0
24   DO 11 J=1,N
25     IF(.NOT. SELPNT) GOTO 13
26     IF(MAT(I,J) .LT. THRESH) GOTO 11
27   13 IF(I .EQ. J) GOTO 11
28     CTR = CTR + 1
29     LIST(CTR,1) = INDEX(J)
30     LIST(CTR,2) = MAT(I,J)
31   11 CONTINUE
32 C
33   IF(CTR .EQ. 0) GOTO 12
34 C
35   WRITE(TTYOUT,2) INDEX(I),(LIST(1,1),LIST(1,2),11 = 1,CTR)
36   GOTO 10
37 C
38 12 WRITE(TTYOUT,3) INDEX(I)
39 C
40 10 CONTINUE
41 RETURN
42 C
43 2 FORMAT(1H ,15,2H=>,8(7(15,1H(,11,2H),)/9X))
44 3 FORMAT(12,2H=>)
45 1 FORMAT(1H-,15X,17H THRESHOLD MATRIX)
46 END

```

```

1      SUBROUTINE QUEST(LL1, LL2)
2 C
3 C      THIS SUBROUTINE PRESENTS THE QUESTIONS
4 C
5 C
6 C      WRITTEN BY:  DAVID R. YINGLING, JR.
7 C                  DECISION SYSTEMS LAB
8 C                  ENGINEERING AND PUBLIC POLICY GROUP
9 C                  UNIVERSITY OF DAYTON
10 C                 DAYTON, OHIO  45469
11 C                 513-229-2238
12 C
13      IMPLICIT INTEGER (A-Z)
14      COMMON /INFJ/  QTYPE, TTYIN, TTYOUT
15      COMMON /BLK1/  L1(160), L2(160)
16 C
17      LOGICAL QTYPE
18 C
19      IF(QTYPE) GOTO 15
20      DEFINE FILE 8(260,160,U,UNUSED)
21      I1 = LL1 + 4
22      I2 = LL2 + 4
23      READ(8'I1) (L1(I),I=1,160)
24      READ(8'I2) (L2(I),I=1,160)
25      I1 = 0
26      I2 = 0
27      DO 12 J=1,10
28      IF(L1(J).EQ. 0) GOTO 12
29      L = L1(J)
30      I1 = I2 + 1
31      I2 = I1 + L - 1
32      WRITE(TTYOUT,2) (L1(K + 10),K = I1,I2)
33 12 CONTINUE
34      WRITE(TTYOUT,3)
35      I1 = 0
36      I2 = 0
37      DO 13 J = 1,10
38      IF(L2(J).EQ. 0) GOTO 13
39      L = L2(J)
40      I1 = I2 + 1
41      I2 = I1 + L - 1
42      WRITE(TTYOUT,2) (L2(K + 10),K = I1,I2)
43 13 CONTINUE
44      WRITE(TTYOUT,1)
45 14 RETURN
46 15 WRITE(TTYOUT,4) LL1, LL2
47      GOTO 14
48 C
49      1 FORMAT(10H WEIGHT ? )
50      2 FORMAT(1X,15A4)
51      3 FORMAT(1H ,20H HAS BEEN RELATED TO)
52      4 FORMAT(1X,15,2H R,14,9H WEIGHT ? )
53      END

```

```

1 SUBROUTINE ADD(N,MAT,INDEX)
2 C
3 C THIS SUBROUTINE ADDS ELEMENTS TO THE ADJACENCY MATRIX
4 C
5 C WRITTEN BY: DAVID R. YINGLING, JR.
6 C DECISION SYSTEMS LAB
7 C ENGINEERING AND PUBLIC POLICY GROUP
8 C UNIVERSITY OF DAYTON
9 C DAYTON, OHIO 45463
10 C 513-229-2238
11 C
12 C IMPLICIT INTEGER (A-Z)
13 C DIMENSION NUMS(3), MAT(50,50), INDEX(50)
14 C
15 C COMMON /BLK1/ L1(160), L2(160)
16 C COMMON /INFO/ QTYPE, TTYIN, TTYOUT
17 C DATA N00/1HN/, N0/2HNO/, Y/1HY/, YES/3HYES/
18 C
19 C LOGICAL FOUND1, FOUND2, QTYPE
20 C
21 C
22 C ASK FOR ADDED ELEMENT NUMBER(3)
23 C
24 2 WRITE(TTYOUT,106)
25 CALL GETNUM(NUMS,1)
26 N = N + 1
27 INDEX(N) = NUMS(1)
28 IF(INDEX(N) .LE. 9999) GOTO 1
29 WRITE(TTYOUT,100) INDEX(N)
30 12 N = N - 1
31 GOTO 2
32 C
33 C CHECK FOR A ZERO INPUT
34 C
35 1 IF(INDEX(N) .GT. 0) GOTO 3
36 N = N - 1
37 RETURN
38 C
39 C CHECK FOR DUPLICATE ELEMENT
40 C
41 3 K = N - 1
42 CALL FINDET(K,INDEX(N),J,J,J,INDEX,FOUND1,FOUND2)
43 IF(.NOT. FOUND1) GOTO 11
44 WRITE(TTYOUT,101) INDEX(N)
45 GOTO 12
46 11 CONTINUE
47 C
48 C SEE IF USER WANTS TO BURDER
49 C
50 WRITE(TTYOUT,102)
51 READ(TTYIN,103) ANSWER
52 IF(ANSWER .EQ. Y) GOTO 4
53 IF(ANSWER .EQ. YES) GOTO 4
54 IF(ANSWER .EQ. N00) GOTO 2
55 IF(ANSWER .EQ. NO) GOTO 2
56 WRITE(TTYOUT,104) ANSWER
57 GOTO 3
58 C

```

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```

59 C      BORDER INPUT
60 C
61      4 DO 5 I=1,N
62      6 CALL QUEST(INDEX(I), INDEX(N))
63      CALL GETNUM(NUMS,I)
64      WEIGHT = NUMS(I)
65      IF(WEIGHT .LE. 9) GOTO 7
66      WRITE(TTYOUT,105) WEIGHT
67      GOTO 6
68      7 MAT(I,N) = WEIGHT
69      5 CONTINUE
70 C
71      DO 9 I=1,N
72      9 CALL QUEST(INDEX(N), INDEX(I))
73      CALL GETNUM(NUMS,I)
74      WEIGHT = NUMS(I)
75      IF(WEIGHT .LE. 9) GOTO 10
76      WRITE(TTYOUT,105) WEIGHT
77      GOTO 9
78      10 MAT(N,I) = WEIGHT
79      8 CONTINUE
80      GOTO 2
81 C
82 C      *****
83 C      FORMATS
84 C      *****
85 C
86      100 FORMAT(1H0,22H ***ELEMENT TOO LARGE>,15,3H***)
87      101 FORMAT(1H0,30H ***ELEMENT ALREADY IN SYSTEM>,15,3H***)
88      102 FORMAT(1H ,31H BORDER ON THIS ELEMENT ? (Y/N))
89      103 FORMAT(A2)
90      104 FORMAT(1H ,21H ***INVALID RESPONSE>,A2,3H***)
91      105 FORMAT(1H ,28H *** WEIGHT VALUE TOO LARGE>,15,3H***)
92      106 FORMAT(1H ,29H ELEMENT NUMBER TO BE ADDED ?)
93      END

```

```

1 SUBROUTINE CHANGE(N,MAT,INDEX)
2 C
3 C THIS SUBROUTINE WILL ALLOW THE USER TO CHANGE
4 C A WEIGHT IN THE ADJACENCY MATRIX
5 C
6 C
7 C WRITTEN BY: DAVID R. YINGLING, JR.
8 C DECISION SYSTEMS LAB
9 C ENGINEERING AND PUBLIC POLICY GROUP
10 C UNIVERSITY OF DAYTON
11 C DAYTON, OHIO 45469
12 C 513-229-2230
13 C
14 C IMPLICIT INTEGER (A-Z)
15 C DIMENSION MAT(50,50), INDEX(50), NUMS(3)
16 C COMMON /INFO/ QTYPE, TTYIN, TTYOUT
17 C
18 C LOGICAL FOUND1, FOUND2
19 C
20 C
21 C 5 WRITE(TTYOUT,100)
22 C CALL GETNUM(NUMS,3)
23 C GOTO 6
24 C 1 CALL GETNUM(NUMS,3)
25 C
26 C CHECK FOR ZERO INPUT
27 C
28 C 6 IF(NUMS(1) .GT. 0 .AND. NUMS(2) .GT. 0) GOTO 8
29 C RETURN
30 C
31 C CHECK FOR EXISTANCE OF ELEMENTS
32 C
33 C 8 CALL FINDIT(N,NUMS(1),NUMS(2),X,Y,INDEX,FOUND1,FOUND2)
34 C
35 C DID WE GET A VALID INPUT ?
36 C
37 C IF(FOUND1 .AND. FOUND2) GOTO 3
38 C IF(Y .EQ. 0) GOTO 4
39 C WRITE(TTYOUT,101) NUMS(1)
40 C IF(Y .GT. 0) GOTO 5
41 C 4 WRITE(TTYOUT,101) NUMS(2)
42 C GOTO 5
43 C
44 C 3 IF(NUMS(3) .LE. 9) GOTO 7
45 C WRITE(TTYOUT,102) NUMS(3)
46 C GOTO 5
47 C
48 C 7 MAT(X,Y) = NUMS(3)
49 C GOTO 1
50 C
51 C *****
52 C FORMATS
53 C *****
54 C
55 C 100 FORMAT(1H ,41H ENTER A REACHES TO R, WEIGHT (3 NUMBERS))
56 C 101 FORMAT(1H ,20H ***ELEMENT DOES NOT EXIST>,15,3H*** )
57 C 102 FORMAT(1H ,20H ***WEIGHT VALUE TOO LARGE>,15,3H*** )
58 C END

```

```

1      SUBROUTINE DELETE(IN,MAT,INDEX)
2
3      IMPLICIT INTEGER (A-Z)
4      DIMENSION MAT(50,50), INDEX(50), NUMS(3)
5      COMMON /INDEX/ JTYPE, TTYIN, TTYOUT
6
7
8      LOGICAL FOUND1, FOUND2
9
10
11      WRITE(TTYOUT,101)
12      1 CALL GETNUM(NUMS,1)
13
14      IS IT ZERO ?
15
16      IF (NUMS(1) .EQ. 0) RETURN
17
18      ELEMENT IN SET ?
19
20      CALL FINDIT(IN,NUMS(1),J,I,J,INDEX,FOUND1,FOUND2)
21      IF (FOUND1) GOTO 3
22      WRITE(TTYOUT,100) NUMS(1)
23      GOTO 1
24
25      3 IF (I .EQ. N) GOTO 4
26
27      N3 = N - 1
28      DO 5 K1=1,N3
29      K2 = K1 + 1
30      DO 6 J=1,N
31      MAT(K1,J) = MAT(K2,J)
32      6 CONTINUE
33      DO 7 J=1,N
34      MAT(J,K1) = MAT(J,K2)
35      7 CONTINUE
36      5 CONTINUE
37      DO 8 K1=1,N3
38      K2 = K1 + 1
39      INDEX(K1) = INDEX(K2)
40      8 CONTINUE
41
42      4 N = N - 1
43      GOTO 1
44
45      *****
46      FORMATS
47      *****
48
49      100 FORMAT(1H ,27H ***ELEMENT DOES NOT EXIST**,15,4H*** )
50      101 FORMAT(1H ,30H ELEMENT NUMBER TO BE ERASED ?)
51      END

```

```

1 SUBROUTINE FILLIN,MAT,INDEX,R1,R2)
2
3 THIS SUBROUTINE ALLOWS THE USER TO FILL UP THE ADJACENCY MATRIX
4
5
6 *WRITTEN BY: DAVID R. YINGLING, JR.
7 DECISION SYSTEMS LAB
8 ENGINEERING AND PUBLIC POLICY GROUP
9 UNIVERSITY OF DAYTON
10 DAYTON, OHIO 45469
11 513-229-2238
12
13 IMPLICIT INTEGER (A-Z)
14 DIMENSION MAT(50,50), INDEX(50), NUMS(3)
15 COMMON /INFO/ QTYPE, TTYIN, TTYOUT
16 COMMON /BLK1/ L1(160), L2(160)
17 LOGICAL QTYPE
18 SEE IF THIS IS A RESTART
19
20 IF(R1 .GT. 0 .OR. R2 .GT. 0) GOTO 1
21 ROW = 1
22 COL = 1
23 R1 = 0
24 R2 = 0
25
26 10 DO 2 I=ROW,N
27 DO 2 J=COL,N
28 IF(I .EQ. J) GOTO 2
29 5 CALL QUEST(INDEX(I),INDEX(J))
30 CALL GETNUM(NUMS,1)
31 IF(NUMS(1) .EQ. 10) GOTO 3
32 IF(NUMS(1) .LE. 9) GOTO 4
33 WRITE(TTYOUT,101) NUMS(1)
34 GOTO 5
35 4 MAT(I,J) = NUMS(1)
36 2 CONTINUE
37 GOTO 6
38
39 RESTART
40
41 1 I = R1
42 DO 7 J=R2,N
43 8 CALL QUEST(INDEX(I),INDEX(J))
44 CALL GETNUM(NUMS,1)
45 IF(NUMS(1) .EQ. 10) GOTO 3
46 IF(NUMS(1) .LE. 9) GOTO 9
47 WRITE(TTYOUT,101) NUMS(1)
48 GOTO 8
49 9 MAT(I,J) = NUMS(1)
50 7 CONTINUE
51 ROW = 1
52 COL = R1 + 1
53 R1 = 0
54 R2 = 0
55 GOTO 10
56
57 GENERATE RESTART PARAMS
58

```

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```

52      3 H1 - I
63      "  P2 - J
61      6 RETURN
62 C
63 C      *****
64 C      FORMATS
65 C      *****
66 C
67      100 FORMAT(10,15,20H,15,10H,WEIGHT ?)
68      101 FORMAT(11,20H ***WEIGHT VALUE TOO LARGE>,15,10H)
69      END

```

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```

1  SUBROUTINE FINDIT(N1,N2,S1,S2,INDEX,FOUND1,FOUND2)
2  C
3  C  THIS SUBROUTINE FINDS ELEMENTS N1, N2 IN THE INDEX SET
4  C
5  C  THE VALUES S1, S2 ARE THE POSITIONS OF N1 AND N2 IN THE
6  C  INDEX SET.
7  C
8  C  FOUND1 AND FOUND2 ARE LOGICAL VALUES AND ARE SET EQUAL
9  C  TO .TRUE. IF N1 OR N2 (RESPECTIVELY) ARE FOUND IN THE
10 C  INDEX SET.
11 C
12 C  WRITTEN BY:  DAVID R. YINGLING, JR.
13 C              DECISION SYSTEMS LAB
14 C              ENGINEERING AND PUBLIC POLICY GROUP
15 C              UNIVERSITY OF DAYTON
16 C              DAYTON, OHIO  45469
17 C              513-229-2238
18 C
19 C  IMPLICIT INTEGER (A-Z)
20 C  DIMENSION INDEX(128)
21 C  LOGICAL FOUND1, FOUND2
22 C
23 C  FOUND1 = .FALSE.
24 C  FOUND2 = .FALSE.
25 C  S1      = 0
26 C  S2      = 0
27 C
28 C  DO 1 I=1,N
29 C    IF(N1 .EQ. INDEX(I)) S1 = I
30 C    IF(N2 .EQ. INDEX(I)) S2 = I
31 C  1 CONTINUE
32 C
33 C  IF(S1 .GT. 0) FOUND1 = .TRUE.
34 C  IF(S2 .GT. 0) FOUND2 = .TRUE.
35 C  RETURN
36 C  END

```

```

1 SUBROUTINE GETNUM(ARRAY,N)
2 C
3 THIS SUBROUTINE WILL READ "N" UNSIGNED INTEGERS FROM
4 THE TERMINAL TYPED IN A FREE FORMAT AND STORE THEM IN
5 "ARRAY"
6 C
7 WRITTEN BY: DAVID R. YINGLING, JR.
8 DECISION SYSTEMS LAB
9 ENGINEERING AND PUBLIC POLICY GROUP
10 UNIVERSITY OF DAYTON
11 DAYTON, OHIO 45469
12 513-229-2238
13 C
14 IMPLICIT INTEGER (A-Z)
15 COMMON /INFO/ QTYPE, TTYIN, TTYOUT
16 DIMENSION ARRAY(1), BUFFER(80), NUMS(10)
17 DATA NUMS/1H0,1H1,1H2,1H3,1H4,1H5,1H6,1H7,1H8,1H9/
18 DATA BLANK/1H /, COMMA/1H,/
19 C
20 DO 8 I=1,N
21 ARRAY(I) = 0
22 8 CONTINUE
23 C
24 1 READ(TTYIN,200) BUFFER
25 C
26 L = N
27 POWER = 0
28 C
29 DO 2 I=1,80
30 K = 81 - I
31 IF(BUFFER(K) .EQ. BLANK .OR. BUFFER(K) .EQ. COMMA) GOTO 3
32 C
33 FOUND A CHARACTER, SEE IF IT'S A VALID NUMERIC
34 C
35 DO 4 J=1,10
36 II = J - 1
37 IF(BUFFER(K) .NE. NUMS(J)) GOTO 4
38 C
39 ITS A NUMBER, ADD IT TO PRESENT SUM
40 C
41 ARRAY(L) = ARRAY(L) + (II * (10**POWER))
42 POWER = POWER + 1
43 GOTO 2
44 4 CONTINUE
45 C
46 COME HERE IF CHARACTER FOUND WAS NOT NUMERIC
47 C
48 WRITE(TTYOUT,100)
49 GOTO 1
50 C
51 FOUND A DELIMITER, SEE IF END OF A NUMBER
52 C
53 3 IF(ARRAY(L) .EQ. 0) GOTO 2
54 L = L - 1
55 IF(L .EQ. 0) GOTO 5
56 ARRAY(L) = 0
57 POWER = 0
58 2 CONTINUE

```

```

59 C
60 C MAKE SURE NUMBER(S) IS/ARE LESS THAN 99999 SIN
61 C WE DON'T EXCEED 15 FORMATS
62 C
63 C 5 DO 6 I=1,N
64 C IF (ARRAY(I) .GT. 99999) GOTO 7
65 C 6 CONTINUE
66 C RETURN
67 C
68 C ERROR MESSAGE
69 C
70 C 7 WRITE (TTYOUT,101)
71 C GOTO 1
72 C
73 C FORMATS
74 C
75 C 100 FORMAT(37H-***ERROR*** INPUT NOT NUMERIC--RETRY)
76 C 101 FORMAT(37H-***ERROR*** NUMBER(S) TOO LARGE--RETRY)
77 C 200 FORMAT(80A1)
78 C END

```

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```

1 SUBROUTINE RESOLVE(N,IN,LIST,THRESH)
2 C
3 C THIS SUBROUTINE RESOLVES THE WEIGHTED MATRIX
4 C
5 C COMMON /INFO/ QTYPE, TTYIN, TTYOUT
6 C
7 C IMPLICIT INTEGER (A-Z)
8 C INTEGER IN(50,50), LIST(50)
9 C LOGICAL ADJ(50,50), QTYPE, RM(50,50)
10 C
11 C T = THRESH
12 C
13 C CONSTRUCT BINARY ADJACENCY MATRIX
14 C
15 DO 10 I=1,N
16 DO 10 J=1,N
17 ADJ(I,J) = .FALSE.
18 IF(I .EQ. J) GOTO 11
19 IF(IN(I,J) .LT. T) GOTO 10
20 11 ADJ(I,J) = .TRUE.
21 10 CONTINUE
22 C
23 C CHECK IF CYCLE IS RESOLVED AT THIS THRESHOLD
24 C
25 C CALL TRNCLS(ADJ,RM,N)
26 C
27 C CHECK TO SEE IF REACHABILITY MAT IS
28 C ALL ONES, IF SO, CYCLE IS RESOLVED
29 C IF NOT, T = T - 1 AND CONSTRUCT NEW ADJ MA RIX
30 C
31 DO 12 I=1,N
32 DO 12 J=1,N
33 IF(.NOT. RM(I,J)) GOTO 13
34 12 CONTINUE
35 C
36 C TELL USER CYCLE IS RESOLVED
37 C
38 C WRITE(TTYOUT,1) T
39 C
40 C PRINT UNIVERSAL MATRIX
41 C
42 C CALL PRNTMT(N,IN,LIST,T,.TRUE.)
43 C
44 C PRINT GEODEDIC OUTPUT
45 C
46 C CALL GEOD(ADJ,N,LIST)
47 C RETURN
48 C
49 C NO GOOD, TRY AGAIN
50 C
51 13 T = T - 1
52 GOTO 15
53 C
54 1 FORMAT(1H-, ' ***CYCLE RESOLVED***', //, ' THRESHOLD=>', 2X, I2)
55 END

```

```

1 SUBROUTINE TRNCLS(A,B,N)
2 C
3 C THIS SUBROUTINE TAKES THE TRANSITIVE CLOSURE
4 C OF MAT A AND PUTS THE ANSWER INTO MAT B.
5 C
6 IMPLICIT INTEGER (A-Z)
7 INTEGER C(50)
8 LOGICAL A(50,50), B(50,50)
9 C
10 DO 1 I=1,N
11 DO 1 J=1,N
12 1 B(I,J) = A(I,J)
13 C
14 I = 0
15 2 N1 = 0
16 NT = 0
17 I = I + 1
18 IF(I .GT. N) RETURN
19 5 DO 3 K=1,N
20 IF(.NOT. B(I,K)) GOTO 3
21 NT = NT + 1
22 C(NT) = K.
23 3 CONTINUE
24 C
25 C IF(NT .EQ. N1) GOTO 2
27 C
28 N1 = NT
29 DO 4 L=1,N1
30 K = C(L)
31 DO 4 J=1,N
32 4 IF(B(K,J)) B(I,J) = .TRUE.
33 C
34 NT = 0
35 GOTO 5
36 END

```

```

1      SUBROUTINE GEODIA,N,INDEX)
2
3      THIS SUBROUTINE PRINTS OUT GEODEDIC CYCLES
4
5      COMMON /INFO/ QTYPE, TTYIN, TTYOUT
6
7      IMPLICIT INTEGER (A-Z)
8      INTEGER B(50,50), C(50,50), G(50,50), LIM(50)
9      INTEGER INDEX(50), PRINT(50)
10     LOGICAL QTYPE, A(50,50)
11
12     DATA G/2500*0/, R/2500*0/
13
14     FORM A - 1
15
16
17     DO 10 I=1,N
18     DO 10 J=1,N
19     IF(.NOT. A(J,I)) GOTO 10
20     R(J,I) = 1
21     G(J,I) = 1
22     IF(J.EQ. 1) G(J,I) = 0
23 10 CONTINUE
24
25     FORM G
26
27     NBN = 1
28 11 DO 12 I=1,N
29     DO 12 J=1,N
30     C(J,I) = 0
31     DO 13 K=1,N
32     MA = 0
33     IF(A(J,K)) MA = 1
34 13 C(J,I) = MA * B(K,I) + C(J,I)
35     IF(C(J,I).NE. 0) C(J,I) = 1
36 12 CONTINUE
37
38     COMPUTE TO THE NTH
39
40     NBN = NBN + 1
41     DO 14 I=1,N
42     DO 14 J=1,N
43     G(I,J) = G(I,J) + NBN * (C(I,J) - B(I,J))
44 14 R(I,J) = C(I,J)
45     IF(NBN.LT. N-1) GOTO 11
46
47     PRINT HEADING
48
49     WRITE(TTYOUT,1)
50
51     COMPUTE PATHS
52
53     NCPT = 0
54     DO 16 IGP=2,N
55     LIM = IGP - 1
56     DO 16 JGP=1,LIM
57     NB = G(IGP,JGP)
58     NA = G(JGP,IGP)

```

```

5) IF (NA .EQ. 0 .OR. NB .EQ. 0) GOTO 16
6) NBND = 1
61 LION(NBND) = IGP
62 IF (NB .GT. 1) GOTO 17
63 C
64 NBND = NBND + 1
65 LION(NBND) = JGP
66 GOTO 21
67 C
68 19 WRITE (TTYOUT,2) NB,IGP,JGP
69 GOTO 21
70 17 LMM = NB - 1
71 LAST = 0
72 DO 20 IN=1,LMM
73 NBND = NBND + 1
74 LION(NBND) = NOTR(G,IGP,NB-IN,JGP,IN,IX,N,LAST)
75 LAST = LION(NBND)
76 20 IF (IX .EQ. 1) GOTO 19
77 C
78 NBND = NBND + 1
79 LION(NBND) = JGP
80 21 IF (NA .GT. 1) GOTO 22
81 LION(NBND+1) = IGP
82 GOTO 23
83 C
84 24 WRITE (TTYOUT,2) NA,IGP,JGP
85 GOTO 23
86 C
87 2' LMM = NA - 1
88 LAST = 0
89 DO 25 IN=1,LMM
90 NBND = NBND + 1
91 LION(NBND) = NOTR(G,JGP,NA-IN,IGP,IN,IX,N,LAST)
92 LAST = LION(NBND)
93 25 IF (IX .EQ. 1) GOTO 24
94 C
95 LION(NBND+1) = IGP
96 23 CONTINUE
97 NCPT = NCPT + 1
98 LMC = NBND + 1
99 C
100 C FIX PRINT OUT
101 C
102 DO 15 NX = 1,LMC
103 NUM = LION(NX)
104 15 PRINT(NX) = INDEX(NUM)
105 NXXX = NA + NB
106 WRITE (TTYOUT,3) NCPT,NXXX,INDEX(IGP),INDEX(JGP),(PRINT(I),I=1,LMC)
107 16 CONTINUE
108 RETURN
109 C
110 C FORMATS
111 C
112 1 FORMAT(' ',NUMBER',5X,'#LINKS',4X,'ELEMENTS',5X,'PATH')
113 2 FORMAT(1X,' PAS TRIUVE ',3I5)
114 3 FORMAT(' ',1X,14,8X,12,3X,15,' ',',',15,3X,6(7(I5),/33X))
115 END

```



```

1  INTEGER FUNCTION NOTRIG,ICL,NBLG,ILG,NBCL,IX,NHND,LAST)
2  C
3  IMPLICIT INTEGER (A-Z)
4  INTEGER G(50,50), LST(50)
5  C
6  IX = 0
7  IBI = 0
8  DO 1 I=1,NHND
9  IF(G(I,ILG) .NE. NBLG) GOTO 1
10 IBI = IBI + 1
11 LST(IBI) = I
12 1 CONTINUE
13 IF(IBI .LE. 0) RETURN
14 DO 2 I=1,NHND
15 IF(G(ICL,I) .NE. NBCL) GOTO 2
16 IF(LAST .EQ. 0) GOTO 3
17 IF(G(LAST,I) .NE. 1) GOTO 2
18 5 DO 3 K=1,IBI
19 3 IF(I .EQ. LST(K)) GOTO 4
20 2 CONTINUE
21 IX = 1
22 RETURN
23 4 NOTR = LST(K)
24 RETURN
25 END

```

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```

1 PROGRAM MAKEIT
2 C
3 C *****
4 C *
5 C *
6 C * ALL RIGHTS RESERVED. NO PART OF THIS PROGRAM MAY BE SOLD,
7 C * REPRODUCED, STORED IN A RETRIEVAL SYSTEM, OR TRANSMITTED
8 C * IN ANY FORM OR BY ANY MEANS, ELECTRONIC, MECHANICAL,
9 C * PHOTOCOPYING, RECORDING, OR OTHERWISE, WITHOUT THE
10 C * PRIOR PERMISSION OF THE
11 C * UNIVERSITY OF DAYTON RESEARCH INSTITUTE.
12 C *
13 C *****
14 C
15 C
16 C PROGRAM AUTHOR: DAVID R. YINGLING, JR.
17 C DATE: DECEMBER 31, 1976
18 C
19 C THIS PROGRAM CONVERTS A SEQUENTIAL ACCESSED EDITOR
20 C CREATED FILE INTO A DIRECT ACCESS FILE FOR
21 C FULL TEXT QUERIES
22 C
23 C IMPLICIT INTEGER (A-Z)
24 C
25 C THE ENGLISH TEXT CAN BE 60 CHARS. ON ONE LINE
26 C AND UP TO TEN LINES
27 C
28 C COMMON TTYIN, TTYOUT, N, TOTAL, POS, K, LENGTH(10), PACKED(200)
29 C DATA YUP/INY/
30 C
31 C SYSTEM DEPENDENT INITIALIZATIONS
32 C
33 C THE DIMENSION OF ARRAY "PACKED" MUST BE
34 C CHANGED TO "NWORDS"
35 C
36 C EACH EBCDIC OR BCD LETTER USES X BITS
37 C
38 C FOR IBM 360/370 X=8
39 C FOR CDC 6600/6400 X=6
40 C
41 C NWORDS IS EQU TO 60 CHARACTERS (ONE LINE) DIVIDED BY
42 C NUMBER OF CHARACTERS ABLE TO BE STORED IN ONE WORD (NCHAR) TIMES
43 C TEN (FOR TEN LINES).
44 C
45 C NCHAR = NUMBER OF BITS PER MACHINE WORD (NBITS) DIVIDED
46 C BY "X".
47 C
48 C DIMENSION CARD(60)
49 C
50 C TTYIN = 1
51 C TTYOUT = 2
52 C NBITS = 32
53 C X = 8
54 C NCHAR = NBITS / X
55 C NWORDS = (60 / NCHAR) * 10
56 C TOTAL = NWORDS + 10
57 C
58 C DEFINE DIRECT ACCESS FILE

```

```

59 C   CONTRA UNIT # = 8
60 C   RECSIZE (FOR 32 BITS) = 160
61 C   NUM OF RECORDS = 260
62 C
63 C
64 C   DEFINE FILE B(260,160,0,UNUSED)
65 C
66 C   CONSTRUCT FILE
67 C
68 C   CALL DUT(CHAR,NUMRDS,CARD)
69 C
70 C   N IS THE NUMBER OF TEXT RECORDS AND FIRST RECORD OF FILE
71 C
72 C   WRITE(B'1') N
73 C
74 C   SEE IF USER WANTS TO DISPLAY ELEMENTS
75 C
76 C   WRITE(TTYOUT,1)
77 C   READ (TTYIN,2)  REPLY
78 C   IF(REPLY.EQ. YUP) CALL SHOW
79 C   CALL EXIT
80 C
81 C   SATISFY THE COMPILER WITH STOP STATEMENT
82 C
83 C   STOP11111
84 C
85 C   FORMATS
86 C
87 C   1 FORMAT(1H0,47HPERMFILE HAS BEEN CREATED FOR FULL TEXT QUERIES/1H ,
88 C   *11MSHOW(Y/N) ?)
89 C   2 FORMAT(A1)
90 C   END

```

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```

1 SUBROUTINE DUIT(NCHAR,NWORDS,CARD)
2 C
3 C THIS SUBROUTINE CREATES THE RANDOM TEXT FILE
4 C
5 IMPLICIT INTEGER (A-Z)
6 COMMON TTYIN,TTYOUT,N,TOTAL,POS,START,LENGTH(10),PACKED(200)
7 LOGICAL FLAG
8 DATA R/1HR/, E/1HE/, L/1HL/, SLANT/1H//, ONE/1H1/, TWO/1H2/
9 DATA THREE/1H3/
10 DIMENSION CARD(60)
11 C
12 START = 1
13 POS = 0
14 FLAG = .FALSE.
15 N = 0
16 C
17 C ZAP OUT LENGTH INDICATORS
18 C
19 DO I=1,10
20 LENGTH(I) = 0
21 1 CONTINUE
22 C
23 C FORTRAN UNIT 10 POINTS TO SEQUENTIAL TEXT FILE
24 C
25 2 READ(10,10,END=3) CARD
26 IF(CARD(1).EQ.SLANT) GOTO 4
27 C
28 C NOT A CONTROL KEYWORD, PACK MORE TEXT INTO "PACKED"
29 C
30 CALL PACK(NWORDS,NCHAR,CARD)
31 GOTO 2
32 C
33 C SEE IF FLAG IS ON, IF SO, WRITE CURRENT "PACKED" AND PROCESS
34 C SLANT RECORD
35 C
36 4 IF(.NOT.FLAG) GOTO 5
37 C
38 C WRITE RECORD TO FILE ACCORDING TO LOGPOS (LOGICAL POSITION)
39 C
40 CALL WRITE(LOGPOS,NWORDS)
41 C
42 C PROCESS SLANT RECORD
43 C
44 5 LOGPOS = 0
45 IF(CARD(2).NE.R) GOTO 6
46 C
47 C
48 C
49 C
50 C
51 C
52 C
53 C
54 C THIS IS A RELATIONAL CLAUSE, WHICH ONE ?
55 C
56 IF(CARD(3).EQ.ONE) LOGPOS = 2
57 IF(CARD(3).EQ.TWO) LOGPOS = 3
58 IF(CARD(3).EQ.THREE) LOGPOS = 4

```

```

59      IF (LOGPIS .LT. 2) GOTO 7
60      FLAG = .TRUE.
61      GOTO 2
62 C
63 C      COULD BE A "VEL" CARD
64 C
65      6 IF (CARD(2) .EQ. E) GOTO 8
66      IF (CARD(3) .EQ. L) LOGPIS = N + 5
67      IF (LOGPIS .LT. 2) GOTO 9
68      FLAG = .TRUE.
69      GOTO 2
70 C
71 C      IT SHOULD BE AN EOF "/" CARD
72 C
73      8 IF (CARD(2) .EQ. SLANT) RETURN
74 C
75 C      ITS NOT ANYTHING RECOGNIZABLE
76 C
77      GOTO 9
78      3 WRITE (TTYOUT,11)
79      9 WRITE (TTYOUT,12)
80      CALL EXIT
81 C
82 C      ISSUE STOP STATEMENT TO SATISFY COMPILER
83 C
84      STOP 22222
85 C
86 C      FORMATS
87 C
88      10 FORMAT(6,A1)
89      11 FORMAT(1H0,43H***MISSING "/" CARD AT END OF TEXT FILE***)
90      12 FORMAT(1H0,45H***ERROR*** INVALID SLANT KEYWORD ENCOUNTERED/11,28
91      +HRE-EDIT TEXT FILE TO CORRECT/1H-,23H***MAKE IT TERMINATED***)
92      END

```

370

```

1 SUBROUTINE WRITE(LOGPOS,NWORDS)
2 C
3 C THIS SUBROUTINE WRITES OUT "PACKED" AND "LENGTH" UNTO
4 C THE DIRECT ACCESS FILE
5 C
6 C IMPLICIT INTEGER (A-Z)
7 COMMON ITYIN,ITYOUT,N,TOTAL,POS,START,LENGTH(10),PACKED(200)
8 IF(LOGPOS.LT. 5) GOTO 1
9 C
10 C THIS IS AN ELEMENT'S TEXT
11 C
12 N = N + 1
13 C
14 C WRITE UNTO FILE
15 C
16 1 WRITE(8,LOGPOS) (LENGTH(I), I=1,TOTAL)
17 2 START = 1
18 2 POS = 0
19 C
20 C ZAP OUT LENGTH INDICATOR
21 C
22 DO 2 I=1,10
23 LENGTH(I) = 0
24 2 CONTINUE
25 RETURN
26 END

```

375

```

1 SUBROUTINE PACK(WORDS,NCHAR,CARD)
2
3 THIS SUBROUTINE PACKS:
4
5 1) AS MANY CHARACTERS AS POSSIBLE INTO ONE MACHINE WORD
6
7 2) THE LINES OF TEXT ON A WORD BOUNDARY INTO "PACKED"
8
9
10 IMPLICIT INTEGER (A-Z)
11 COMMON ITRYIN,ITRYOUT,N,TOTAL,POS,START,LENGTH(10),PACKED(200)
12 DIMENSION CARD(60)
13 DATA BLANK/1H /
14
15 C FIGURE OUT LENGTH
16
17 POS = POS + 1
18 DO 1 I=1,60
19 IF(CARD(I) - 1) .NE. BLANK) GOTO 2
20 1 CONTINUE
21
22 I = 59
23 2 LEN = 61 - I
24
25 C NOW COMES THE TRICKY PART. RE-READ INFORMATION TO PACK.
26 C WITH CDC EXTENDED FORTRAN USE AN ENCODE STATEMENT
27 C ON OTHER COMPUTERS USE CORE TO CORE I/O OR A SCRATCH
28 C I/O UNIT. FORTRAN UNIT 20 FOR MY SPECTRA 70 IS VIRTUAL MEMORY
29
30 CORE = 20
31 REWIND CORE
32 WRITE(CORE,3) (CARD(I),I=1,LEN)
33 REWIND CORE
34
35 C FIGURE OUT LENGTH OF PACKED LINE
36
37 LENGTH(POS) = (LEN - 1) / NCHAR + 1
38 END = START + LENGTH(POS)
39 READ(CORE,4) (PACKED(I),I=START,END)
40 START = START + LENGTH(POS)
41 RETURN
42
43 C BE SURE TO CHECK FORMAT N 4
44
45 3 FORMAT(60A1)
46 4 FORMAT(15A4)
47 END

```

```

1 SUBROUTINE SHOW
2 C
3 C THIS SUBROUTINE SHOWS THE ELEMENTS AS THEY MIGHT
4 C APPEAR DURING AN ISM SESSION. THIS IS HELPFUL
5 C FOR DETERMINING THE READABILITY OF THE TEXT
6 C
7 IMPLICIT INTEGER (A-Z)
8 COMMON ITYIN, ITYOUT, N, TOTAL, PDS, START, I1(10), DUM(200)
9 DATA YUP/144/
10 COMMON /SHOWB/ R1(160), L1(160), R2(160), L2(160), R3(160)
11 C
12 C COMMON "SHOWB" AND VECTOR "BLOCK" ARE EQU AND WILL
13 C CONTAIN THE TEXT
14 C
15 LOGICAL ALLS
16 DIMENSION BLOCK(850), NUMS(2)
17 EQUIVALENCE (BLOCK,R1), (EL1,NUMS(1)), (EL2,NUMS(2))
18 FIND (8'2)
19 ALLS = .FALSE.
20 EL1 = 0
21 EL2 = 0
22 C
23 C READ IN RELATIONAL CLAUSES 1,2, + 3
24 C
25 READ(8'2) (R1(I),I=1,TOTAL)
26 READ(8'3) (R2(I),I=1,TOTAL)
27 READ(8'4) (R3(I),I=1,TOTAL)
28 C
29 C SEE IF USER WANTS TO DISPLAY ALL ELEMENTS
30 C
31 WRITE(ITYOUT,9)
32 READ (ITYIN,10) REPLY
33 IF(REPLY .EQ. YUP) GOTO 2
34 C
35 C SEE WHICH ELEMENTS THE USER WANTS TO SHOW
36 C
37 1 WRITE(ITYOUT,5)
38 CALL GETNUM(NUMS,2)
39 IF(EL1 .EQ. 0 .OR. EL2 .EQ. 0) RETURN
40 IF(EL1 .LE. N .AND. EL2 .LE. N) GOTO 8
41 IF(EL1 .GT. N) WRITE(ITYOUT,11) EL1
42 IF(EL2 .GT. N) WRITE(ITYOUT,11) EL2
43 GOTO 1
44 8 I1 = EL1 + 4
45 FIND (R1(I1))
46 I2 = EL2 + 4
47 READ(8'11) (L1(I),I=1,TOTAL)
48 READ(8'12) (L2(I),I=1,TOTAL)
49 OFFSET = 0
50 DO 3 I=1,5
51 I1 = 0
52 I2 = 0
53 DO 4 J=1,10
54 IF(BLOCK(J + OFFSET) .EQ. 0) GOTO 4
55 LENGTH = BLOCK(J + OFFSET)
56 I1 = I1 + 1
57 I2 = I1 + LENGTH - 1
58 WRITE(ITYOUT,7) (BLOCK(C + OFFSET + 10),C=I1,I2)

```



```

59 4 CONTINUE
60   OFFSET = OFFSET + TOTAL
61 3 CONTINUE
62   IF (.NOT. ALLS) GOTO 1
63
64 2 EL1 = EL2 + 1
65   EL2 = EL1 + 1
66   IF (EL1 .GT. N) RETURN
67   IF (EL2 .GT. N) EL2 = 1
68   ALLS = .TRUE.
69   GOTO 8
70 C
71 C
72 C   FORMATS
73 C
74 5 FORMAT(1H ,25HSHOW WHICH TWO ELEMENTS ?)
75 6 FORMAT(2I5)
76 7 FORMAT(1H ,15A4)
77 9 FORMAT(1H-,24HSHOW ALL ELEMENTS(Y/N) ?)
78 10 FORMAT(A1)
79 11 FORMAT(1H-,11H***ERROR***,15,24H NOT FOUND ON QUERY FILE)
80   END

```

375

```

1 SUBROUTINE GLTNUM(ARRAY,N)
2 C
3 THIS SUBROUTINE WILL READ "N" UNSIGNED INTEGERS FROM
4 THE TERMINAL TYPED IN A FREE FORMAT AND STORE THEM IN
5 "ARRAY"
6 C
7 C WRITTEN BY: DAVID R. YINGLING, JR.
8 C DECISION SYSTEMS LAB
9 C ENGINEERING AND PUBLIC POLICY GROUP
10 C UNIVERSITY OF DAYTON
11 C DAYTON, OHIO 45469
12 C 513-229-2238
13 C
14 IMPLICIT INTEGER (A-Z)
15 COMMON TTYIN, TTYOUT, N, TOTAL, PDS, START, LENG(10), DUM(200)
16 DIMENSION ARRAY(1), BUFFER(80), NUMS(10)
17 DATA NUMS/1H0,1H1,1H2,1H3,1H4,1H5,1H6,1H7,1H8,1H9/
18 DATA BLANK/1H /, COMMA/1H,/
19 C
20 DO 8 I=1,N
21 ARRAY(I) = 0
22 8 CONTINUE
23 C
24 1 READ(TTYIN,200) BUFFER
25 C
26 L = 0
27 POWER = 0
28 C
29 DO 2 I=1,80
30 K = 81 - I
31 IF(BUFFER(K) .EQ. BLANK .OR. BUFFER(K) .EQ. COMMA) GOTO 3
32 C
33 FOUND A CHARACTER, SEE IF IT'S A VALID NUMERIC
34 C
35 DO 4 J=1,10
36 II = J - 1
37 IF(BUFFER(K) .NE. NUMS(J)) GOTO 4
38 C
39 ITS A NUMBER, ADD IT TO PRESENT SUM
40 C
41 ARRAY(L) = ARRAY(L) + (II * (10**POWER))
42 POWER = POWER + 1
43 GOTO 2
44 4 CONTINUE
45 C
46 COME HERE IF CHARACTER FOUND WAS NOT NUMERIC
47 C
48 WRITE(TTYOUT,100)
49 GOTO 1
50 C
51 FOUND A DELIMITER, SEE IF END OF A NUMBER
52 C
53 3 IF(ARRAY(L) .EQ. 0) GOTO 2
54 L = L + 1
55 IF(L .EQ. 0) GOTO 5
56 ARRAY(L) = 0
57 POWER = 0
58 2 CONTINUE

```

```

59 C
60 C MAKE SURE NUMBER(S) IS/ARE LESS THAN 99999 SO
61 C WE DON'T EXCEED 15 FORMATS
62 C
63 5 DO 6 I=1,N
64 IF (ARRAY(I) .GT. 99999) GOTO 7
65 6 CONTINUE
66 RETURN
67 C
68 C ERROR MESSAGE
69 C
70 7 WRITE (TTYOUT,101)
71 GOTO 1
72 C
73 C FORMATS
74 C
75 100 FORMAT(37H-***ERROR*** INPUT NOT NUMERIC--RETRY)
76 101 FORMAT(39H-***ERROR*** NUMBER(S) TOO LARGE--RETRY)
77 200 FORMAT(80A1)
78 END

```

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